

Program Testing with Polymorphic Types

Favonia, 2021.12.06 J.W.W. Zhuyang Wang

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Same program for different kinds of data

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$$\text{map } (*2) [1, 2, 3] = [2, 4, 6]$$

Polymorphism

Same program for different kinds of data

```
map : (a -> b) -> list(a) -> list(b)
```

```
map (not) [true, false] = [false, true]
```

Polymorphism

Same program for different kinds of data

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Question: how can we test this function?

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$$\text{proj} (\text{true}, \text{false}) = ?$$

proj : (a * a * a) -> a

There are only **three** possibilities!

proj₁ (x, y) = x

proj₂ (x, y) = y

proj (true, false) = ?

$$\text{proj} : (a * a * a) \rightarrow a$$

There are only **three** possibilities!

$$\text{proj}_1 (x, y, z) = x$$

$$\text{proj}_2 (x, y, z) = y$$

$$\text{proj}_3 (x, y, z) = z$$

$$\text{proj} : (a * a * a) \rightarrow a$$

There are only **three** possibilities!

$$\text{proj}_1 (x, y, z) = x$$

$$\text{proj}_2 (x, y, z) = y$$

$$\text{proj}_3 (x, y, z) = z$$

$$\text{proj} (0, 1, 2) = ?$$

$$\text{proj} : (a * a * \dots * a) \rightarrow a$$

There are exactly **n** possibilities!

proj : $(a * a * \dots * a) \rightarrow a$

There are exactly n possibilities!

proj $(0, 1, 2, \dots, n-1) = ?$

$$\text{proj} : (a * a * \dots * a) \rightarrow a$$

There are exactly **n** possibilities!

$$\text{proj} (\emptyset, 1, 2, \dots, n-1) = ?$$

Any type with **n** distinct values works

$$\text{proj} : (a * a * \dots * a) \rightarrow a$$

There are exactly **n** possibilities!

$$\log_a(a * a * \dots * a) = \log_a(a^n) = n$$

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counting elements of type **a** in the input

$$f : \alpha(a) \rightarrow H(a)$$

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one of the best types for testing*

$$\log_a \alpha(a)$$

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testing is still complete**

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** the empty type might need to be separately tested

2018 **I joined the U**

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- 2021** **Paper published at POPL 2022**
“Logarithm and Program Testing”

2010 *“Testing Polymorphic Properties”*
by Bernardy, Jansson, and Claessen

f : **(F(a) -> a) * (G(a) -> K) -> H(a)**
(other cases manually massaged into this form)

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Logarithm-like operations have been discovered repeatedly
in the literature but no one connected logarithm to testing

e.g., [Abbott et al. 2003; Altenkirch et al. 2015]

$f : (a \rightarrow a) * a \rightarrow a$

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Suppose the input is the pair (s, x)

The output must be $s(s(\dots s(x)\dots))$

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The output must be $s(s(\dots s(x)\dots))$

$$f ((+1), \emptyset) = ?$$

this reveals the number of s
in its output $s(s(\dots s(x)\dots))$

$f : (a \rightarrow a) * a \rightarrow a$

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$$\begin{aligned} & \log_a(a^a * a) \\ &= \log_a(a^a) + \log_a(a) \\ &= a * \log_a(a) + 1 \\ &= a * 1 + 1 \\ &\simeq a + 1 \end{aligned}$$

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Problem: a appears in the logarithm of $a^a * a$ indicating recursion

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Solution: recursive types!
 $\mu a.a+1$ is the naturals \mathbb{N}

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Problem: a appears in the logarithm of $a^a * a$ indicating recursion

Solution: recursive types!
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Sufficient for $f ((+1), \emptyset) = ?$

$$f : \alpha(a) \rightarrow H(a)$$

one of the best types for testing*

$$\mu a. \log_a \alpha(a)$$

a sufficiently large type to describe all possible ways to generate an a -element

* see our paper for remaining caveats

Theorem*

For any two functions $f, g : \alpha(a) \rightarrow H(a)$, if they agree on all inputs when type a is instantiated with $\mu a. \log_a \alpha(a)$ and with the empty type**, then they are exactly the same function!

* see our paper for omitted conditions

** see our paper for why the empty type is needed

Logarithm for All

$$\log_a(a) = 1$$

$$\log_a(1) = 0$$

$$\log_a(\alpha_1(a) * \alpha_2(a)) = \log_a(\alpha_1(a)) + \log_a(\alpha_2(a))$$

$$\log_a(\alpha_2(a)^{\alpha_1(a)}) = \alpha_1(a) * \log_a(\alpha_2(a))$$

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$$\log_a(b) = \log_a(0) = 0$$

$$\log_a(\alpha_1(a) + \alpha_2(a)) = \log_a(\alpha_1(a)) + \log_a(\alpha_2(a))$$

$$\log_a(\mu b \dots) = \dots$$

Use over-approximation to cover more types

$$\text{proj} : (a * a) \rightarrow a$$

Theorem: instantiate a with $\mu a. \log_a(a*a) = \mu a. 2 \simeq 2$

proj : (a * a) -> a

Theorem: instantiate **a** with $\mu a. \log_a(a*a) = \mu a. 2 \simeq 2$

proj (true, true) = ?

proj (true, false) = ?

proj (false, false) = ?

proj (false, true) = ?

$$\text{proj} : (a * a) \rightarrow a$$

Theorem: instantiate a with $\mu a.\log_a(a*a) = \mu a.2 \simeq 2$

~~$\text{proj}(\text{true}, \text{true}) = ?$~~

$$\text{proj}(\text{true}, \text{false}) = ?$$

~~$\text{proj}(\text{false}, \text{false}) = ?$~~

$$\text{proj}(\text{false}, \text{true}) = ?$$

Recall that logarithm is large enough to index all a -elements. The best tests are those containing only distinct a -elements.

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one of the best types and inputs

$a = b = \mathbb{N}$

$\text{map } \text{id}_{\mathbb{N}} [0, 1, \dots, n-1] = ?$

intuition: an evil programmer can only do these three things:
duplicating, omitting, or permuting elements in the output list

the above test case detects all possible deviations
 $\text{map id } [1,1,1,1]$ in comparison is much less useful

Theorem 2.0*

For any two functions $f, g : \alpha(a) \rightarrow H(a)$, if they agree on **optimal inputs with distinct a-elements**** when type a is instantiated with $\mu a.\log_a \alpha(a)$ and on all inputs with the empty type***, then they are exactly the same function!

* see our paper for omitted conditions

** see our paper for the precise definition

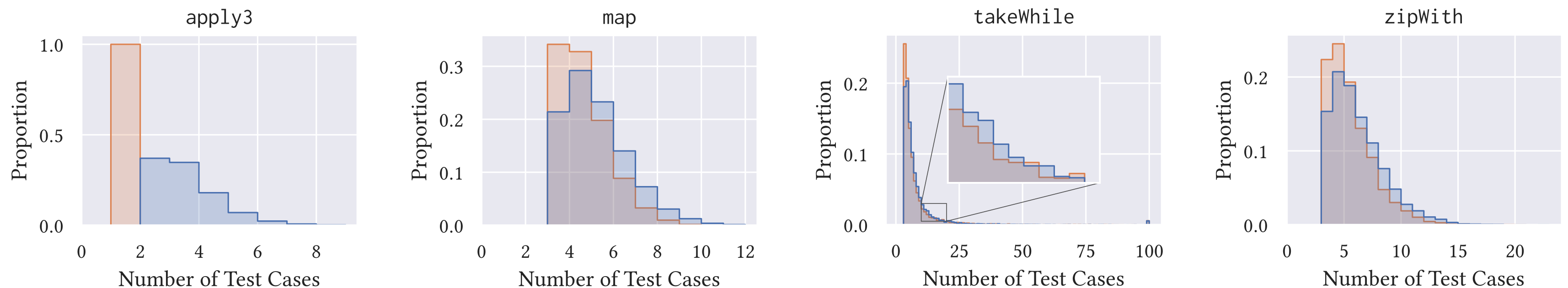
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Implementation

Haskell library: github.com/hawnzug/polycheck
automatically specializing types and inputs

Can work with either QuickCheck or SmallCheck

Implementation



Implementation

-  POLYCHECK with QUICKCHECK
-  Original QUICKCHECK

Preliminary experiments showed it requires fewer test cases to find counterexamples

Future Work

1. Incorporate information other than the typing
e.g., `sort cmp list` expects `cmp` to form a total order

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e.g., `sort cmp list` expects `cmp` to form a total order
2. Extend work to test an API, not just one function
3. Further optimize the theorem
e.g., for `length : list(a) -> int`
the best choice is the unit type, not natural numbers