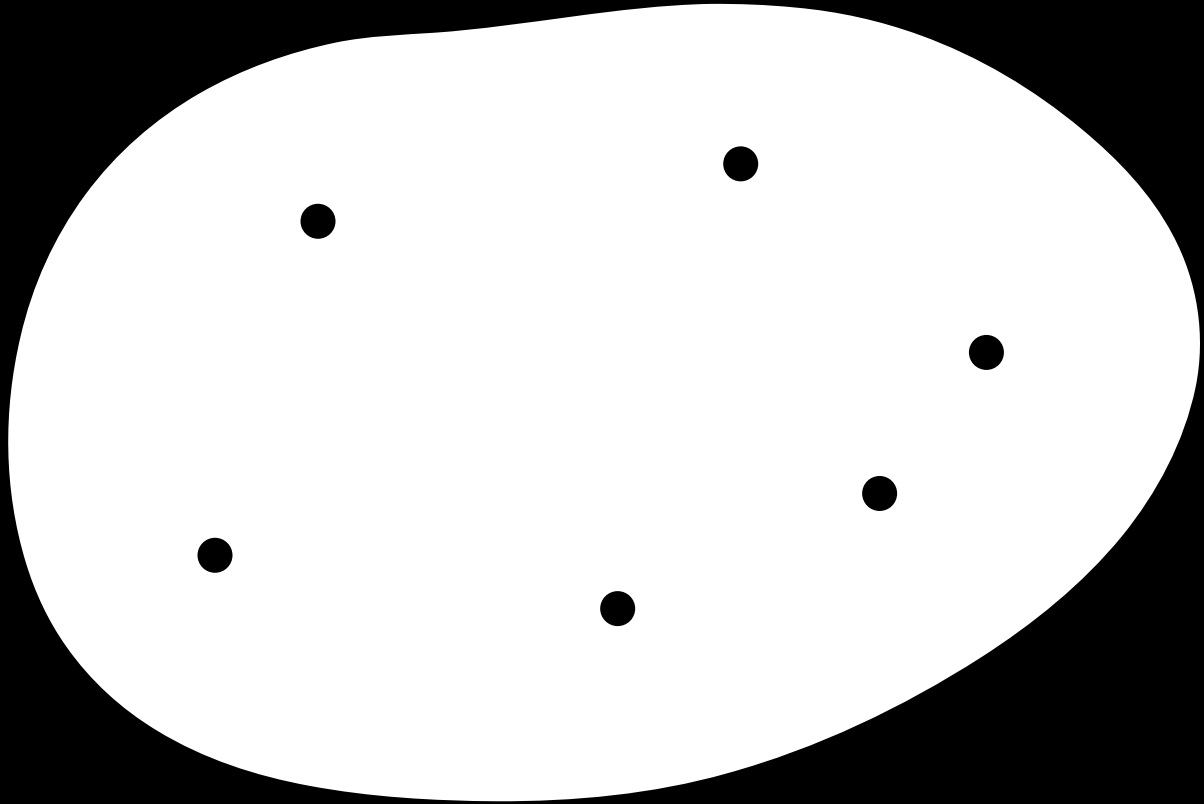
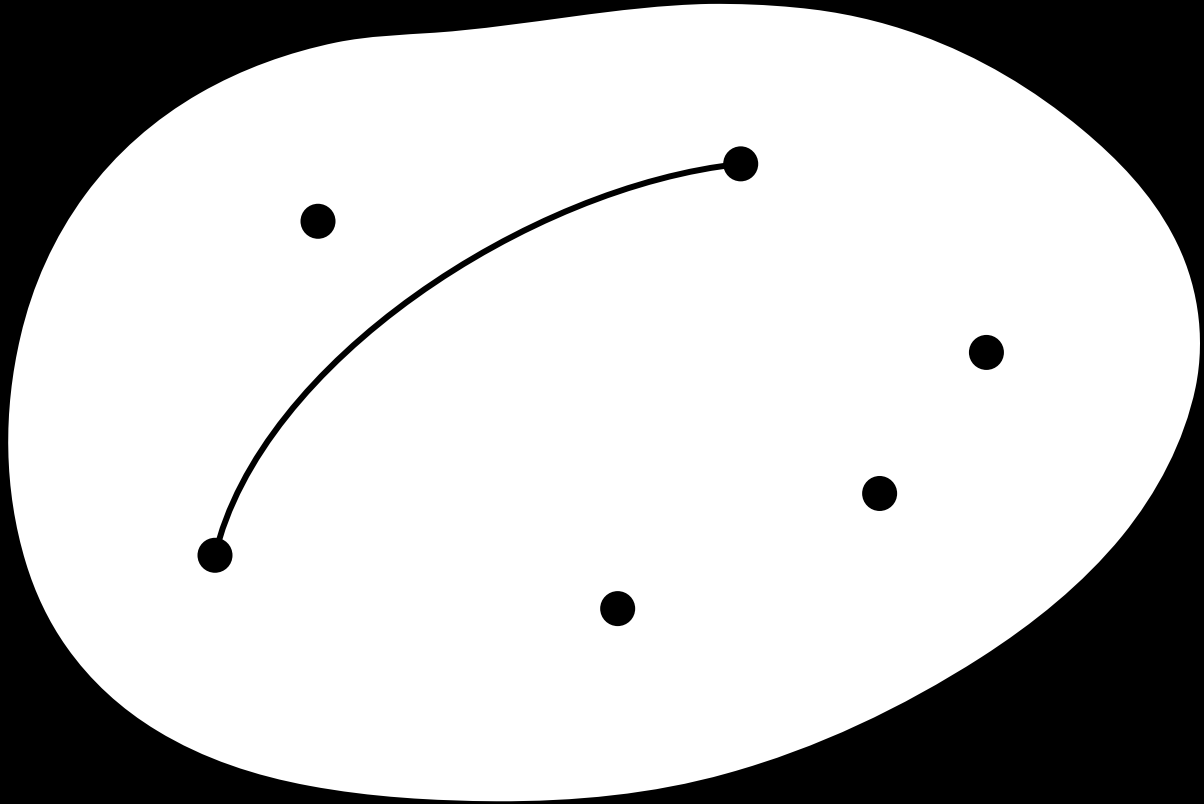
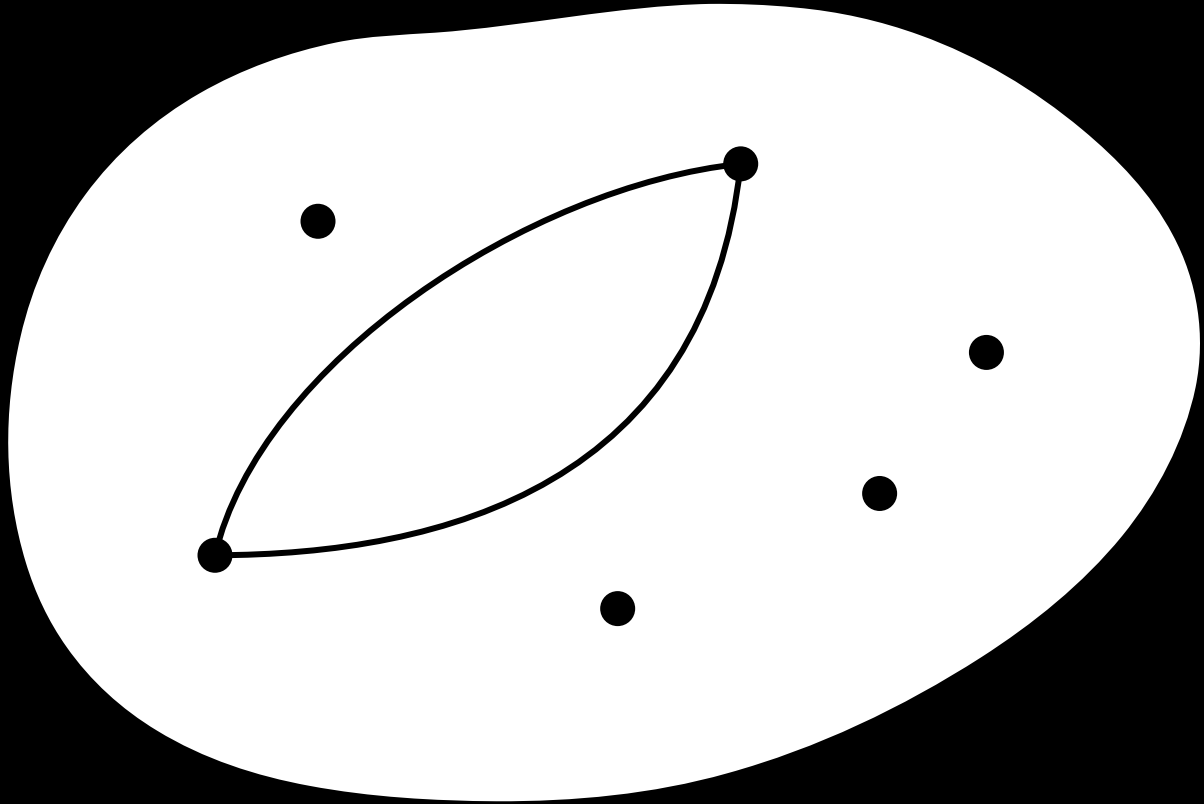


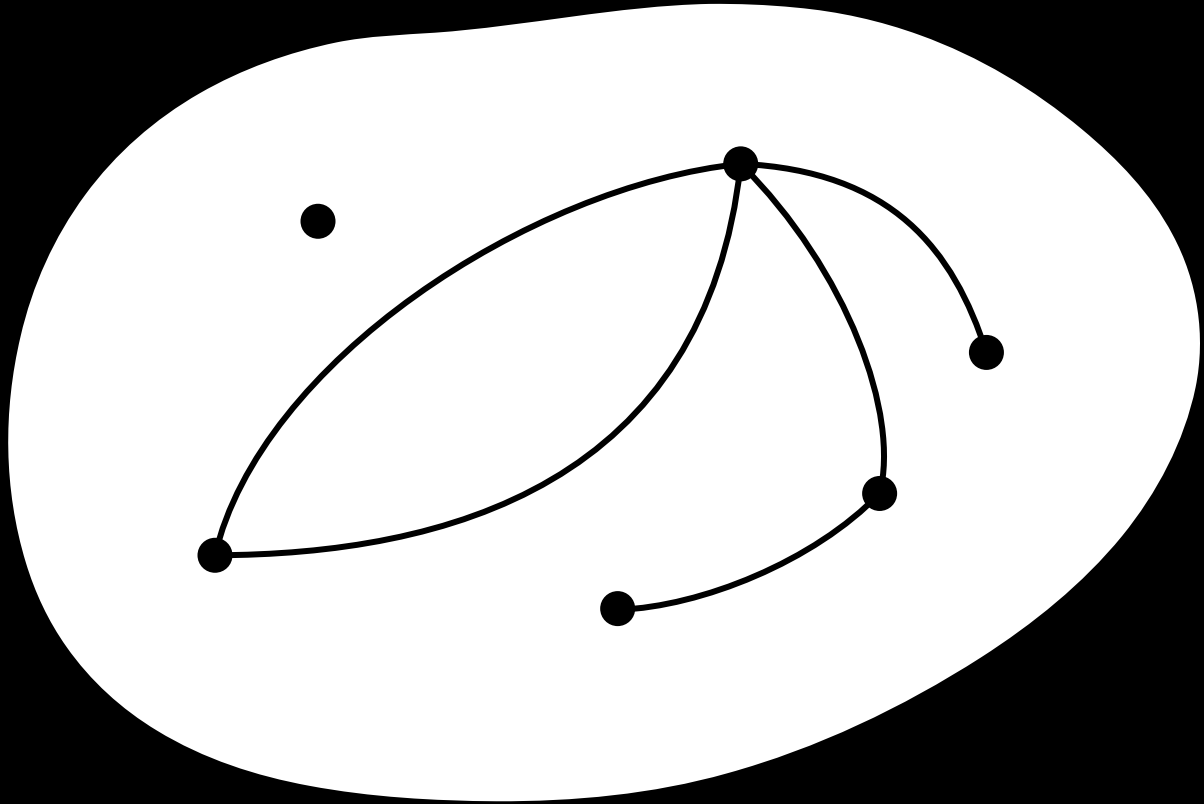
\mathbb{N}

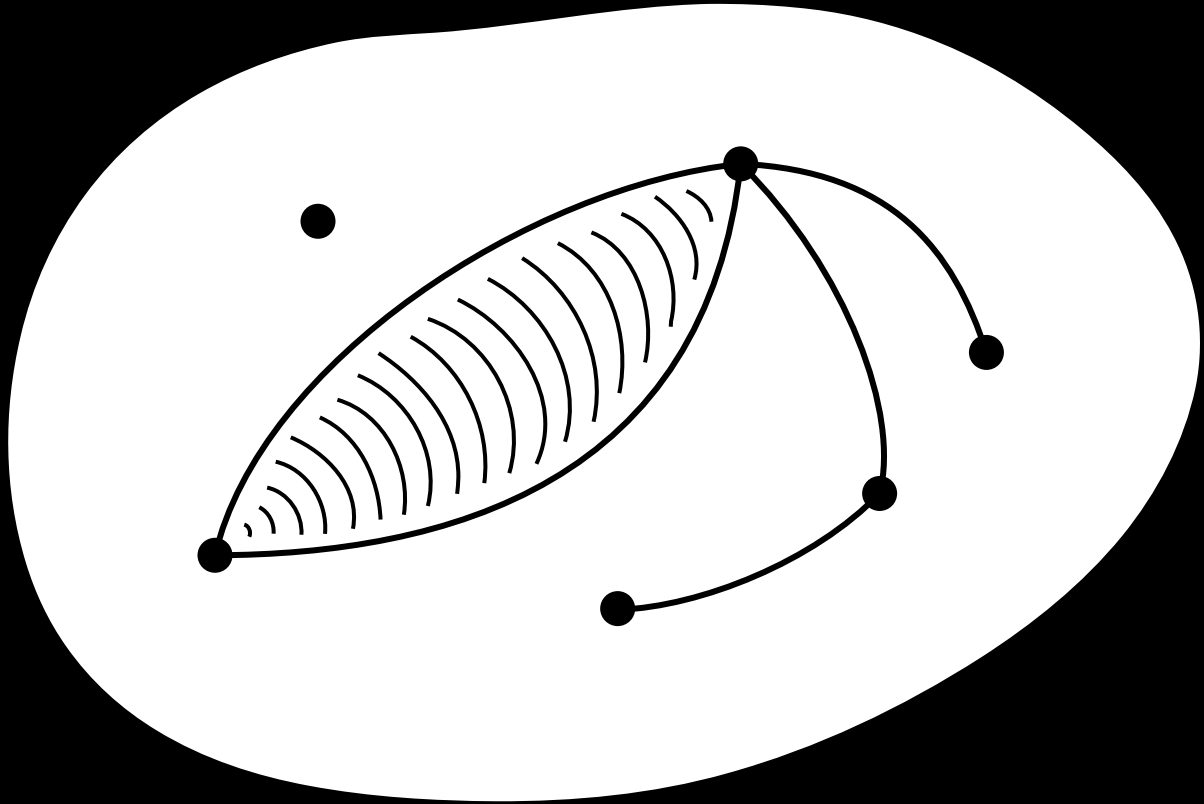
$\mathbb{N} \rightarrow \mathbb{N}$

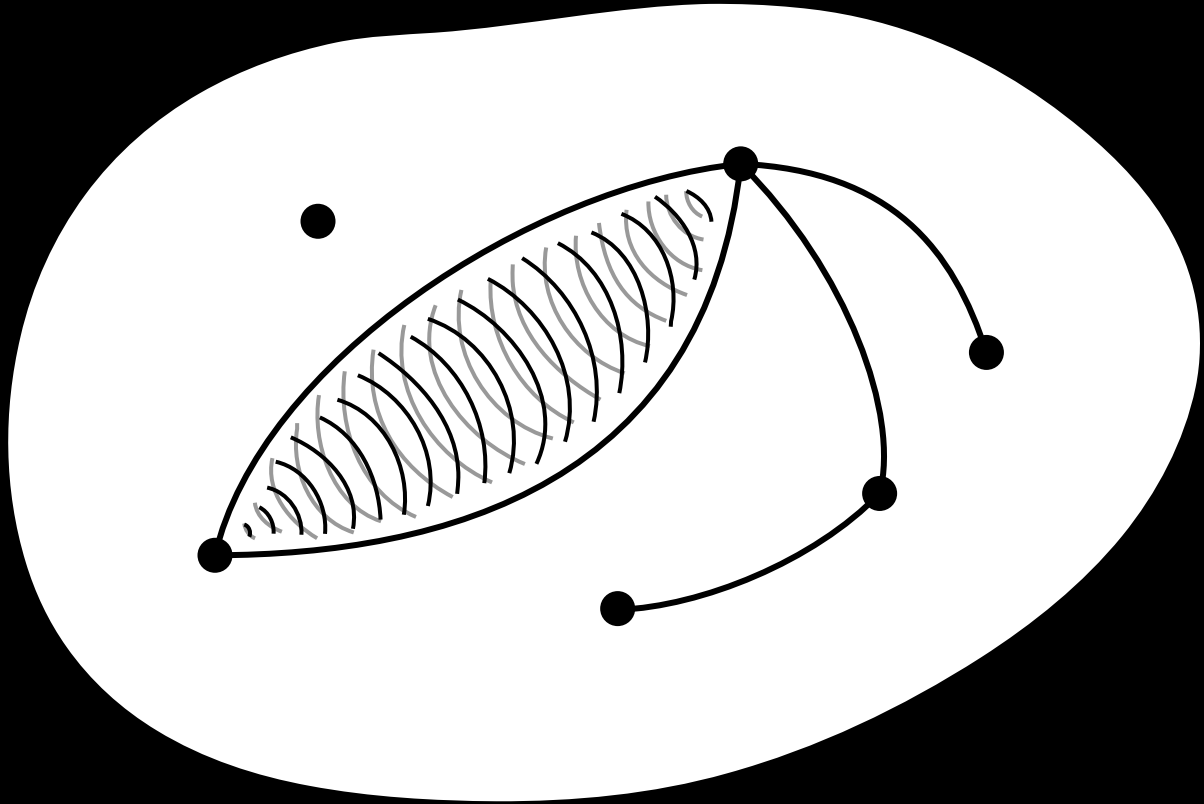


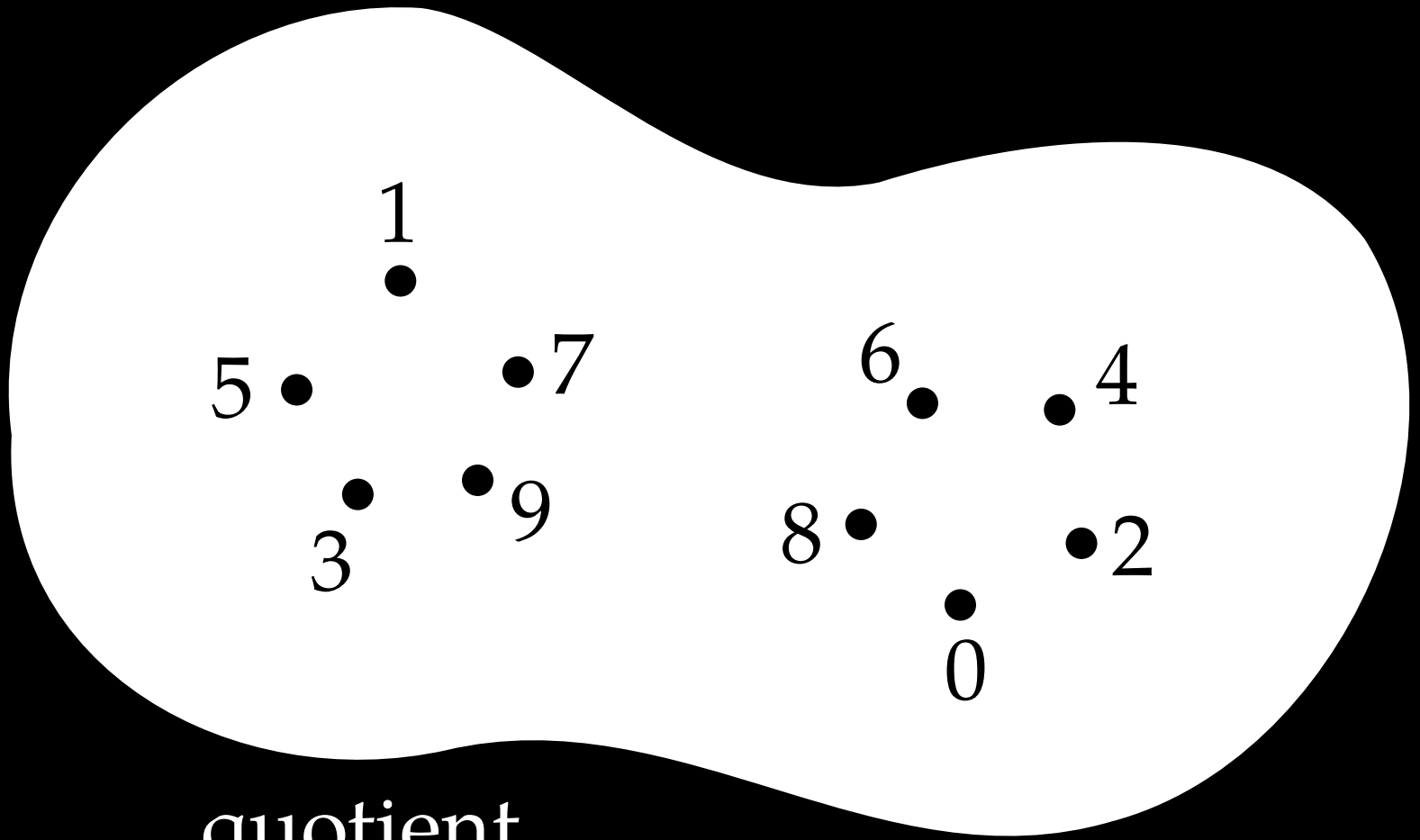




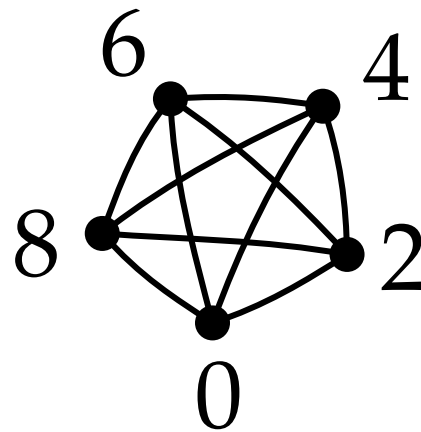
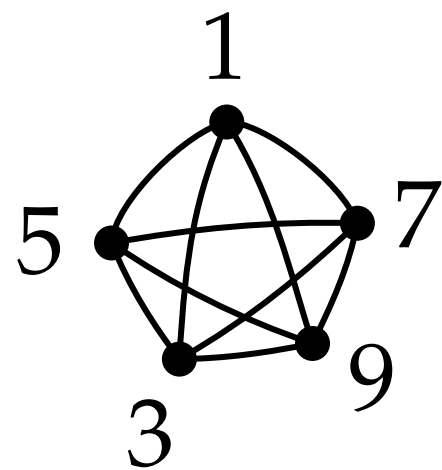




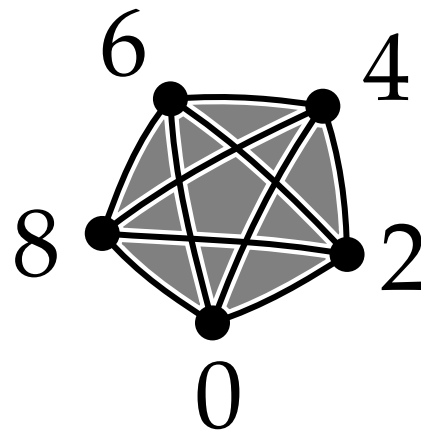
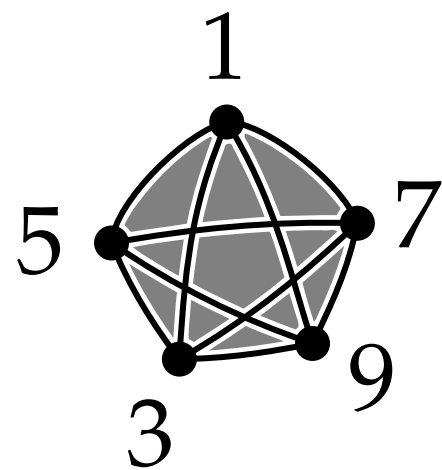




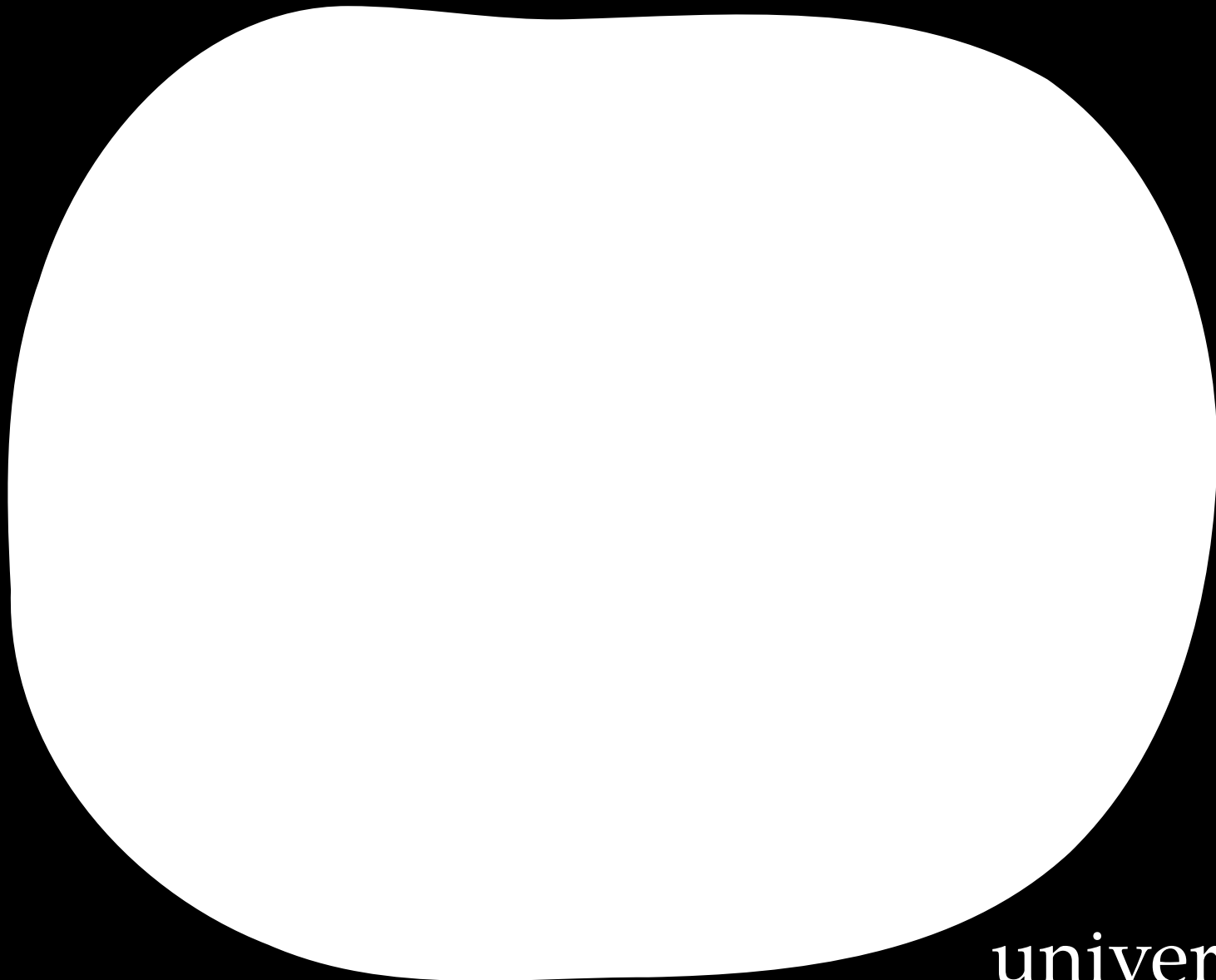
quotient



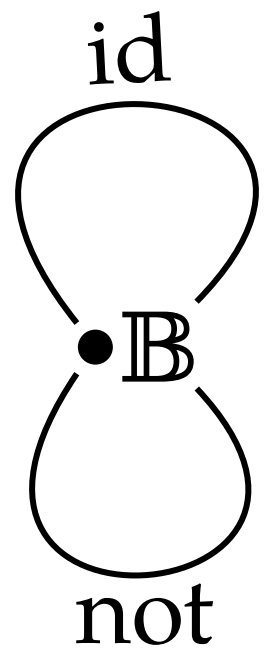
quotient



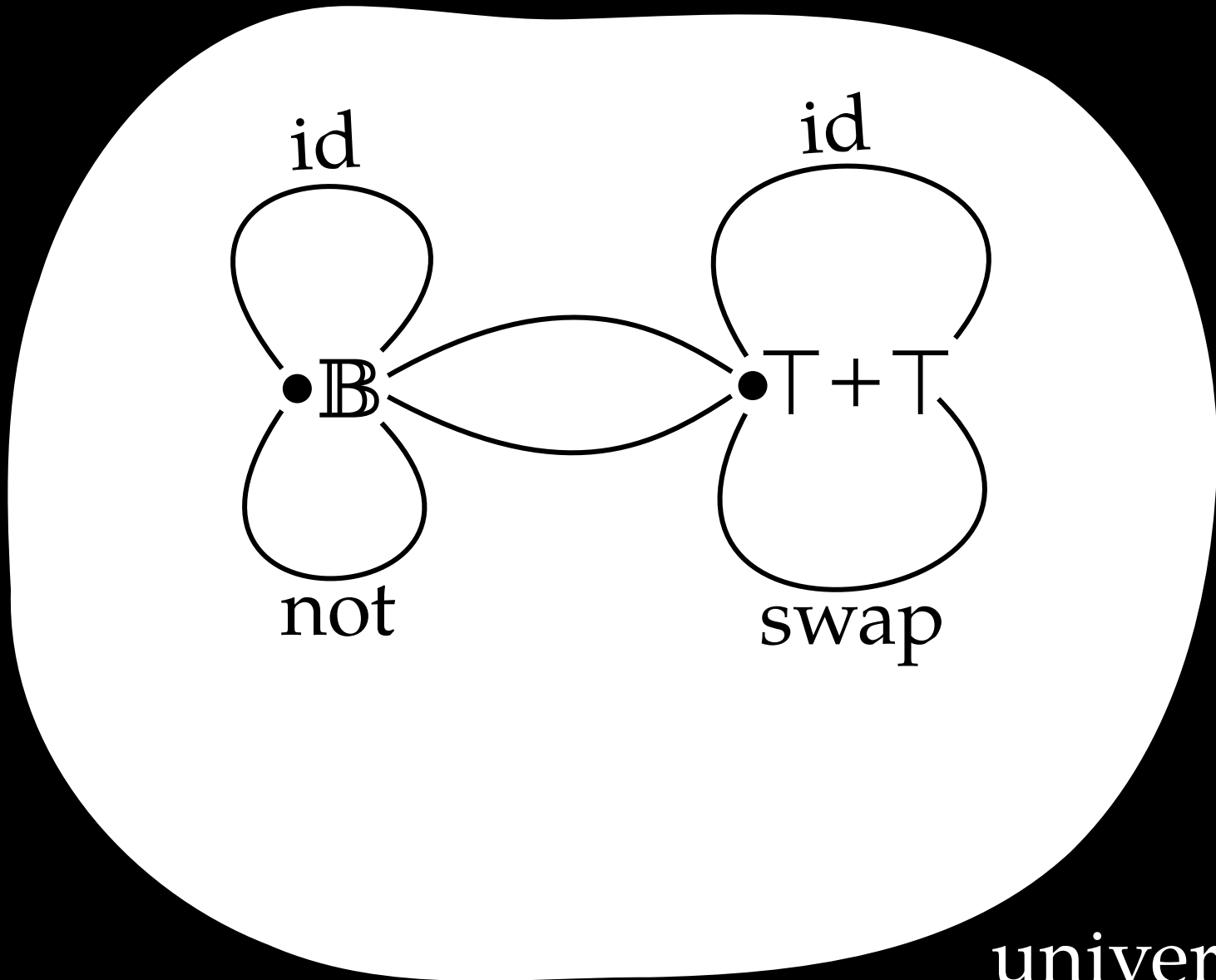
quotient



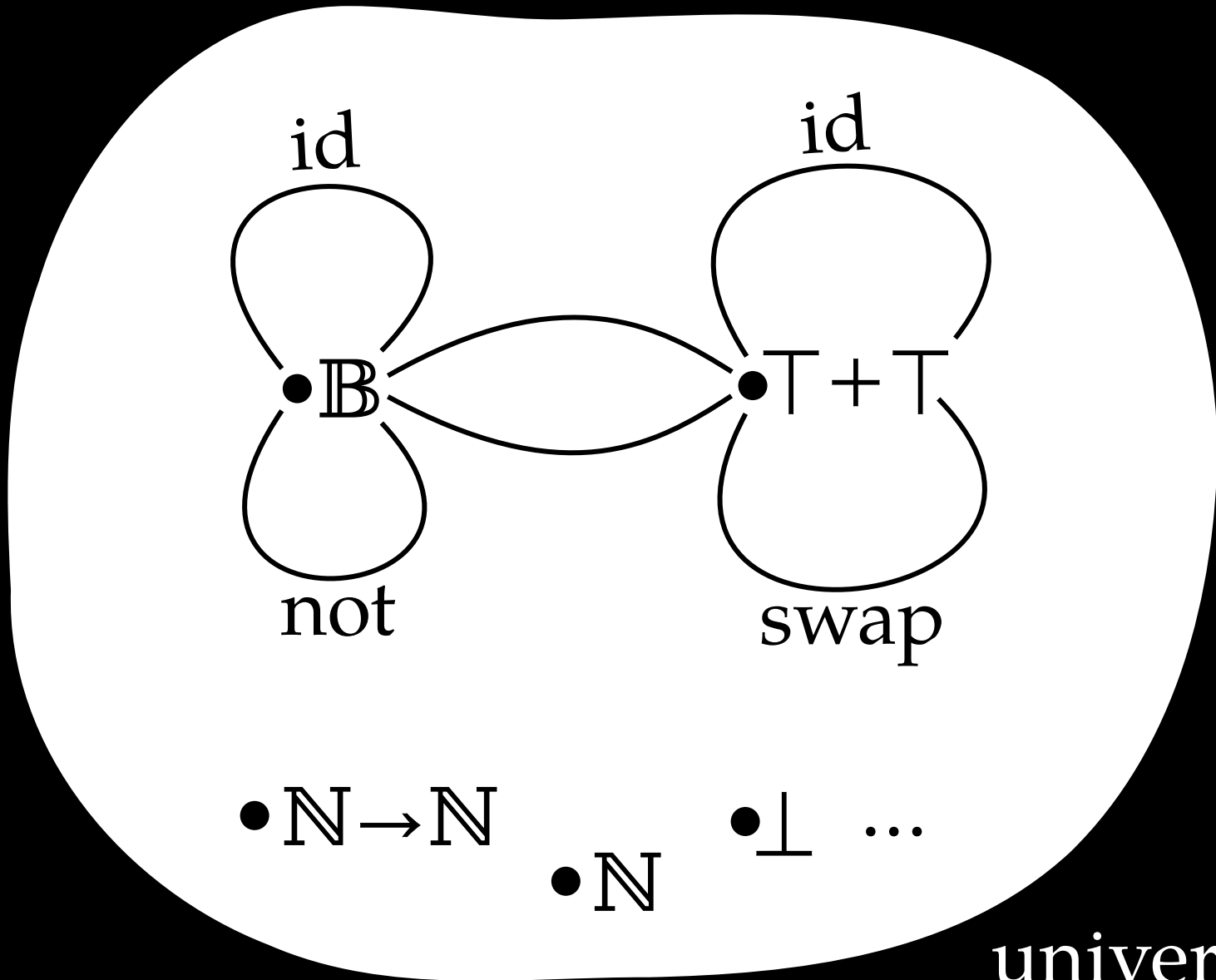
universe

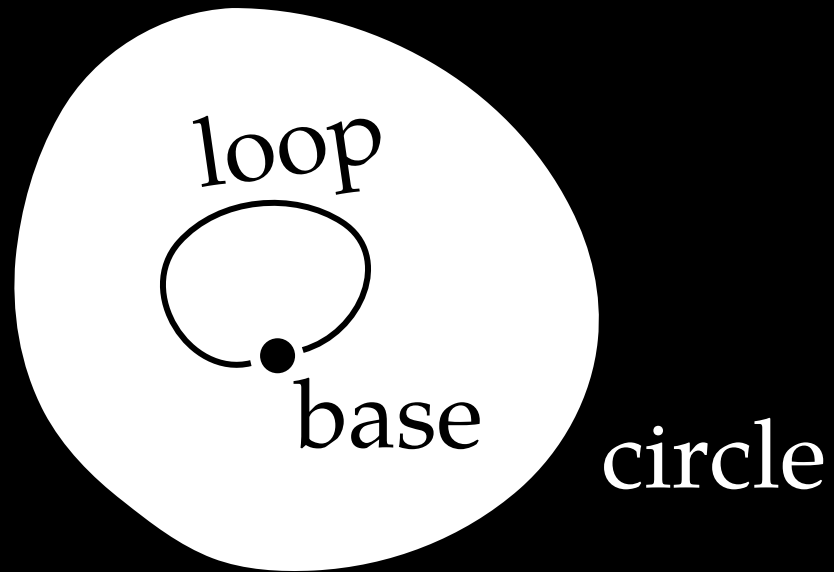


universe



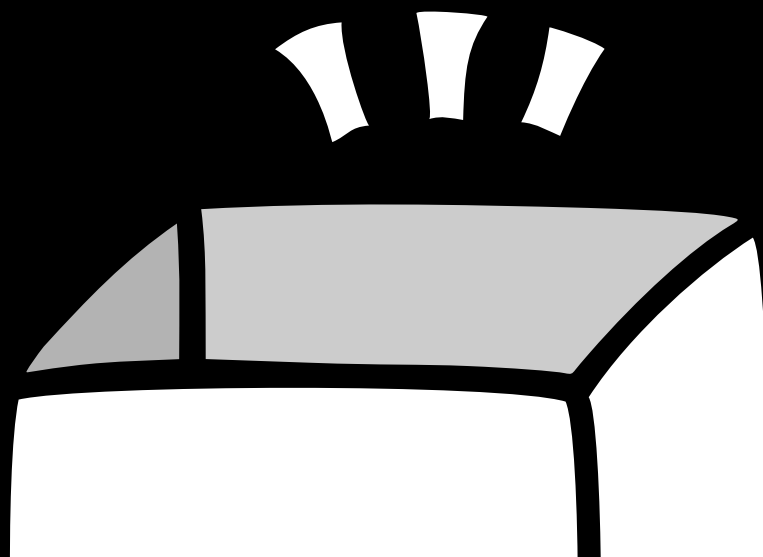
universe

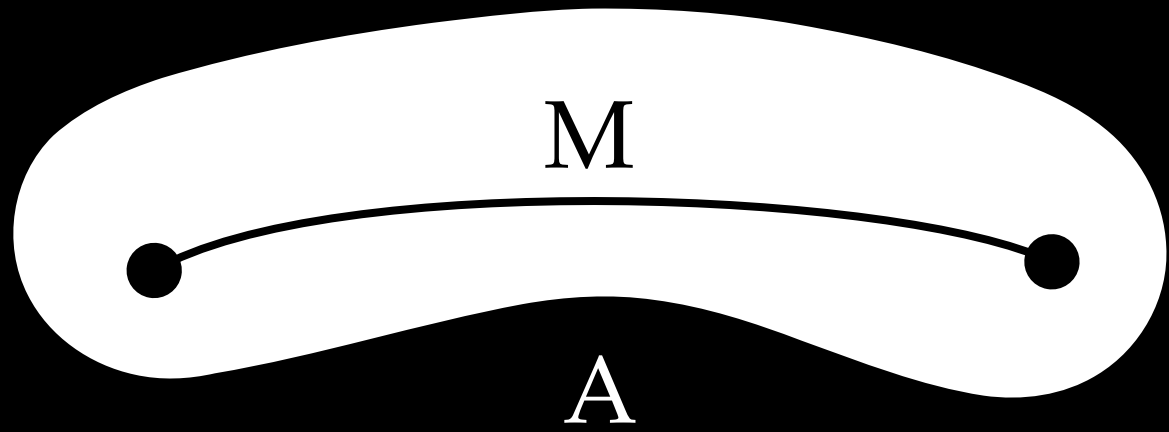


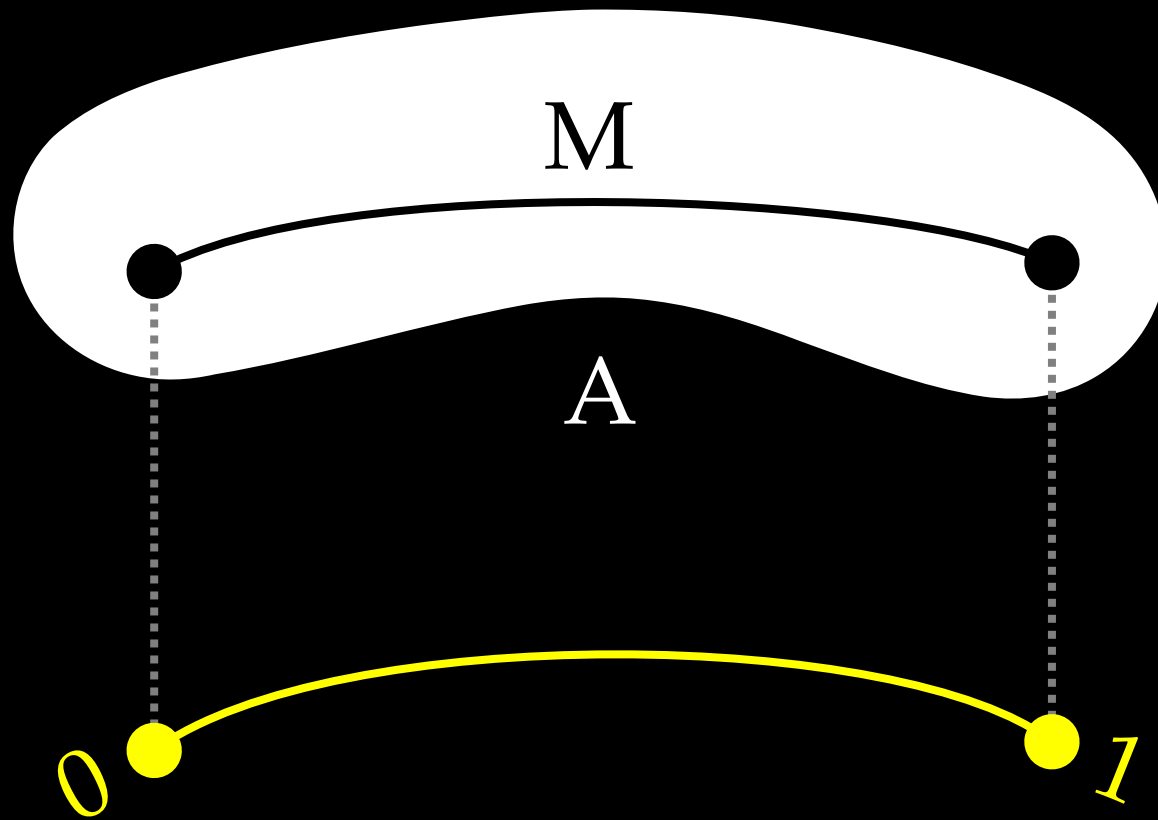


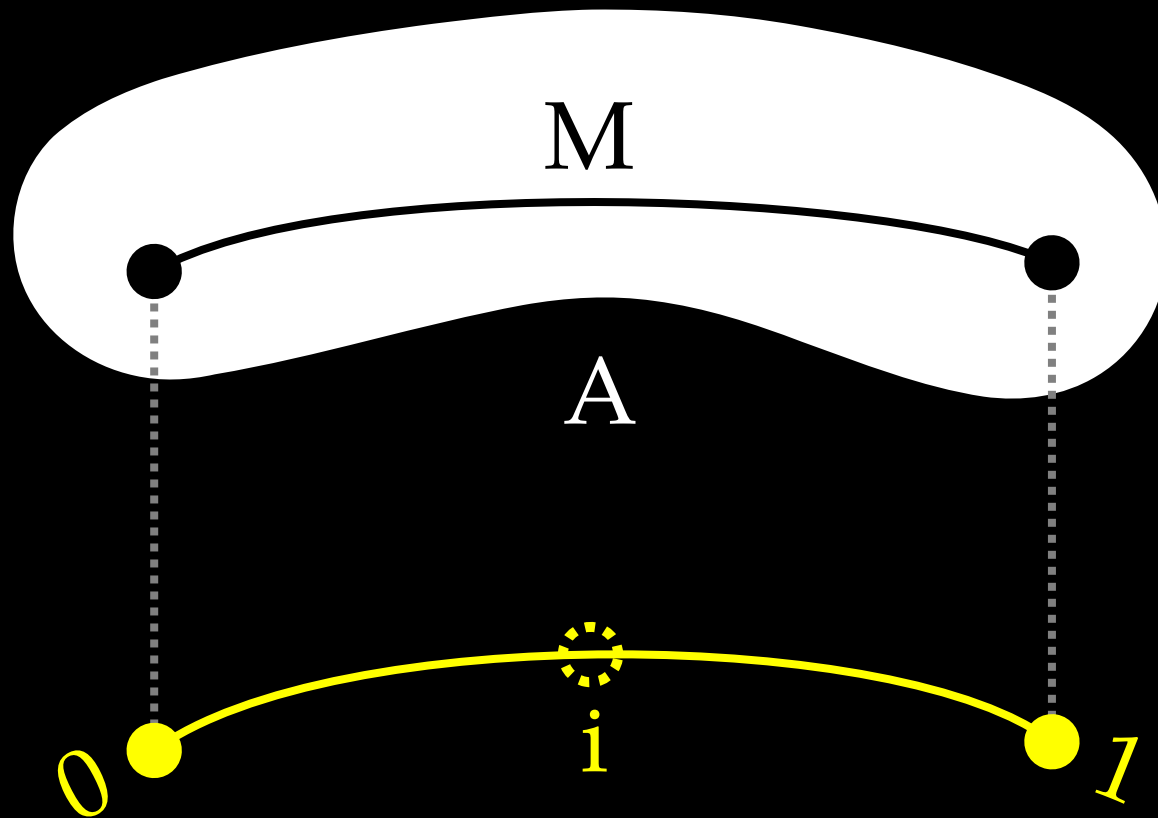
(homotopy theory)

I. Cubes

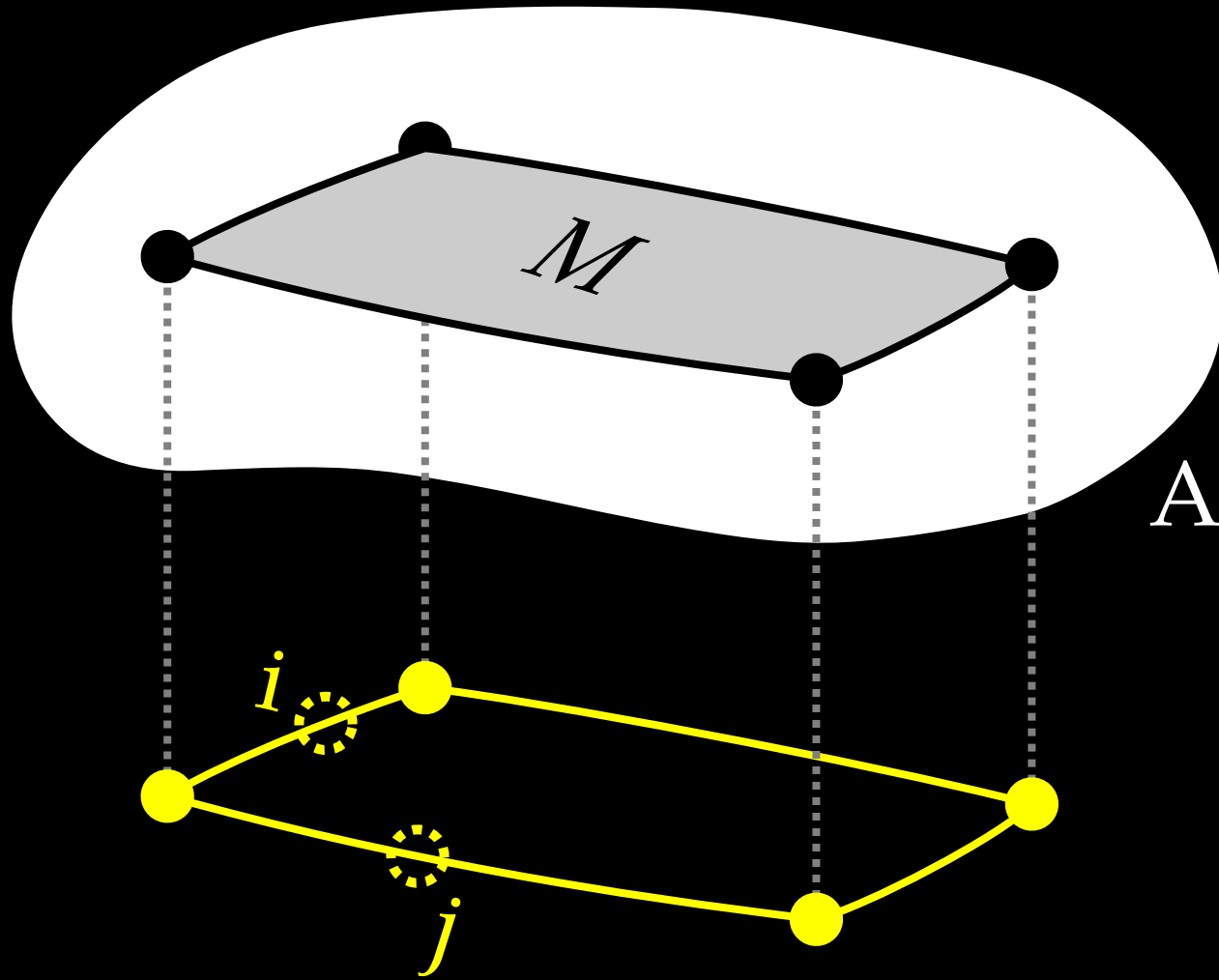




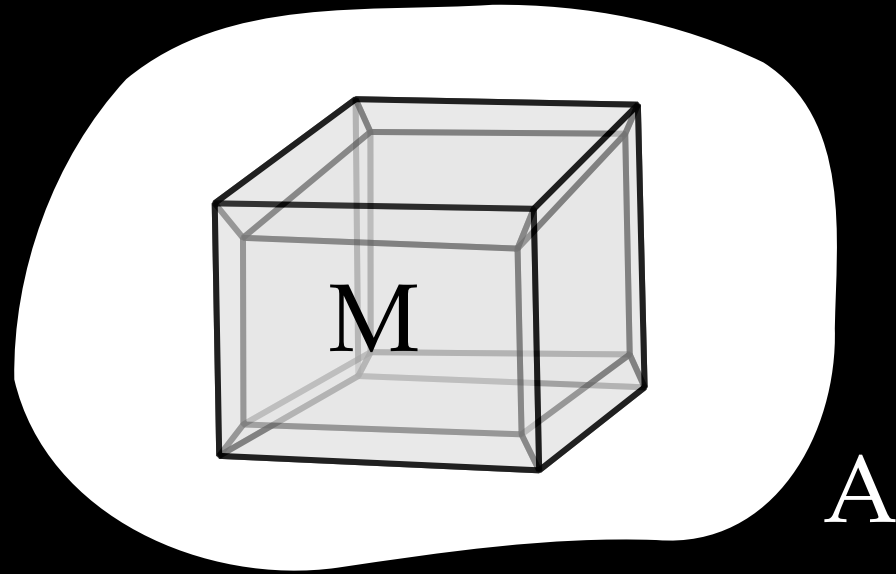




$$i:\mathbb{I} \vdash M : A$$

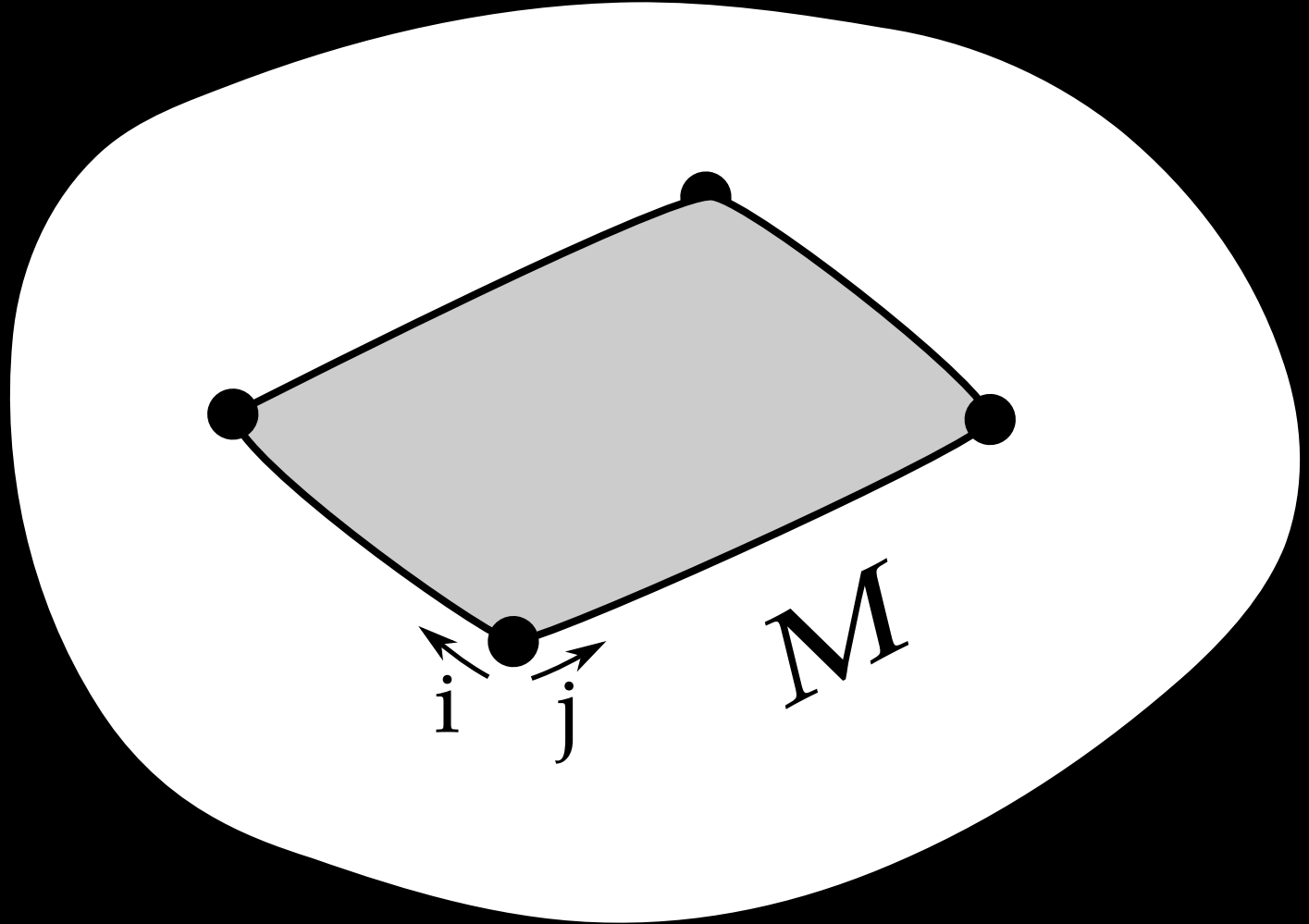


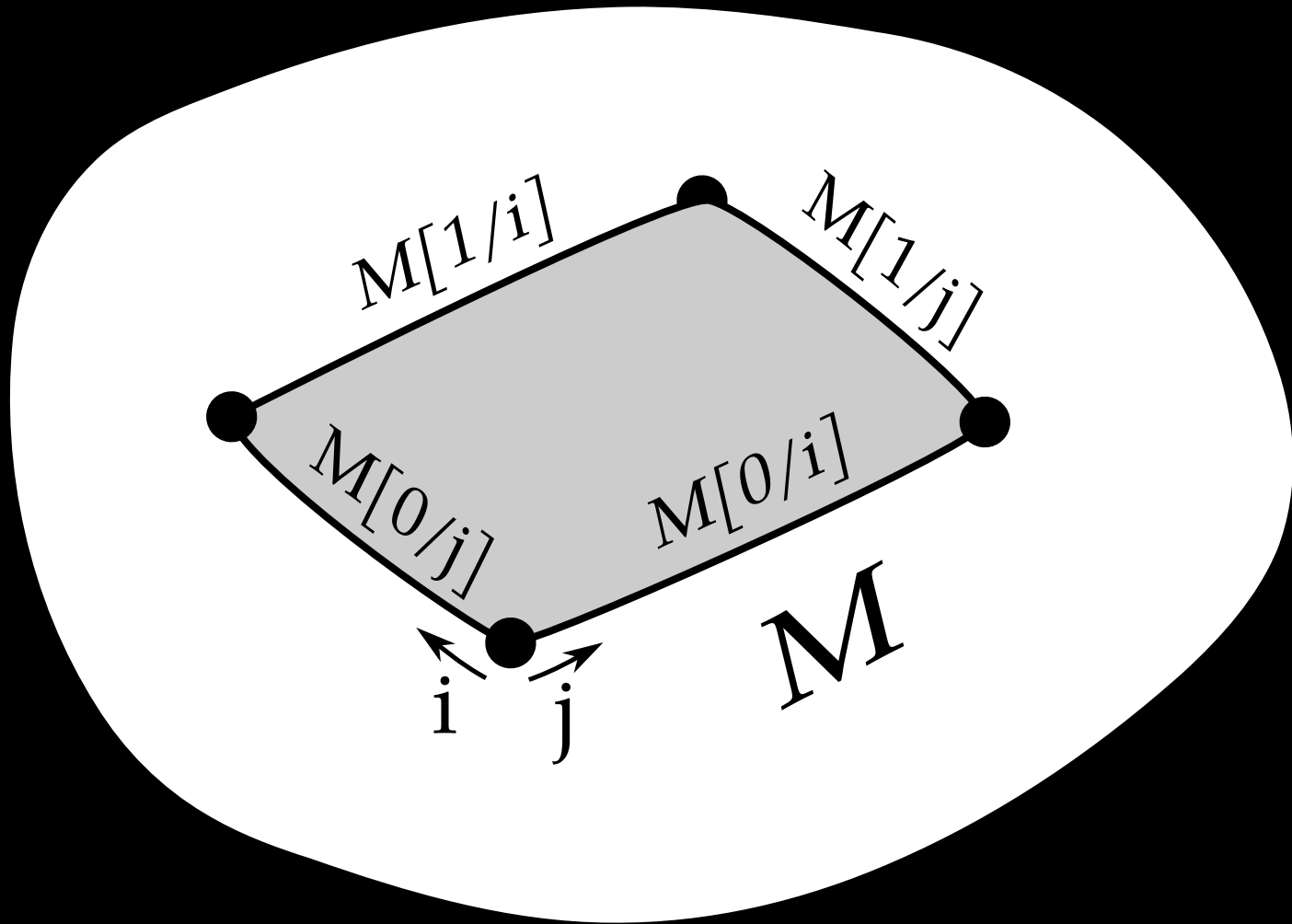
$$i:\mathbb{I}, j:\mathbb{I} \vdash M : A$$

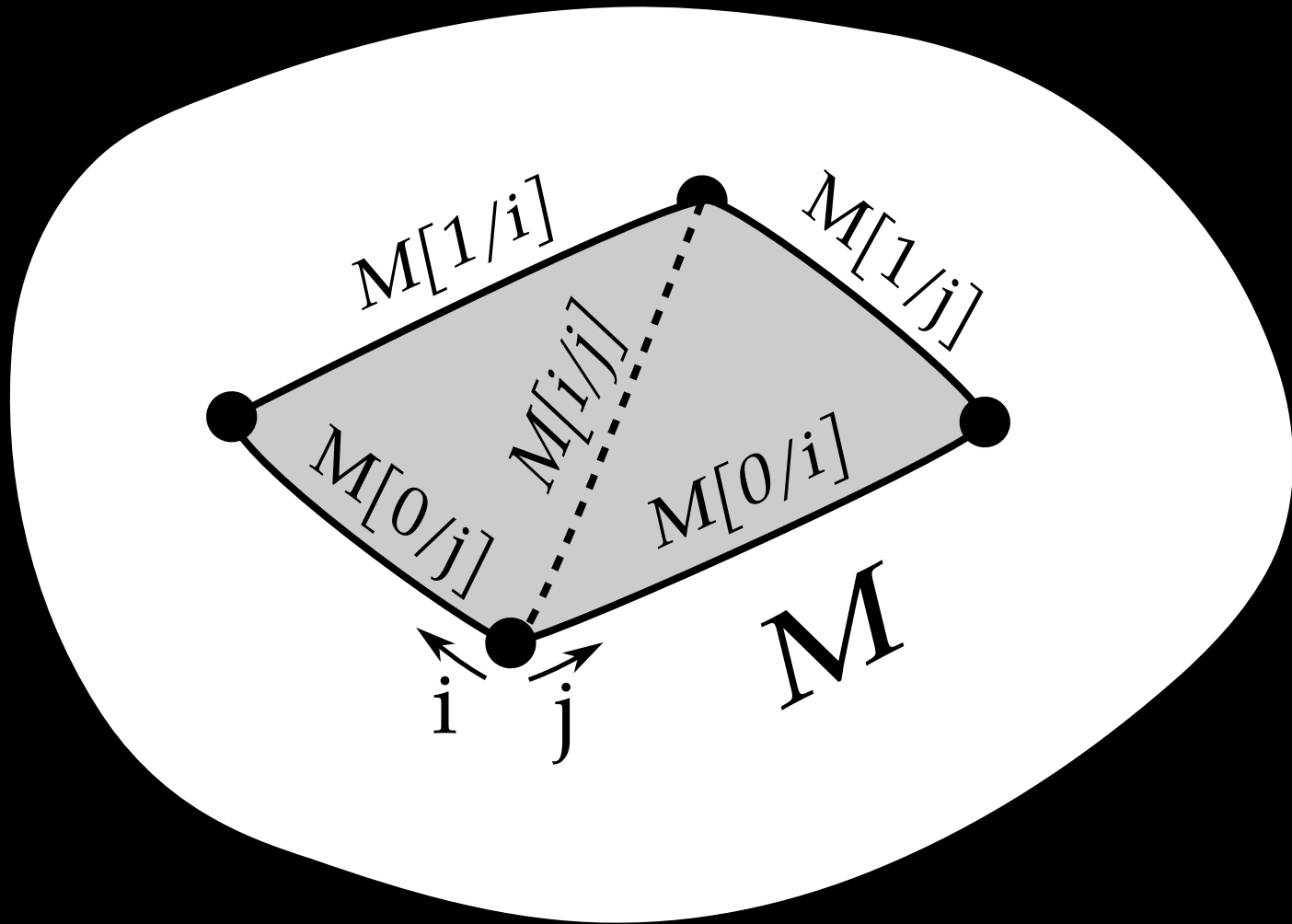


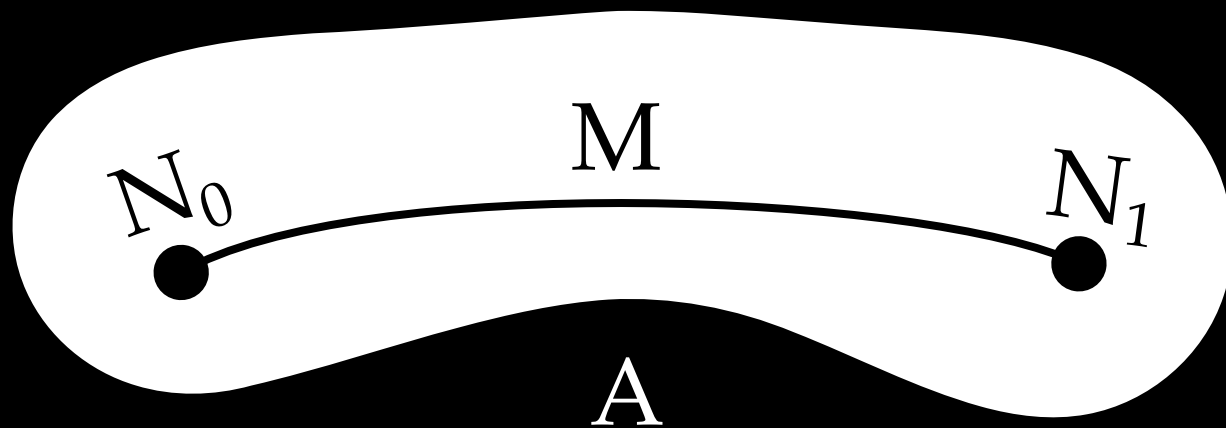
n-cube

$$i_1:\mathbb{I}, i_2:\mathbb{I}, \dots, i_n:\mathbb{I} \vdash M : A$$

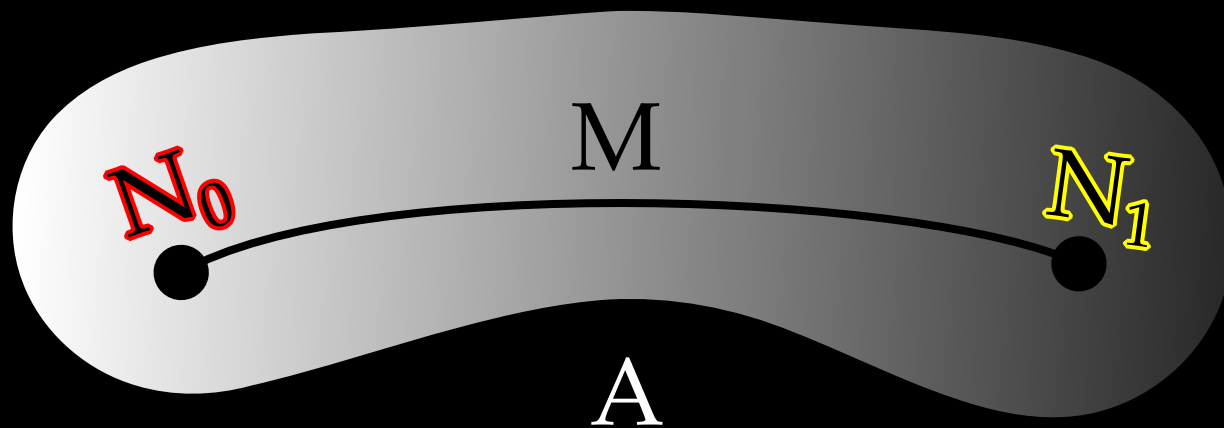






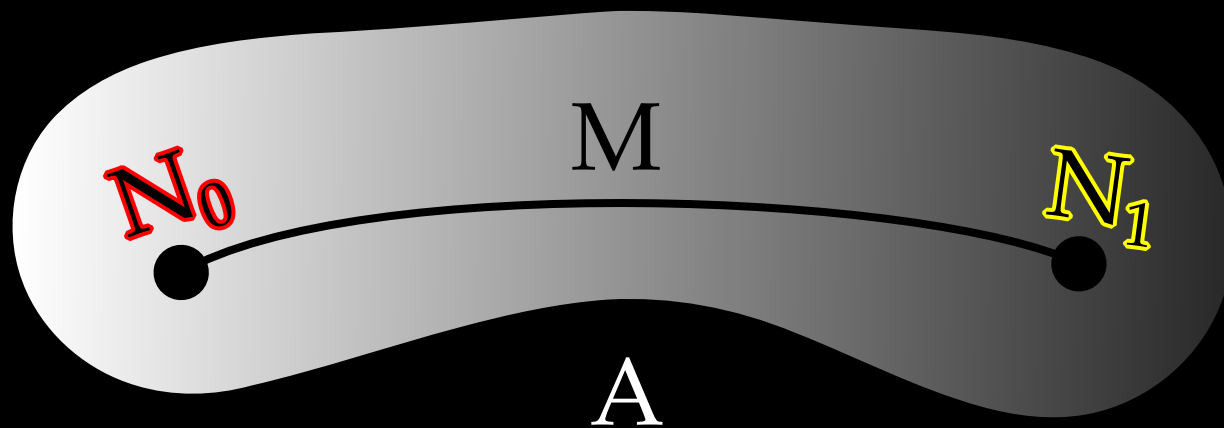


path types
functions from \mathbb{I}



$$i:\mathbb{I} \vdash M : A \quad \begin{array}{l} M[0/i] \equiv N_0 : A[0/i] \\ M[1/i] \equiv N_1 : A[1/i] \end{array}$$

$$\lambda i.M : \text{Path}_{i,A}(N_0, N_1)$$



$$i:\mathbb{I} \vdash M : A \quad \begin{array}{l} M[0/i] \equiv N_0 : A[0/i] \\ M[1/i] \equiv N_1 : A[1/i] \end{array}$$

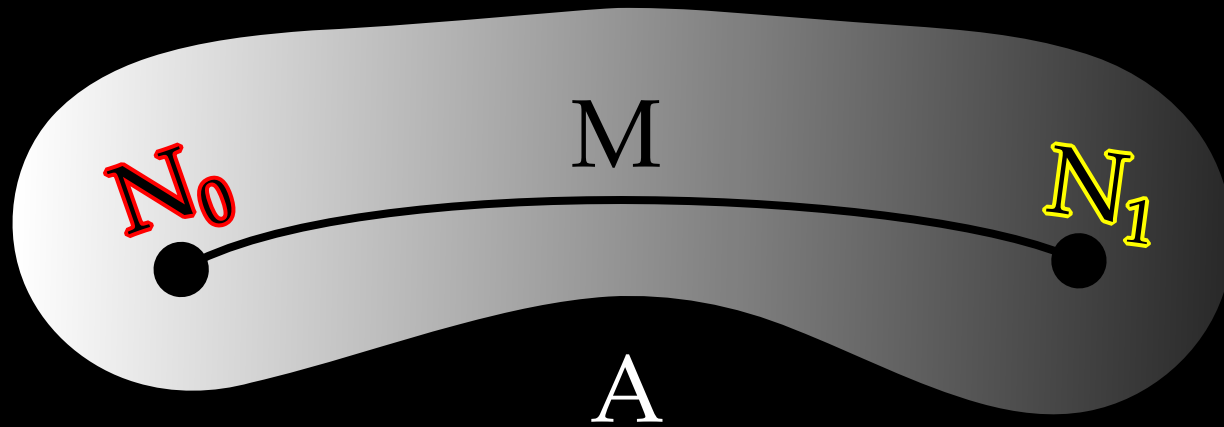
$$\lambda i.M : \text{Path}_{i.A}(N_0, N_1)$$

$$P : \text{Path}_{i.A}(N_0, N_1) \quad r:\mathbb{I}$$

$$P@r : A[r/i]$$

$$P@0 \equiv N_0 : A[0/i]$$

$$P@1 \equiv N_1 : A[1/i]$$



$$i:\mathbb{I} \vdash M : A \quad \begin{array}{l} M[0/i] \equiv N_0 : A[0/i] \\ M[1/i] \equiv N_1 : A[1/i] \end{array}$$

$$\lambda i.M : \text{Path}_{i.A}(N_0, N_1)$$

$$P : \text{Path}_{i.A}(N_0, N_1) \quad r:\mathbb{I}$$

$$P@r : A[r/i]$$

$$P@0 \equiv N_0 : A[0/i]$$

$$P@1 \equiv N_1 : A[1/i]$$

$$(\lambda i.M)@r \equiv M[r/i] : A[r/i]$$

$$P \equiv \lambda i.P@i : \text{Path}_{i.A}(N_0, N_1)$$

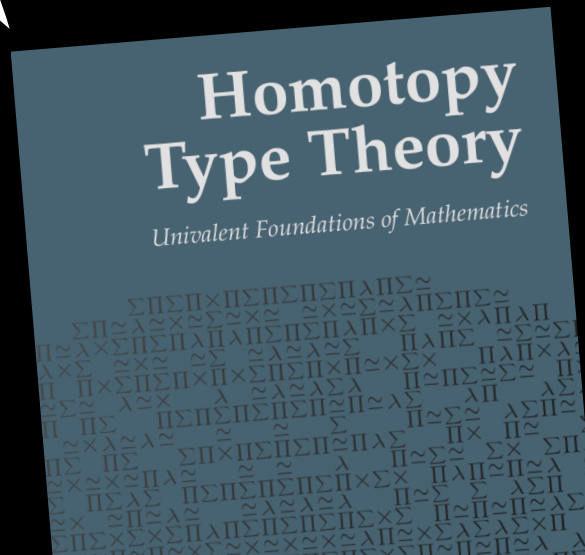
function extensionality

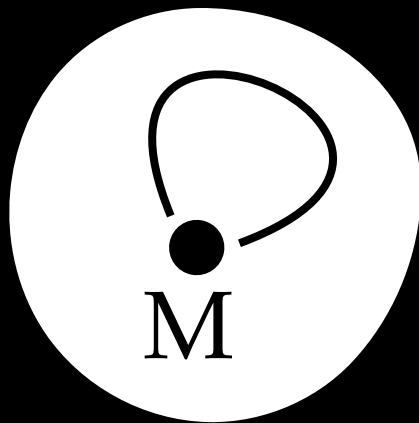
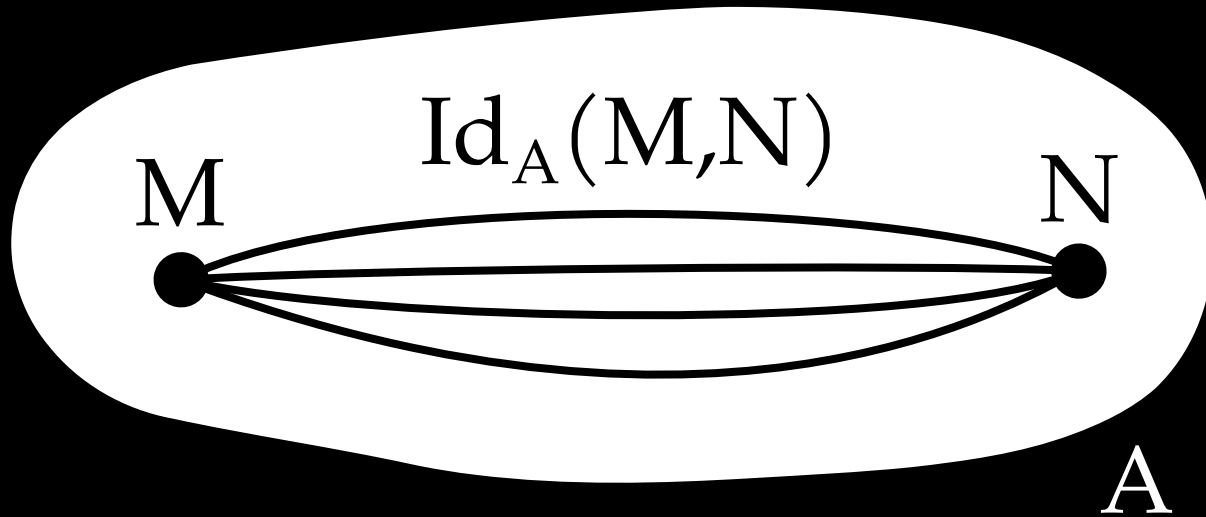
almost trivial

$\mathbf{h} : \prod(x:A). \text{Path}(F(x), G(x))$

$\lambda i. \lambda x. \mathbf{h}(x) @ i : \text{Path}(F, G)$

II. the Book





$\text{refl}_M : \text{Id}_A(M, M)$
the only generator

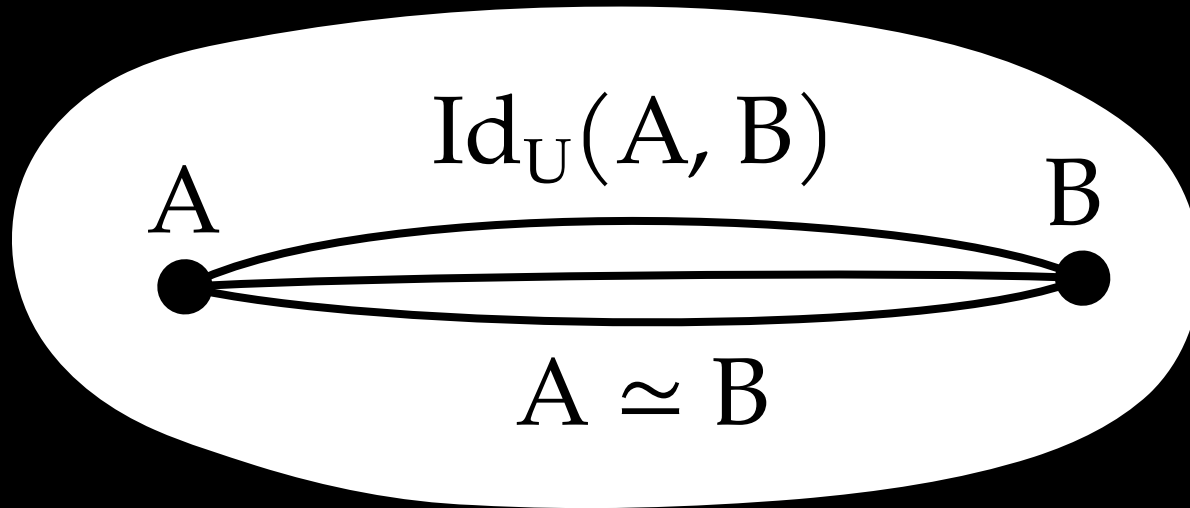
zero ←→ base case
suc ←→ induction step

mathematical induction

refl ←→ refl case

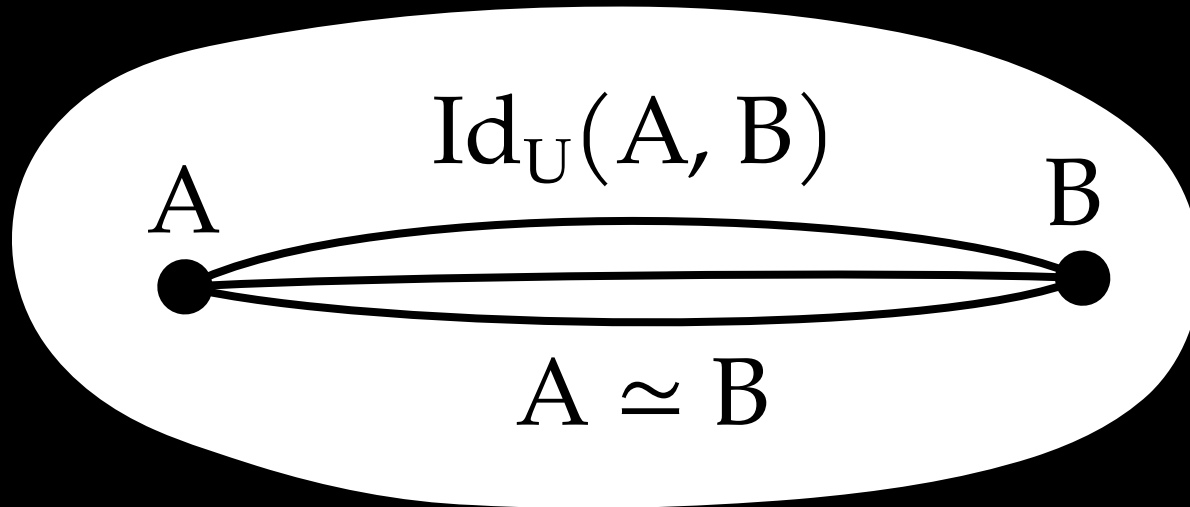
“J”: induction principle for Id

univalence as axiom



$$(A \simeq B) \simeq \text{Id}_U(A, B)$$

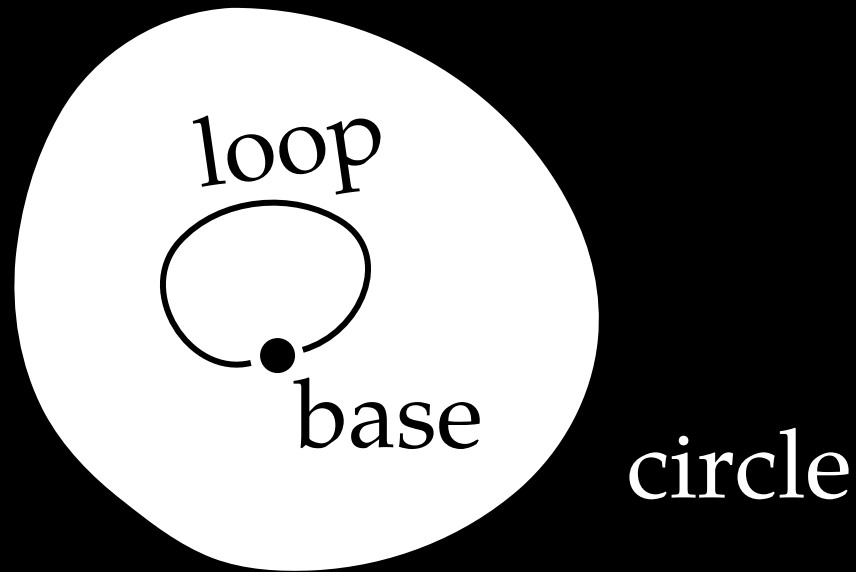
univalence as axiom



$$(A \simeq B) \simeq \text{Id}_U(A, B)$$

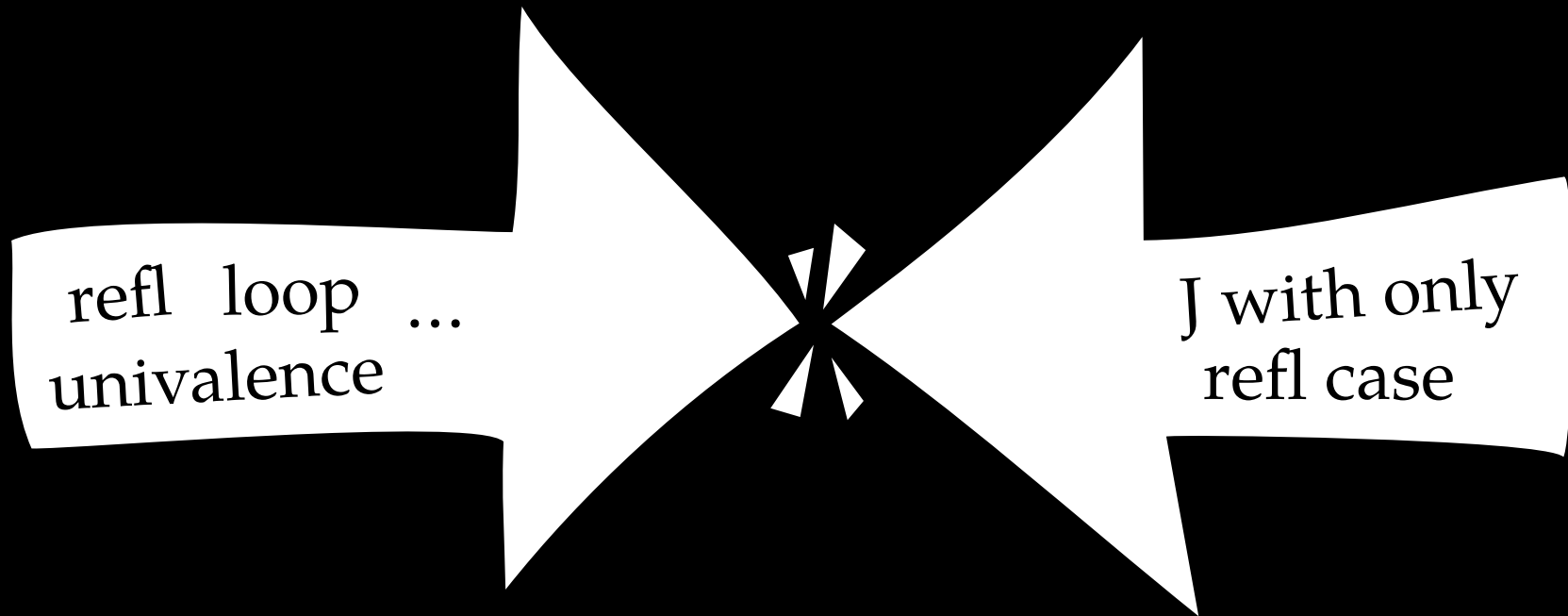


$$(A \simeq B) \rightarrow \text{Id}_U(A, B)$$



base : circle

loop : $\text{Id}_{\text{circle}}(\text{base}, \text{base})$



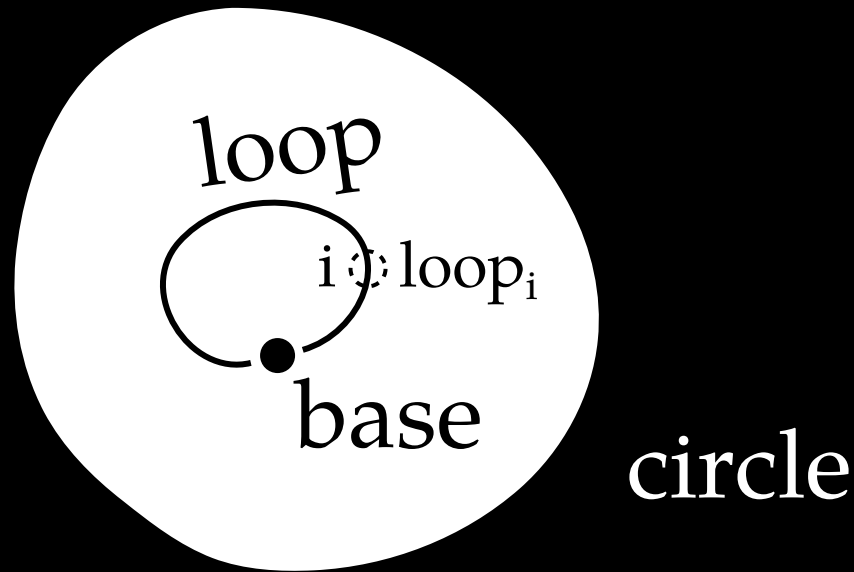
normalization in danger*

$J(\text{loop}) \equiv ???$

*most experts believe the normalization fails

Leave Id alone!

Id/Path should reflect existing paths,
not inducing new ones



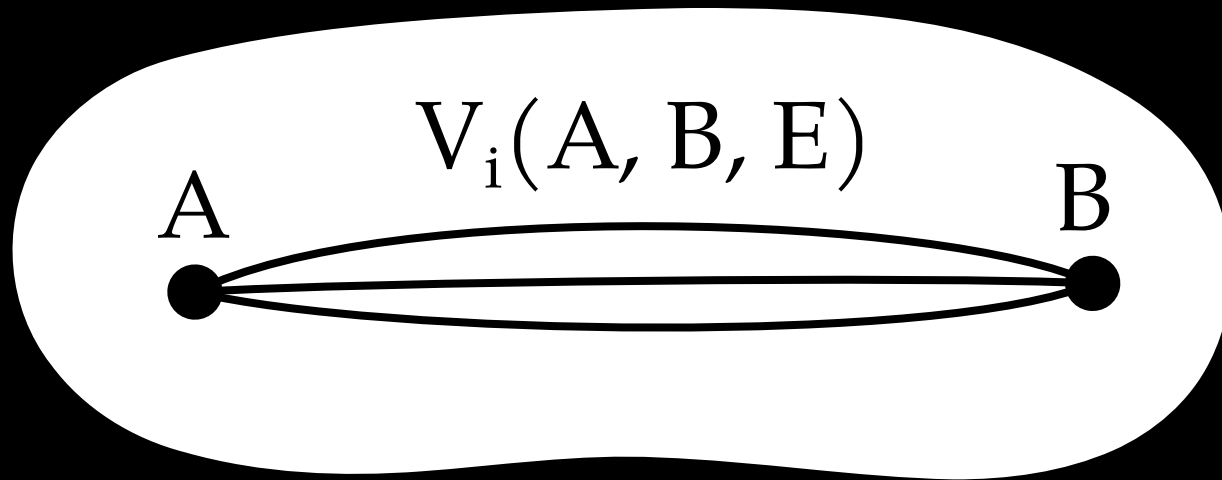
$\text{base} : \text{circle}$

$i : \mathbb{I} \vdash \text{loop}_i : \text{circle}$

$\text{loop}_0 \equiv \text{base} : \text{circle}$

$\text{loop}_1 \equiv \text{base} : \text{circle}$

univalence as a type



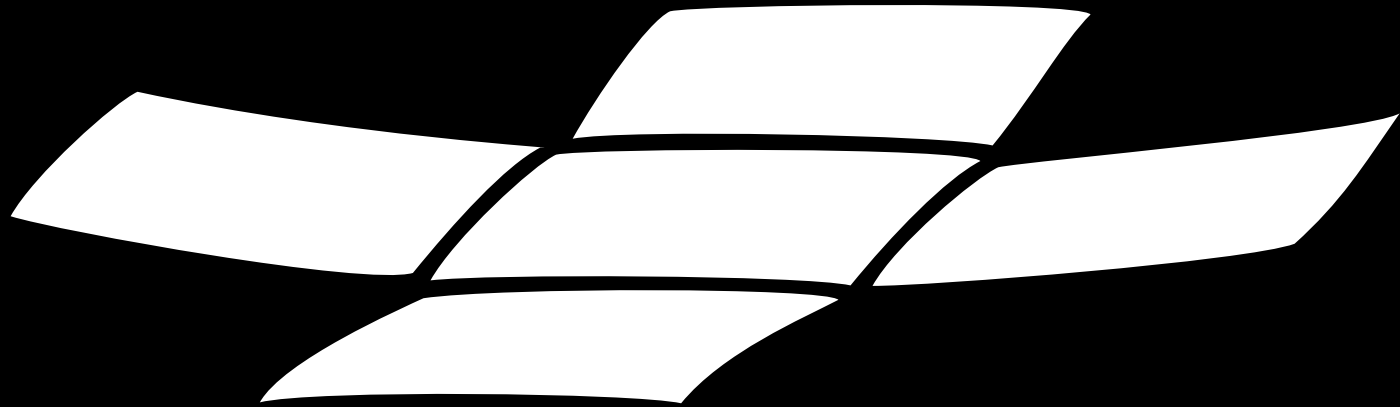
$$E : A \simeq B \text{ (when } i = 0 \text{)}^*$$

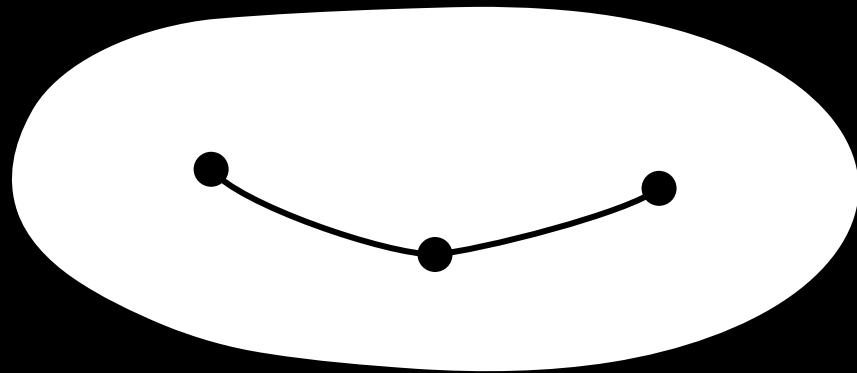
*see [AFH] (not recommended as the first paper to read)

Judgmental framework of paths

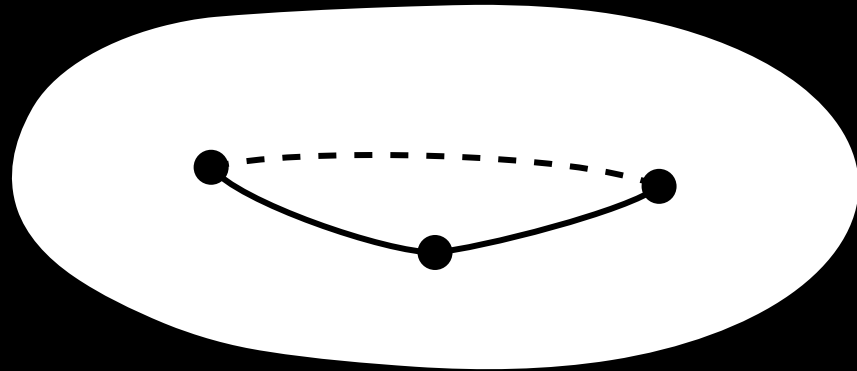
(then internalized by Path/Id)

III. Compositions

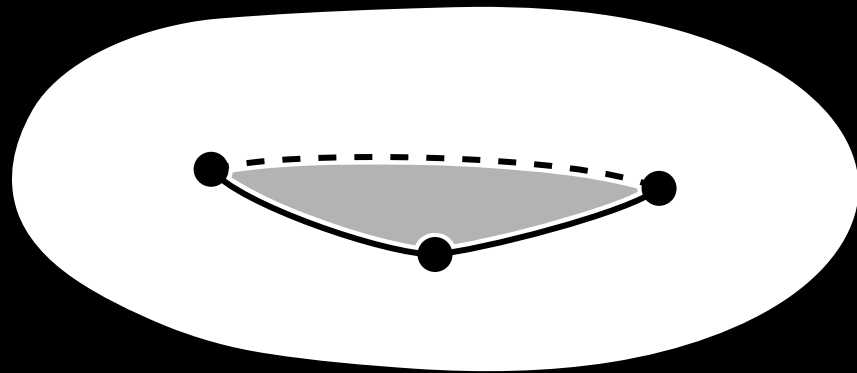




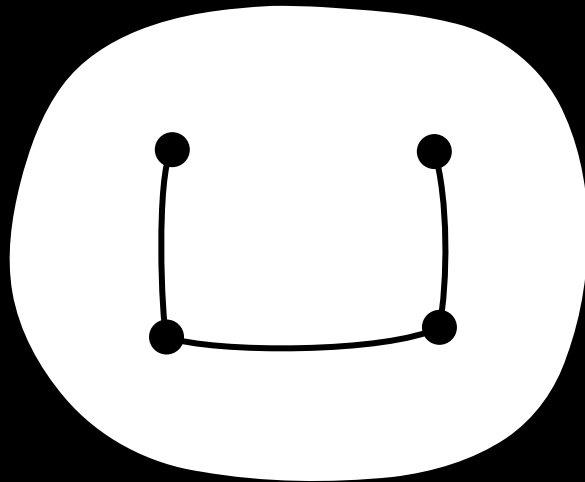
concatenation?



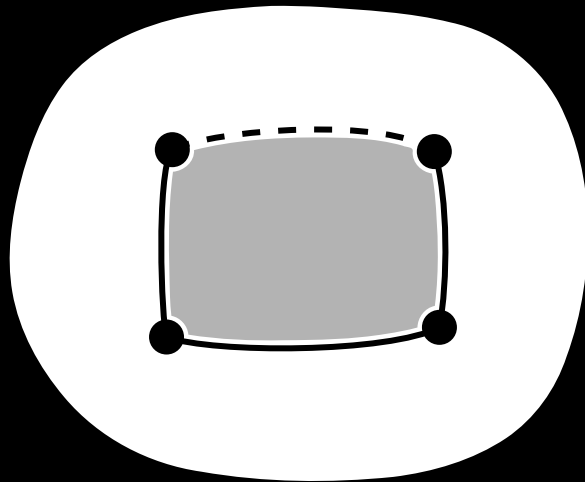
concatenation?



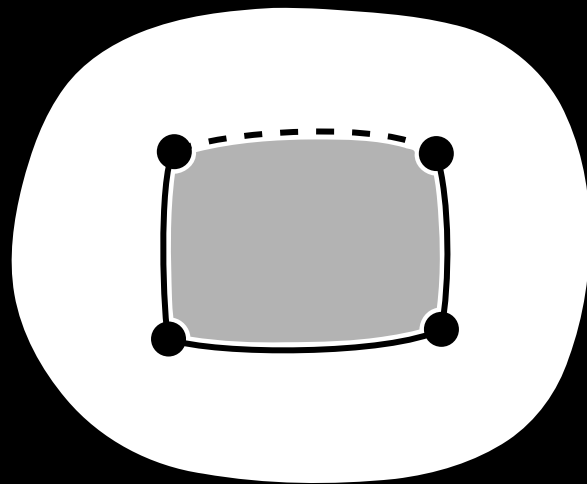
concatenation?



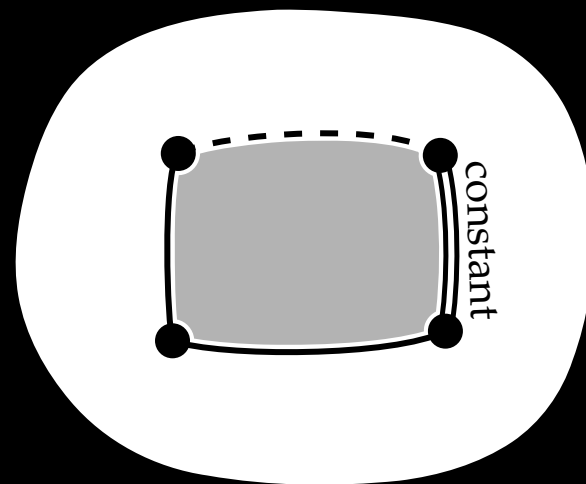
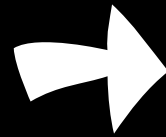
Kan filling
for cubes



Kan filling
for cubes

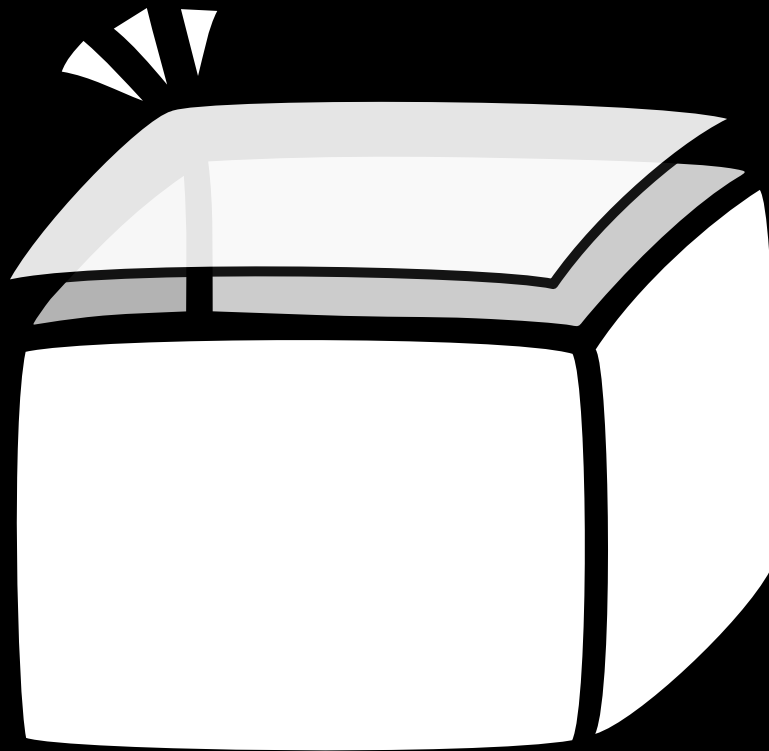


Kan filling
for cubes



concatenation

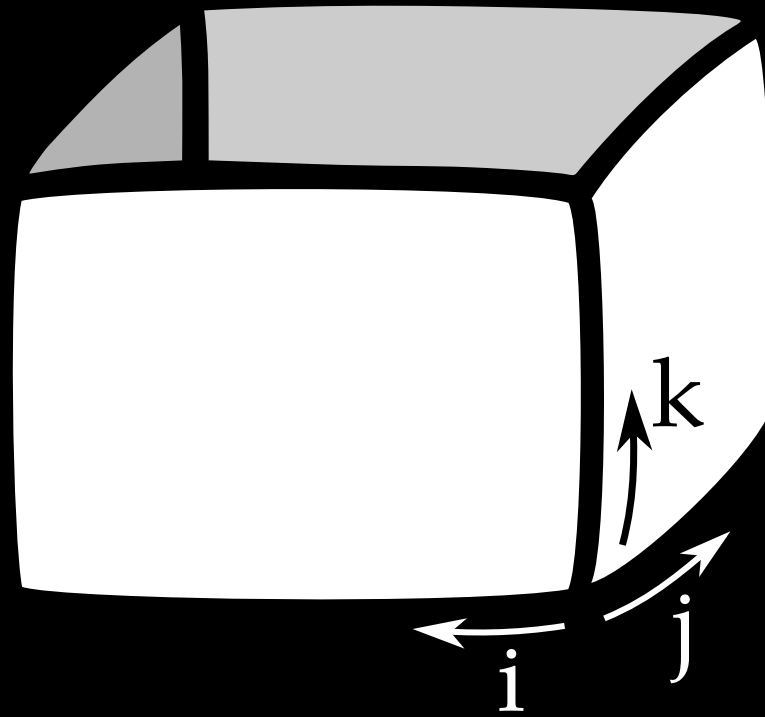
fillers can be done by
higher-dimensional composition*



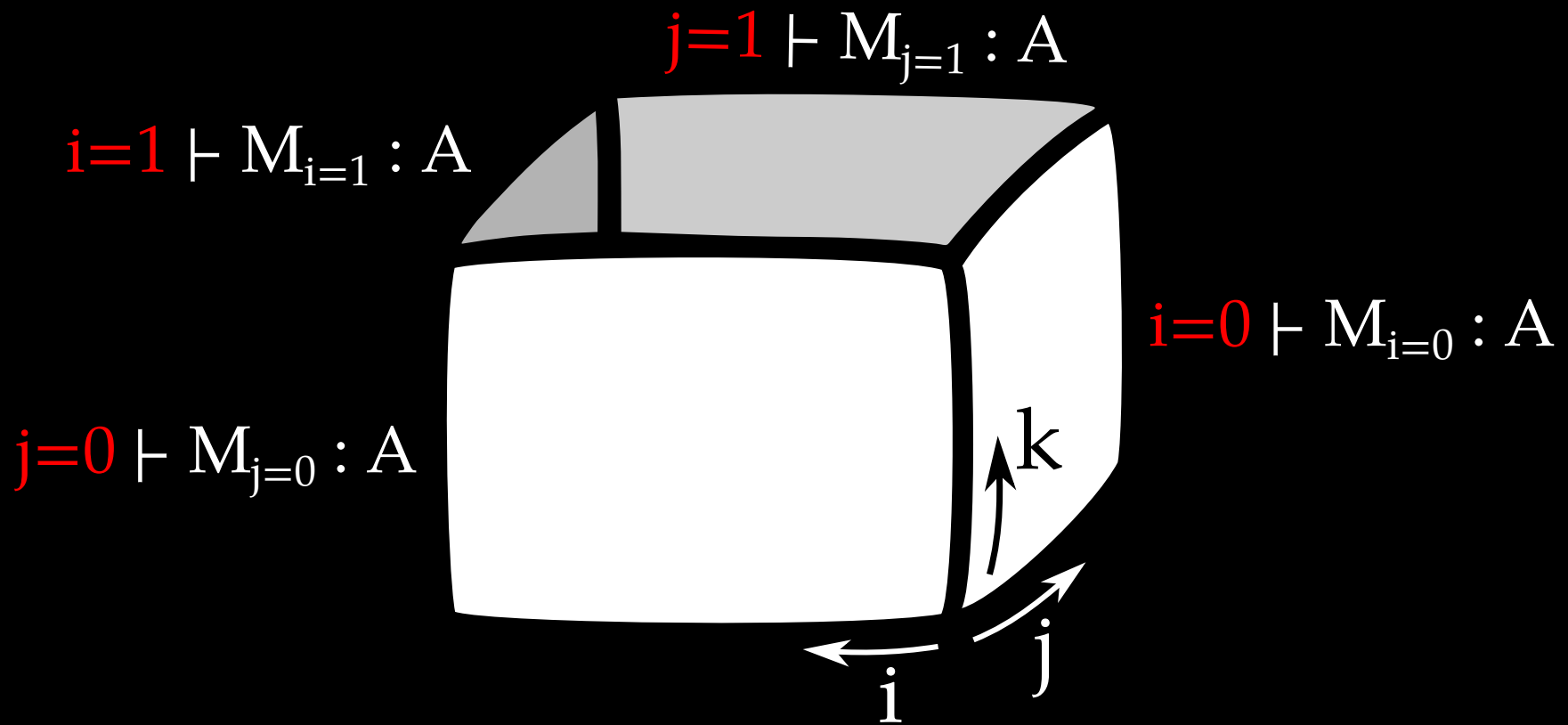
Kan composition

*technical limitations apply; see papers for real (!) math

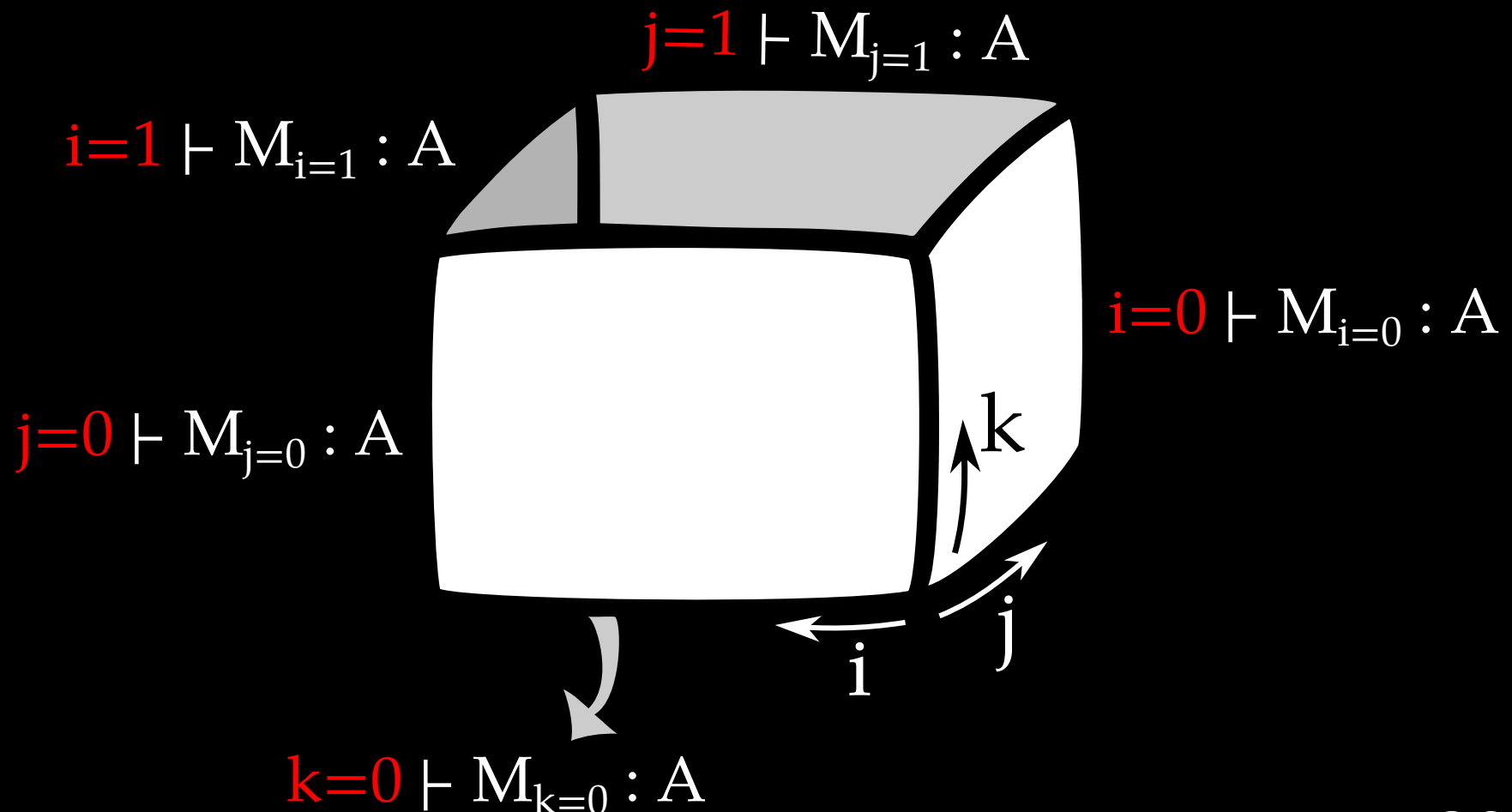
fillers can be done by
higher-dimensional composition*



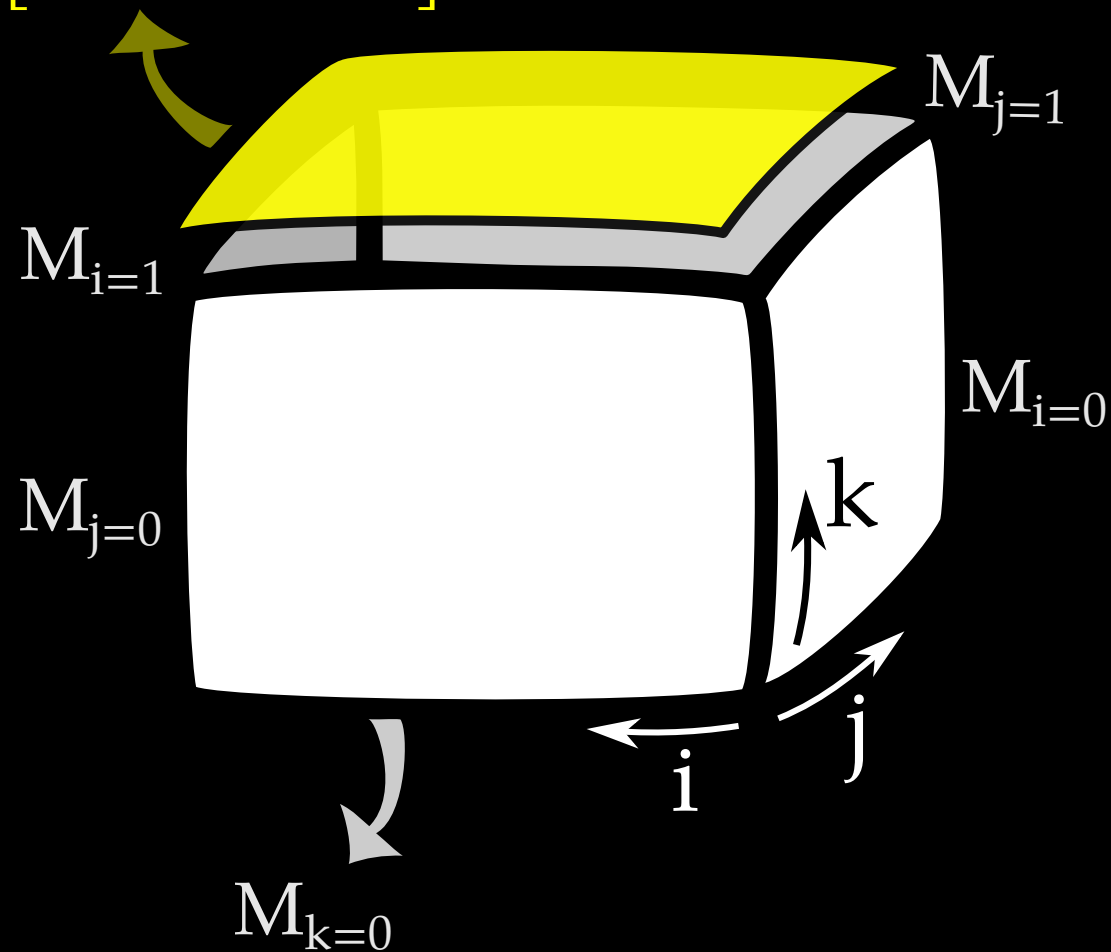
fillers can be done by
higher-dimensional composition*

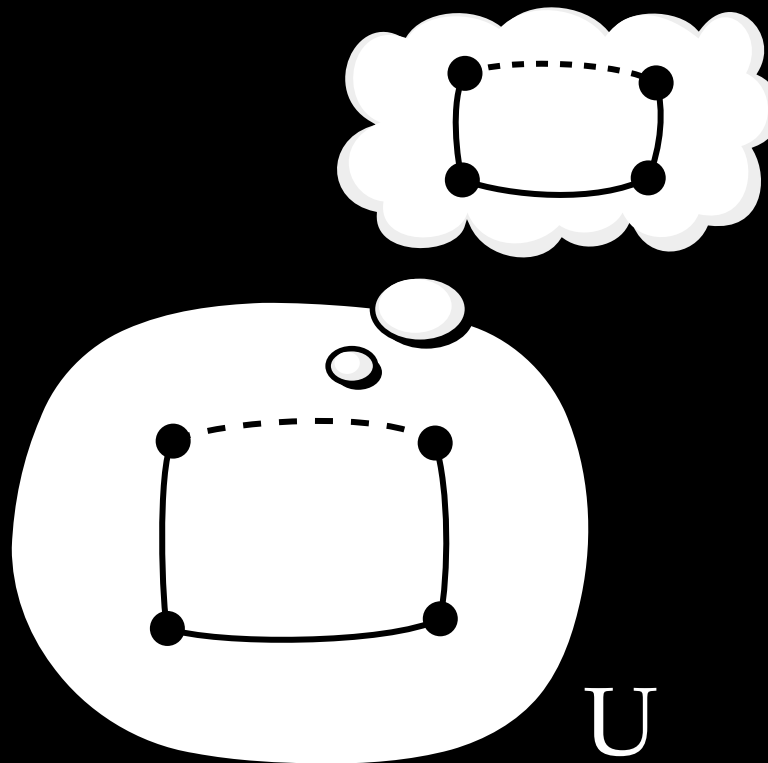


fillers can be done by
higher-dimensional composition*



$$\text{comp}_{k.A} M_{k=0} \begin{bmatrix} i=0 \hookrightarrow M_{i=0} \\ i=1 \hookrightarrow M_{i=1} \\ j=0 \hookrightarrow M_{j=0} \\ j=1 \hookrightarrow M_{j=1} \end{bmatrix} : A[1/k]$$





a composite in a universe is a type itself
which has its own composition operator

$i:\mathbb{I} \vdash M : A$

syntax

models in other
higher toposes?

a model equivalent to
the “Standard”
homotopy theory?

Agda --cubical

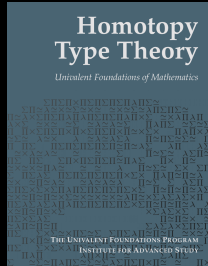
<https://github.com/agda/cubical>

redtt

<https://github.com/RedPRL/redtt>

One Path to Enlightenment

(in this order)



Homotopy Type Theory: Univalent Foundations of Mathematics

Syntax and Models of Cartesian Cubical Type Theory [ABCFLH]

<https://github.com/dlicata335/cart-cube/blob/master/cart-cube.pdf>

Axioms for Modelling Cubical Type Theory in a Topos [OP]

(expanded version)

<https://arxiv.org/abs/1712.04864>

Computational Semantics of Cartesian Cubical Type Theory [A]

(chapter 3, still changing everyday)

<https://www.cs.cmu.edu/~cangiuli/thesis/>