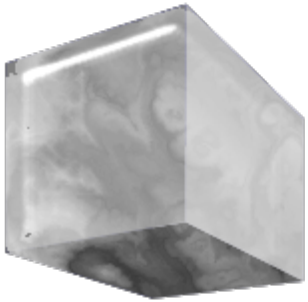
2018.07.07 LFMTP

Cubical Computational Type Theory & RedPRL

>> redprl.org >>



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Cubical

features of homotopy type theory

univalence, higher inductive types

+

Computational

features of Nuprl and PVS

strict equality, strict quotients,
predicative subtypes...

Cartesian Cubical

features of homotopy type theory

univalence, higher inductive types

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strict equality, strict quotients,
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Computational Types

programs/
realizers

computation

Computational Types

programs/ realizers <------ **computational type theory**

computation — theory of computation

Computational Types

programs/ realizers <----- **computational type theory**

computation — theory of computation

meaning explanation <----- Martin-Löf type theory

pre-mathematical in M-L's work

A Minimum Example

```
M := a | bool | true | false | if(M,M,M)
```

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bool val       if(M,Mt,Mf) ↦ if(M',Mt,Mf)
true val       if(true,M,_) ↦ M
false val      if(false,_,M) ↦ M
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The Language

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The Language

What are the types in canonical forms? **{bool}**

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What are the canonical forms of the types?

bool: {true, false}

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What are the types in canonical forms? **{bool}**

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bool: {true, false}

How they are equal? **syntactic equality**

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The Language

What are the types in canonical forms? **{bool}**

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One Theory

A Minimum Example

M := a | bool | true | false | if(M,M,M)

types: {bool} with syntactic equality $\approx$

bool: {true, false} with syntactic equality $\approx_{\text{bool}}$

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$$A \doteq B \text{ type}$$

$$A \Downarrow A' \quad B \Downarrow B' \text{ and } A' \approx B'$$

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$$\text{bool} \doteq \text{bool type}$$

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types: {bool} with syntactic equality ≈

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A ≑ B type

A⇓A' B⇓B' and A'≈B'

---

bool ≑ bool type

if(true,bool,bool) ≑ bool type
⇓bool

A Minimum Example

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types: {bool} with syntactic equality $\approx$

bool: {true, false} with syntactic equality $\approx_{\text{bool}}$

A ≐ B type

A⇓A' B⇓B' and A'≈B'

bool ≐ bool type

if(true,bool,bool) ≐ bool type
⇓bool

if(true,bool,*any closed term*) ≐ bool type

A Minimum Example

```
M := a | bool | true | false | if(M,M,M)
```

types: {bool} with syntactic equality $\approx$

bool: {true, false} with syntactic equality $\approx_{bool}$

$$M \doteq N \in A$$

$A \doteq A$ type, $M \Downarrow M'$, $N \Downarrow N'$, $A \Downarrow A'$ and $M' \approx_{A'} N'$

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$$\mathrm{false} \doteq \mathrm{false} \in \mathrm{bool}$$

$$\underset{\Downarrow \mathrm{true}}{\mathrm{if(true,true,bool)}} \doteq \mathrm{true} \in \underset{\Downarrow \mathrm{bool}}{\mathrm{if(true,bool,bool)}}$$

A Minimum Example

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types: {bool} with syntactic equality $\approx$

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$$a{:}A \gg M \doteq N \in B$$

$P \doteq Q \in A$ implies $M[P/a] \doteq N[Q/a] \in B[P/a]$

A Minimum Example

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M := a | bool | true | false | if(M,M,M)
```

types: {bool} with syntactic equality $\approx$

bool: {true, false} with syntactic equality $\approx_{\mathrm{bool}}$

a:A >> M ≐ N ∈ B

P ≐ Q ∈ A implies M[P/a] ≐ N[Q/a] ∈ B[P/a]

---

b:bool >> b ≐ if(b,true,false) ∈ bool?

A Functional Example

```
M := a | M1⟶M2 | \a.M | M1 M2 | ...

(M1⟶M2) val  \a.M val  (\a.M1)M2 ↦ M1[M2/a]
```

Another Language

A Functional Example

```
M := a | M1⟶M2 | \a.M | M1 M2 | ...

(M1⟶M2) val   \a.M val   (\a.M1)M2 ↦ M1[M2/a]
```

Another Language

What are the types in canonical forms?

```
  the least fixed point of
  S ↦ {M⟶N | M⇓, N⇓ in S} union ...
```

What are the canonical forms of the types?

```
  A⟶B: {\a.M}
```

How they are equal?

```
  A1⟶B1 ≈ A2⟶B2 if A1 ≐ A2 and B1 ≐ B2
```

$\backslash a.M1 \approx_{A\longrightarrow B} \backslash a.M2$ if $a{:}A \gg M1 \doteq M2 \in B$

Variables

Nuprl/...	Coq/Agda/...
Vars range over closed terms	Vars are indet.
Defined by transition b/w closed terms	Defined by conversion b/w open terms

Open-endedness

Proof theory/tactics/editors

↓

Computational type theory

↓

Programming language

Open-endedness

Proof theory/tactics/editors

↓

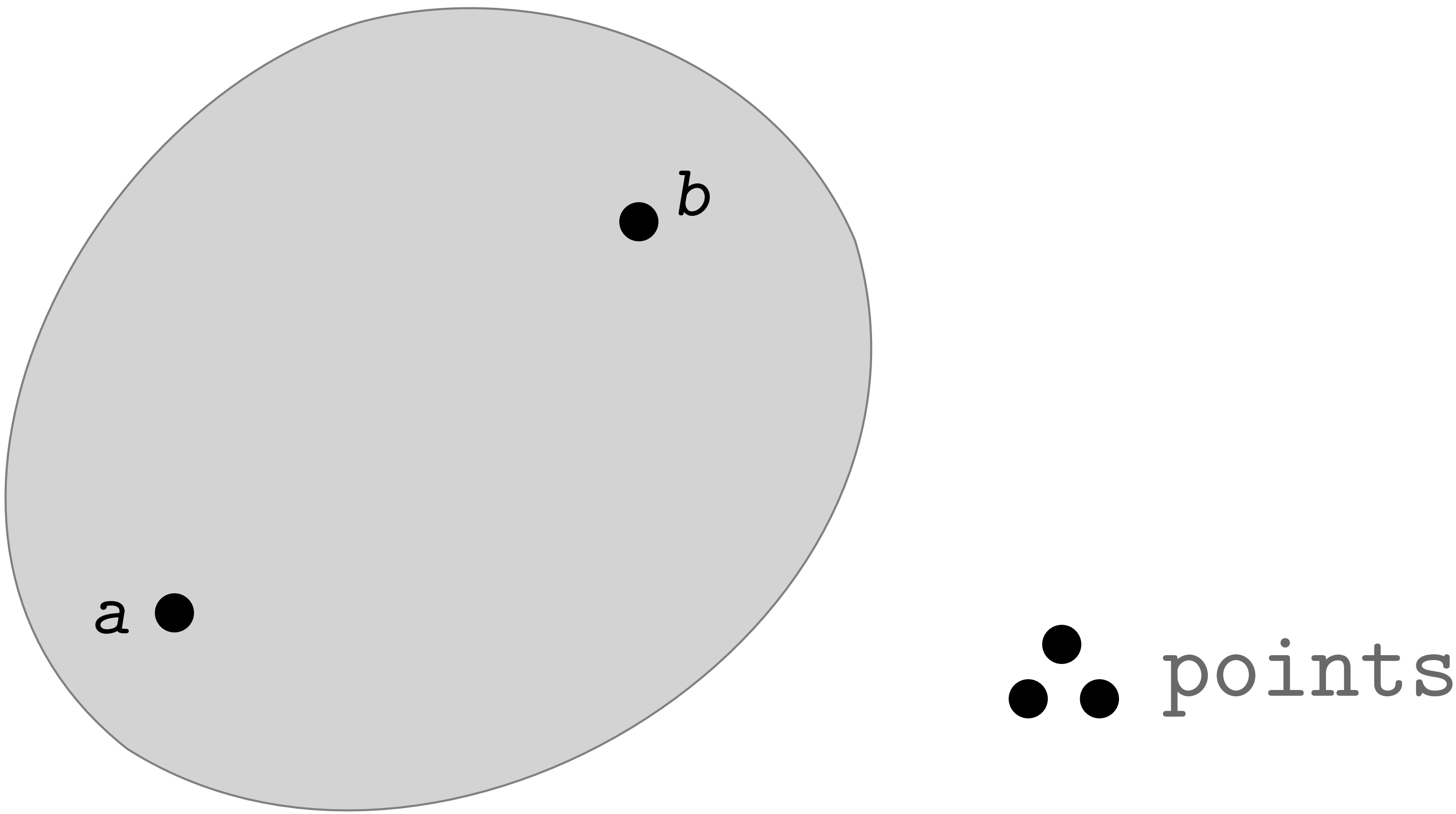
Computational type theory

↓

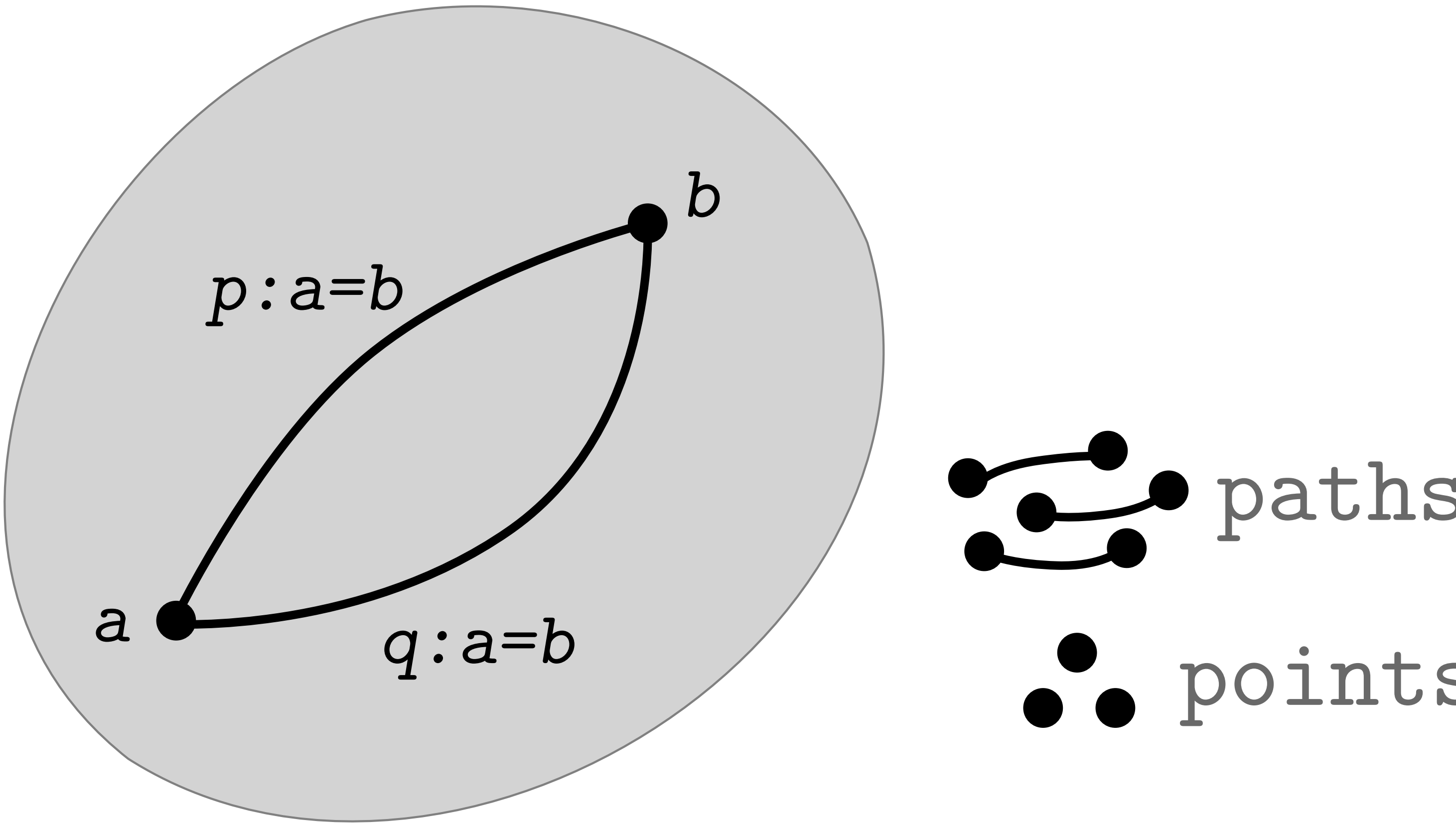
Programming language

Canonicity always holds

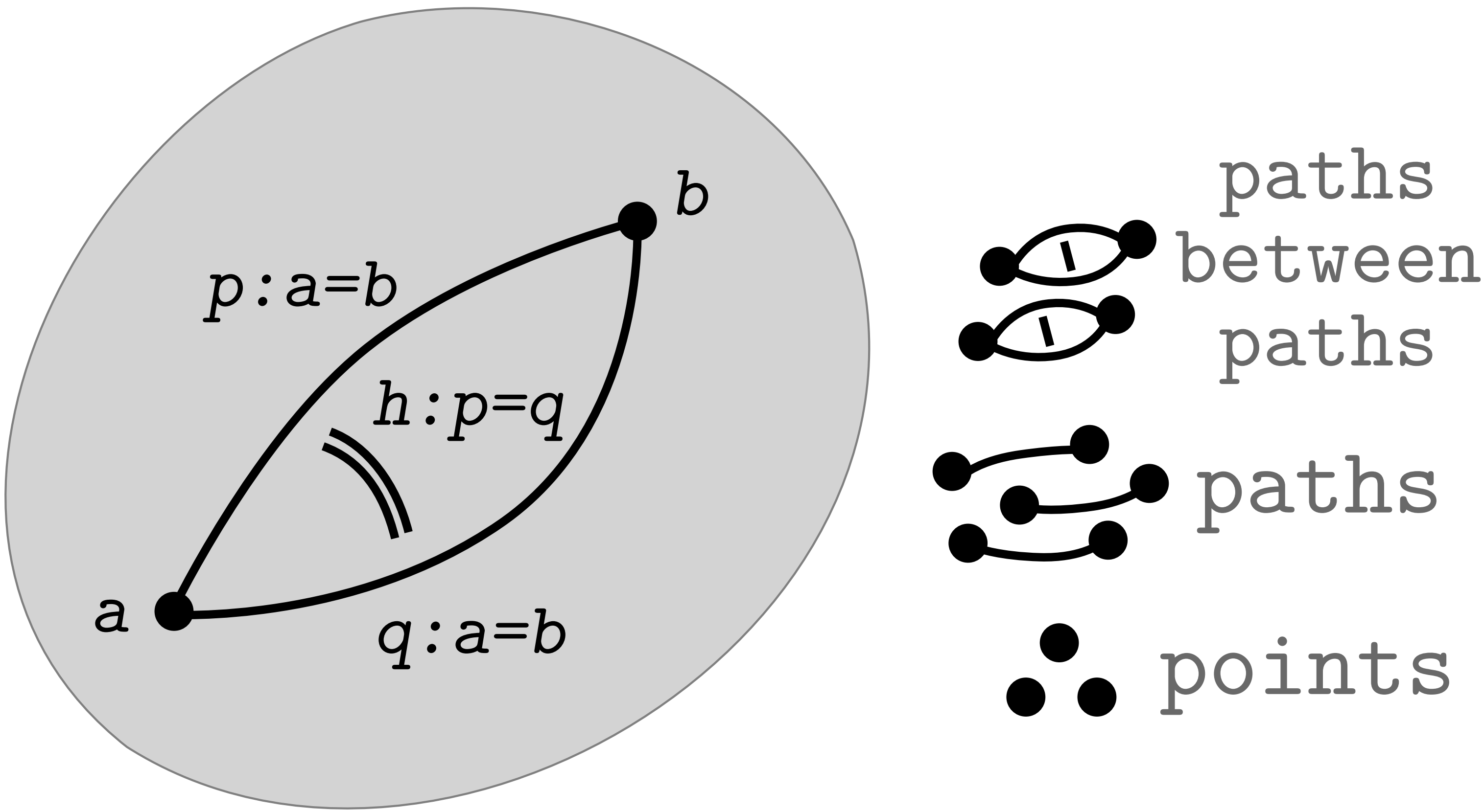
Homotopy Type Theory



Homotopy Type Theory



Homotopy Type Theory



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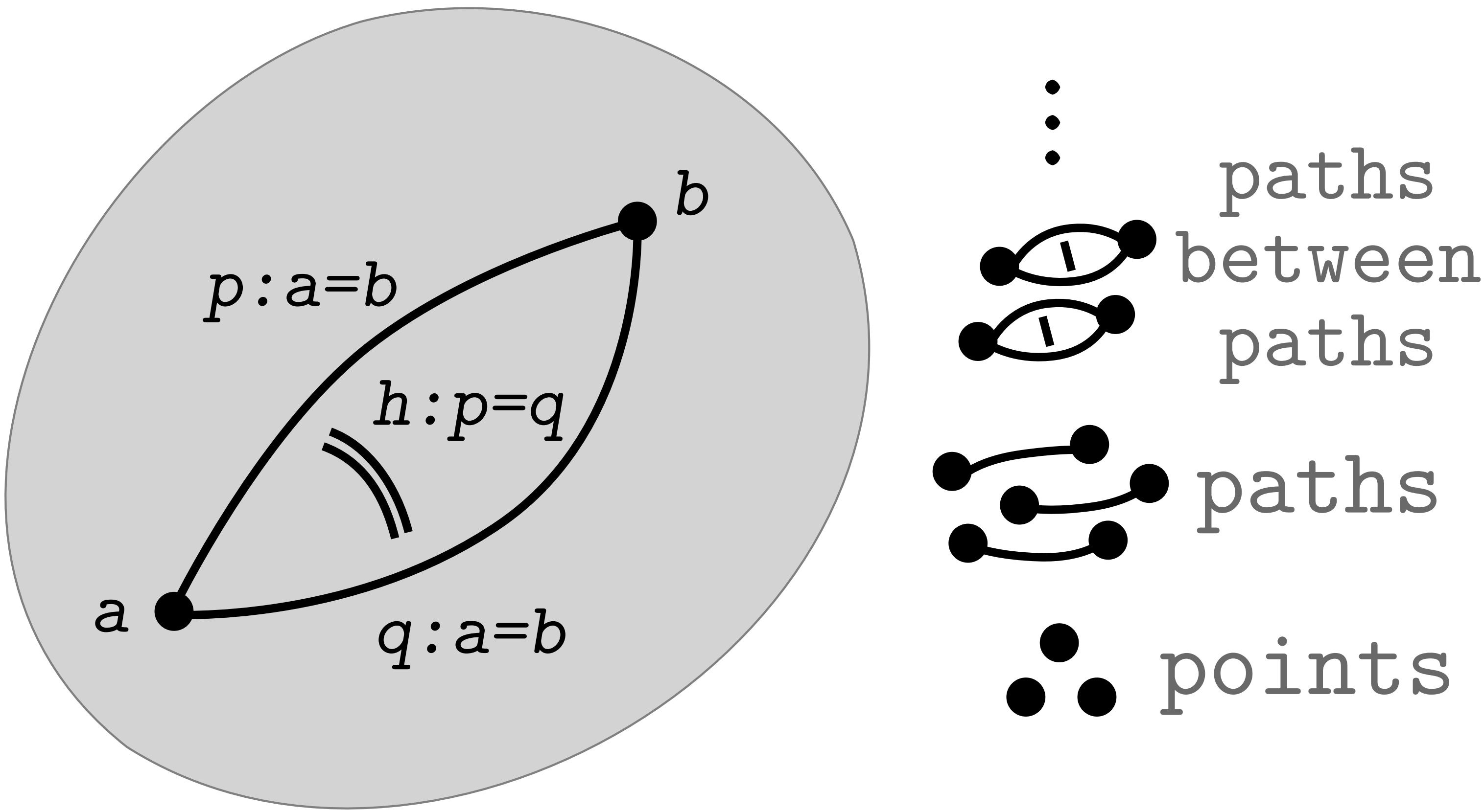

Equality and Paths

Equality ($\equiv$)

Silent in theory

$$2 + 3 \equiv 5$$

$$fst\ \langle M, N \rangle \equiv M$$

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If $A \equiv B$ and $M : A$ then $M : B$

Equality and Paths

Equality (≡)

Silent in theory

$2 + 3 \equiv 5$

$fst\ \langle M,N \rangle \equiv M$

If $A \equiv B$ and $M : A$ then $M : B$

Paths (=)

Visible in theory

If $P : A = B$ and $M : A$ then $transport(M,P) : B$

Homotopy Type Theory

[Awodey and Warren] [Voevodsky *et al*] [van den Berg and Garner]

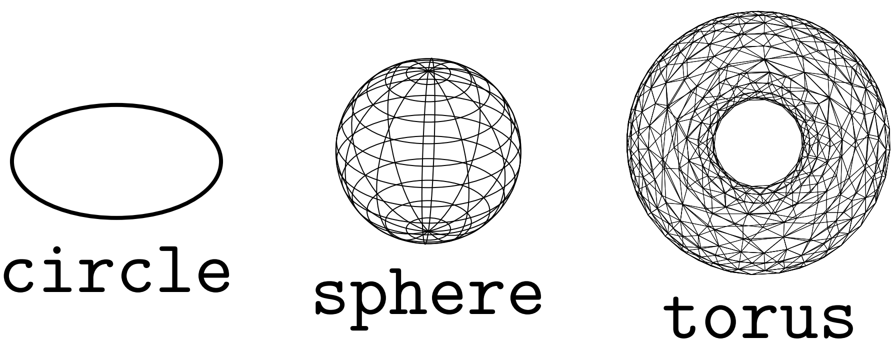
A	Type	Space
$a : A$	Element	Point
$f : A \to B$	Function	Continuous Mapping
$C : A \to Type$	Dependent Type	Fibration
$a =_A b$	Identification	Path

Features of HoTT

Univalence

If E is an equivalence between types A and B, then $ua(E) : A = B$

Higher Inductive Types



Canonicity?

Canonicity broken by
new features stated as axioms!

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Canonicity

For any $M : bool$, either
$M \equiv true : bool$ or $M \equiv false : bool$

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Canonicity broken by
new features stated as axioms!

Canonicity

For any $M : bool$, either
$M \equiv true : bool$ or $M \equiv false : bool$

$ua(not) : bool = bool$

$transport(ua(not), true) \not\equiv false$

Canonicity for All

Canonicity for bool means
canonicity for *everyone*

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$$M : bool \times A$$

$$fst(M) \equiv ??? : bool$$

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Wants $M \equiv \langle P,Q \rangle$ and then
$fst(M) \equiv fst\langle P,Q \rangle \equiv P \equiv$ true or false

Canonicity for Paths?

$$\frac{M : A}{\mathit{refl}(M) : M =_A M}$$

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$$\frac{M : A}{\mathit{refl}(M) : M =_A M}$$

$$\frac{a{:}A \vdash R : C(a,a,\mathit{refl}(a)) \qquad P : M = N}{\mathit{path\text{-}ind}[C](a.R,P) : C(M,N,P)}$$

Canonicity for Paths?

$$\frac{M : A}{refl(M) : M =_A M}$$

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$$\frac{a{:}A \vdash R : C(a,a,refl(a)) \qquad M : A}{\begin{array}{c} path\text{-}ind[C](a.R,refl(M)) \equiv R[M/a] \\ : C(M,M,refl(M)) \end{array}}$$

Canonicity for Paths?

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$$\mathtt{path\text{-}ind}[C](a.R,\mathtt{ua}(E)) \equiv \;???$$

Restore Canonicity

Can we have a new TT with canonicity + univalence?

Yes with De Morgan cubes [CCHM 2016]
Yes with Cartesian cubes [AFH 2017]

... and higher inductive types?

Examples with De Morgan cubes [CHM 2018]
Yes with Cartesian cubes [CH 2018]

Restore Canonicity

Idea: each type manages its own paths

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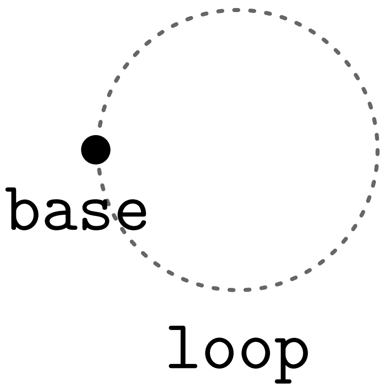


```
base : S1
```

Restore Canonicity

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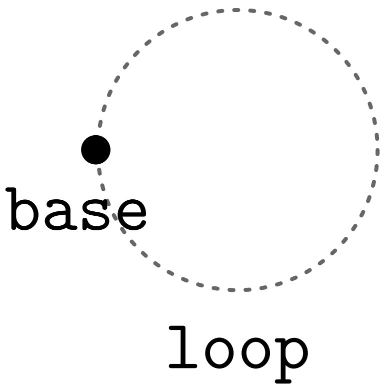


```
base : S1
loop : base = base
```

Restore Canonicity

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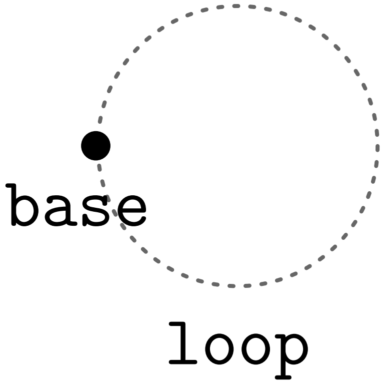

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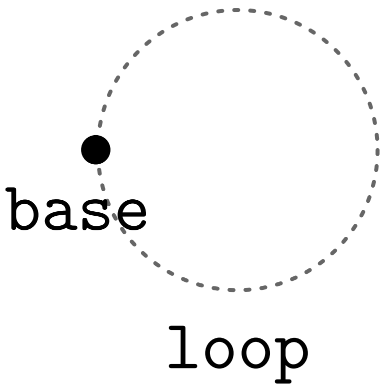

base : S1

~~loop : base = base~~

x:𝕀 ⊢ loop{x} : S1

loop{0} ≡ base : S1

loop{1} ≡ base : S1

Restore Canonicity

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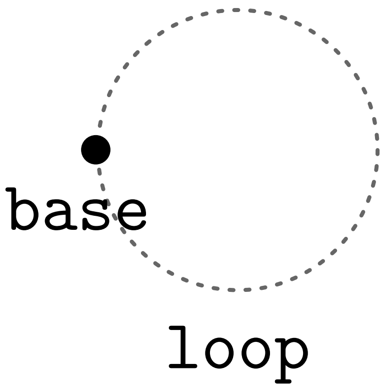

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Kan structure:
sufficient to implement path-ind

Kan types: types with Kan structure

Cartesian Cubes

Introducing $\mathbb{I}$ the formal interval

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Introducing $\mathbb{I}$ the formal interval

$$\Gamma \vdash 0 : \mathbb{I} \qquad \Gamma \vdash 1 : \mathbb{I}$$

$$\Gamma,\ x : \mathbb{I},\ \Gamma' \vdash x : \mathbb{I}$$

Cartesian Cubes

Introducing $\mathbb{I}$ the formal interval

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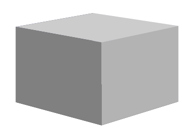
$$x_1 : \mathbb{I},\ x_2 : \mathbb{I},\ \ldots,\ x_n : \mathbb{I} \vdash M : A$$

$\Leftrightarrow$ M is an n-cube in A 

Cartesian Cubes

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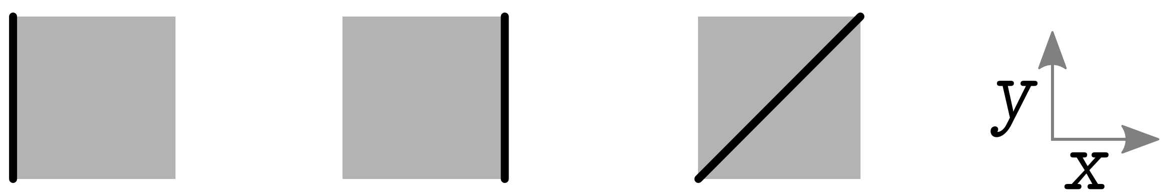
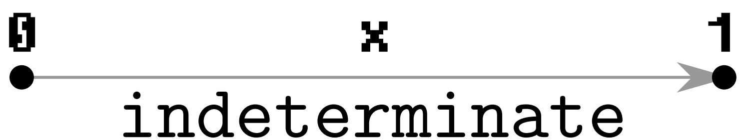
Cartesian: works as normal contexts

$$M\langle 0/x \rangle \quad M\langle 1/x \rangle \quad M\langle y/x \rangle$$

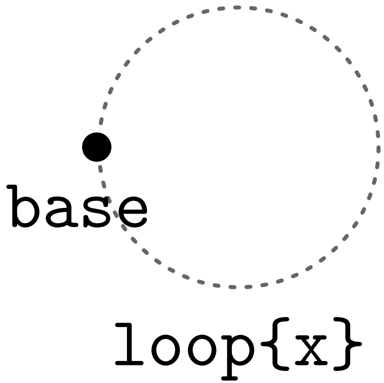


Cubical Programming

```
dim expr r := 0 | 1 | x
```



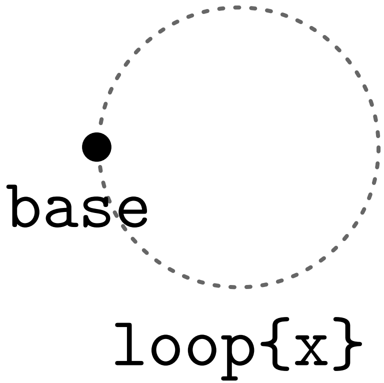
Circle



Circle

```
M := S1 | base | loop{r}
   | S1elim(a.M, M, M, x.M) | ...
```

dim expr



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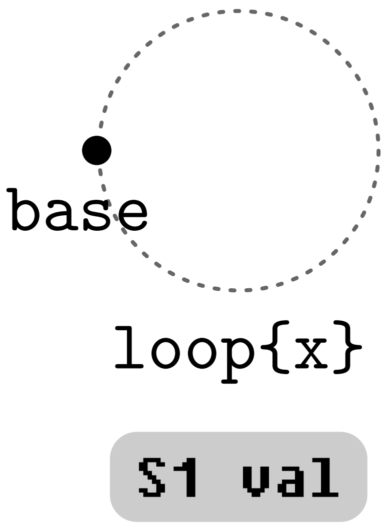

Circle

```
M := S1 | base | loop{r}
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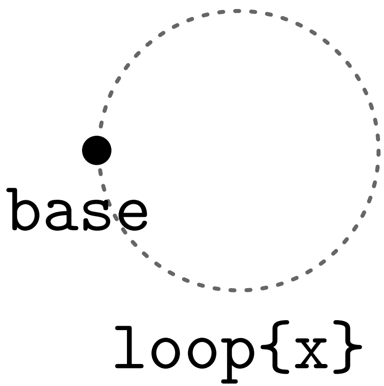
```

(r → dim expr)



`S1 val`

`base val`

Circle

```
M := S1 | base | loop{r}
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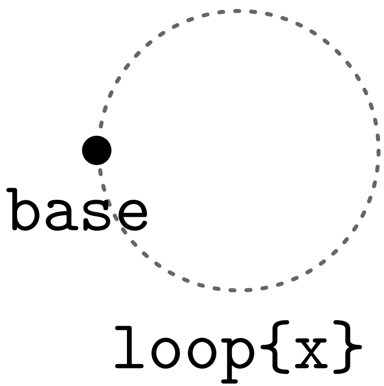
dim expr



`S1 val`

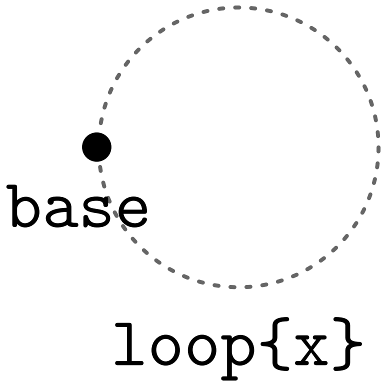
`base val`

`loop{x} val`

`loop{0} ↦ base`

`loop{1} ↦ base`

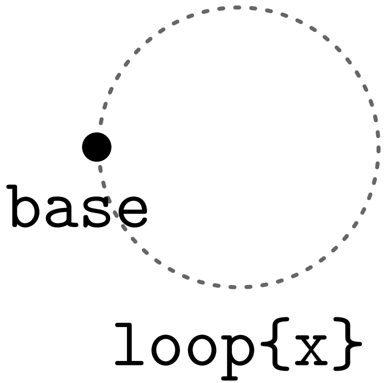
Circle



S1 val

```
M ↦ M'
──────────────────────────────
S1elim(a.A, M, B, x.L)
  ↦ S1elim(a.A, M', B, x.L)
```

Circle

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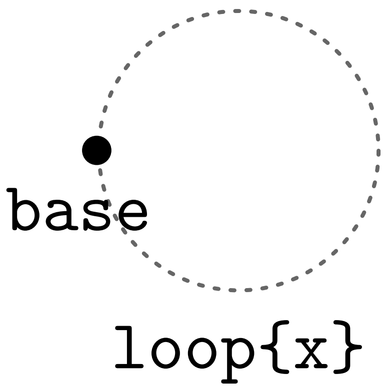
```

```
S1elim(a.A, base, B, x._)
  ↦ B
```

Circle

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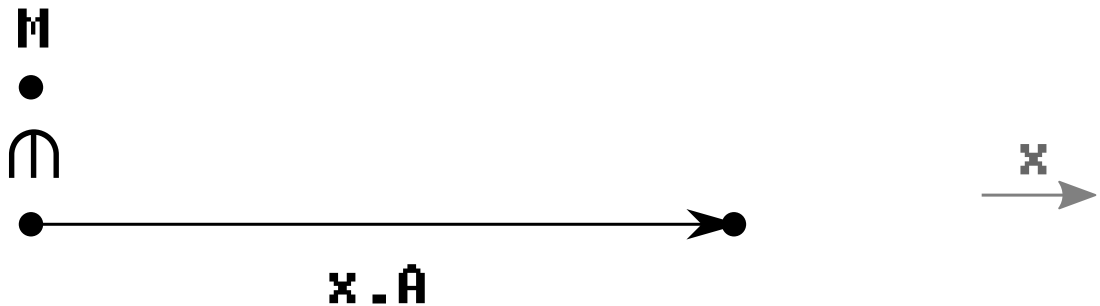
```

```
S1elim(a.A, base, B, x._)
  ↦ B
```

```
S1elim(a.A, loop{x}, _, y.L)
  ↦ L<x/y>
```

Kan 1/2: Coercion



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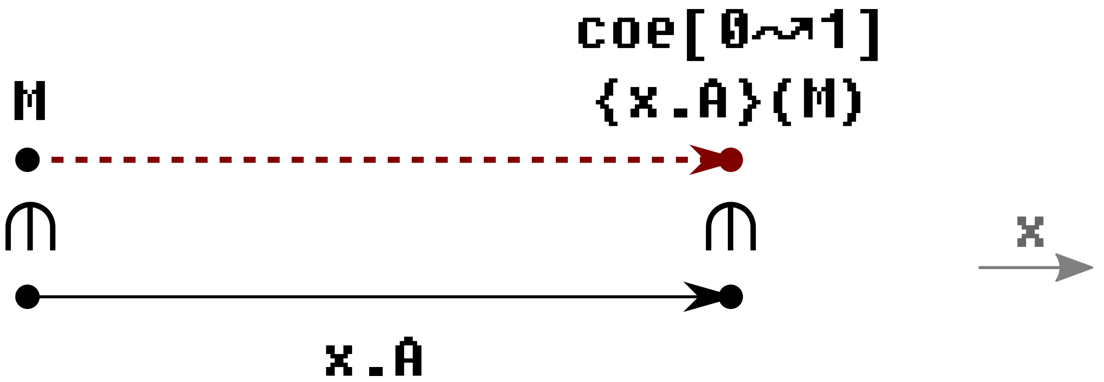

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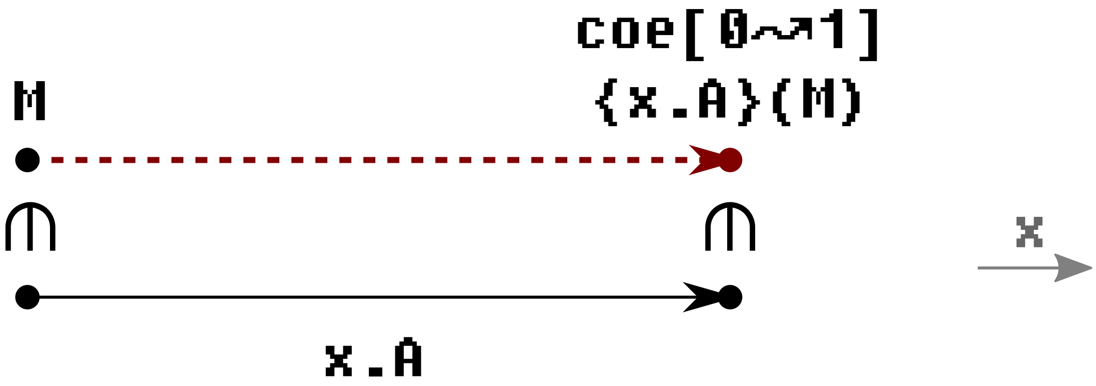
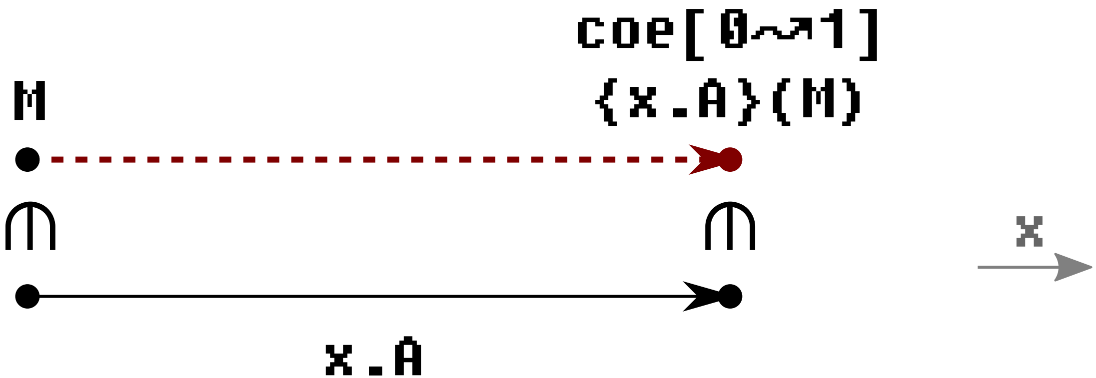


coe[r⇝r']{x.A}(M) ∈ A<r'/x>

⋔

A<r/x>

Kan 1/2: Coercion



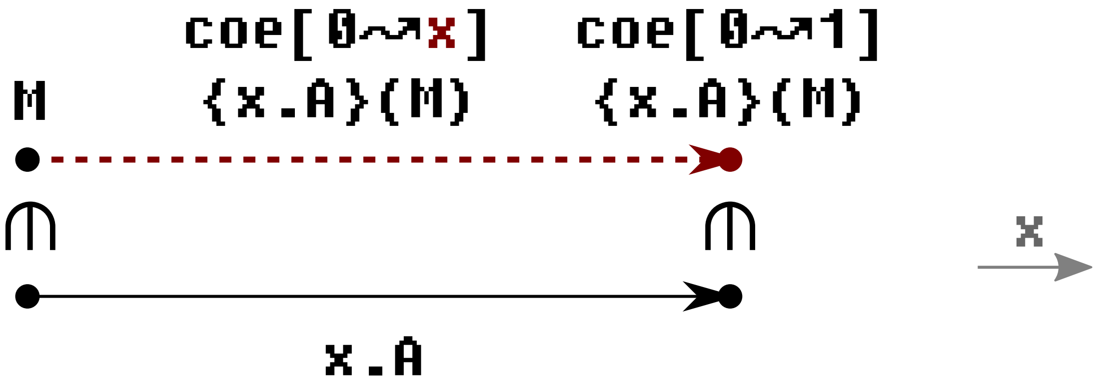
coe[r↝r']{x.A}(M) ∈ A<r'/x>

⋒

A<r/x>

coe[r↝r]{x.A}(M) ≐ M ∈ A<r/x>

Kan 1/2: Coercion



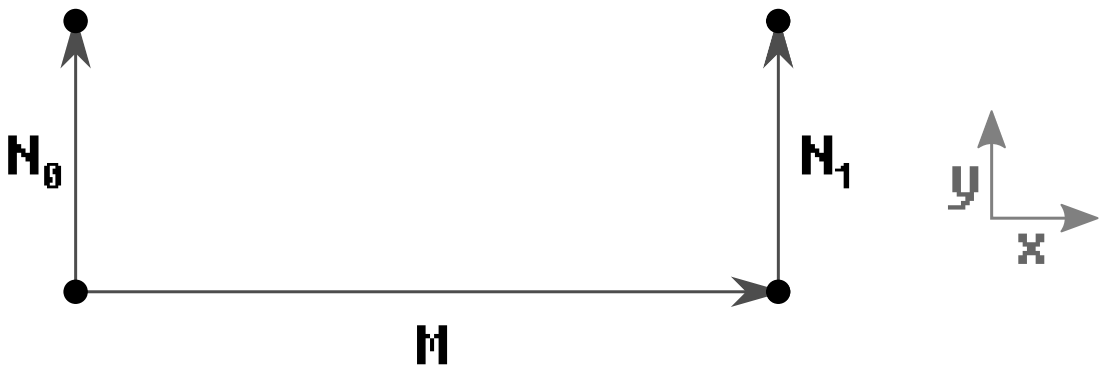
coe[r⤳r']{x.A}(M) ∈ A<r'/x>

⋔

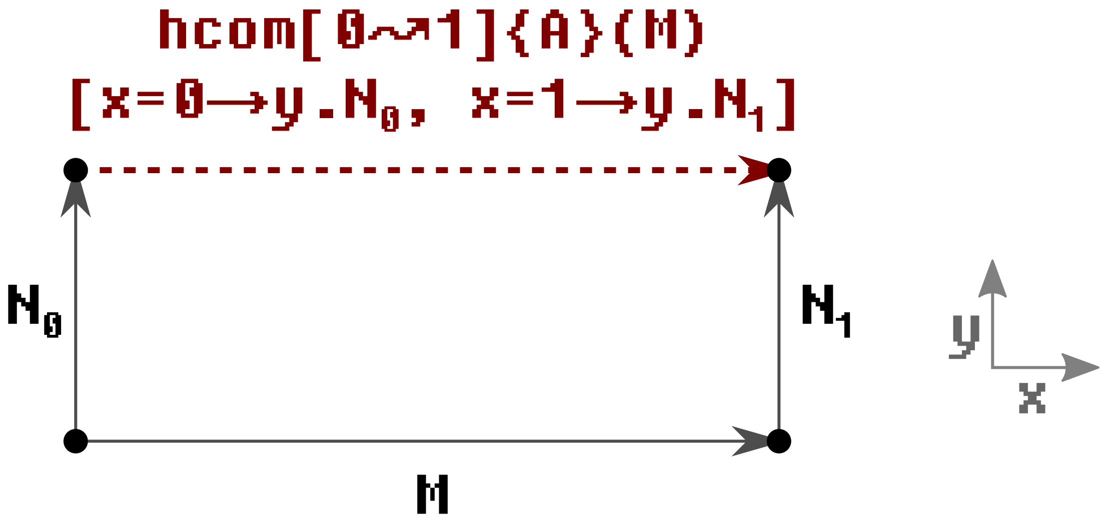
A<r/x>

coe[r⤳r]{x.A}(M) ≐ M ∈ A<r/x>

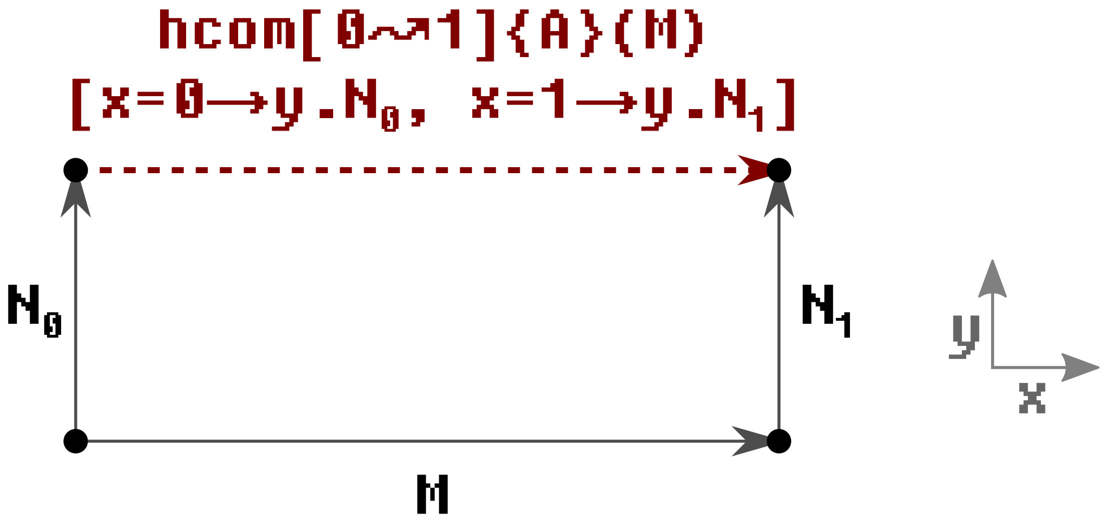
Kan 2/2: Homogeneous Comp.



Kan 2/2: Homogeneous Comp.

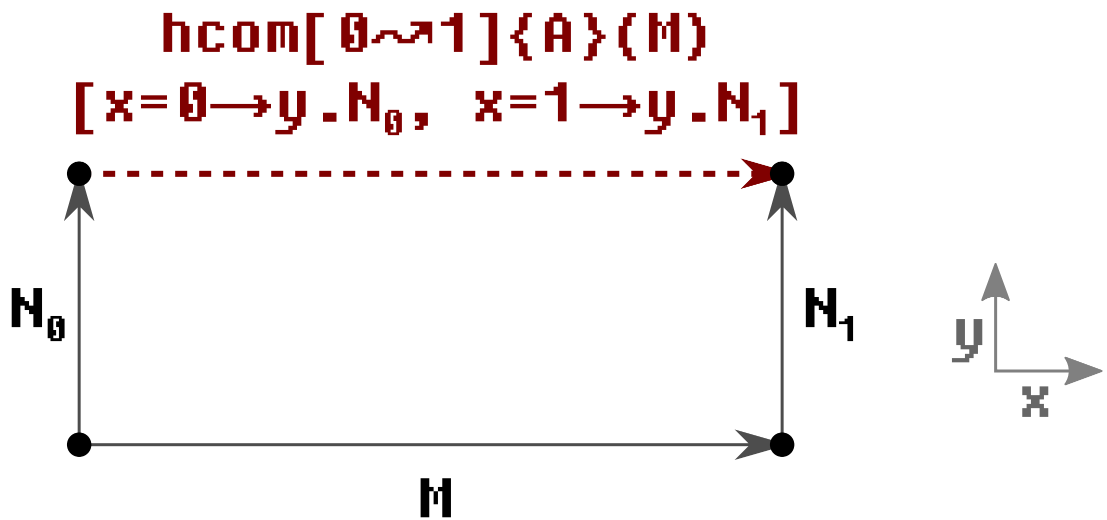


Kan 2/2: Homogeneous Comp.



hcom[r⤳r']{A}(M) [..., r_i=r'_i⟶y.N_i, ...] ∈ A

Kan 2/2: Homogeneous Comp.

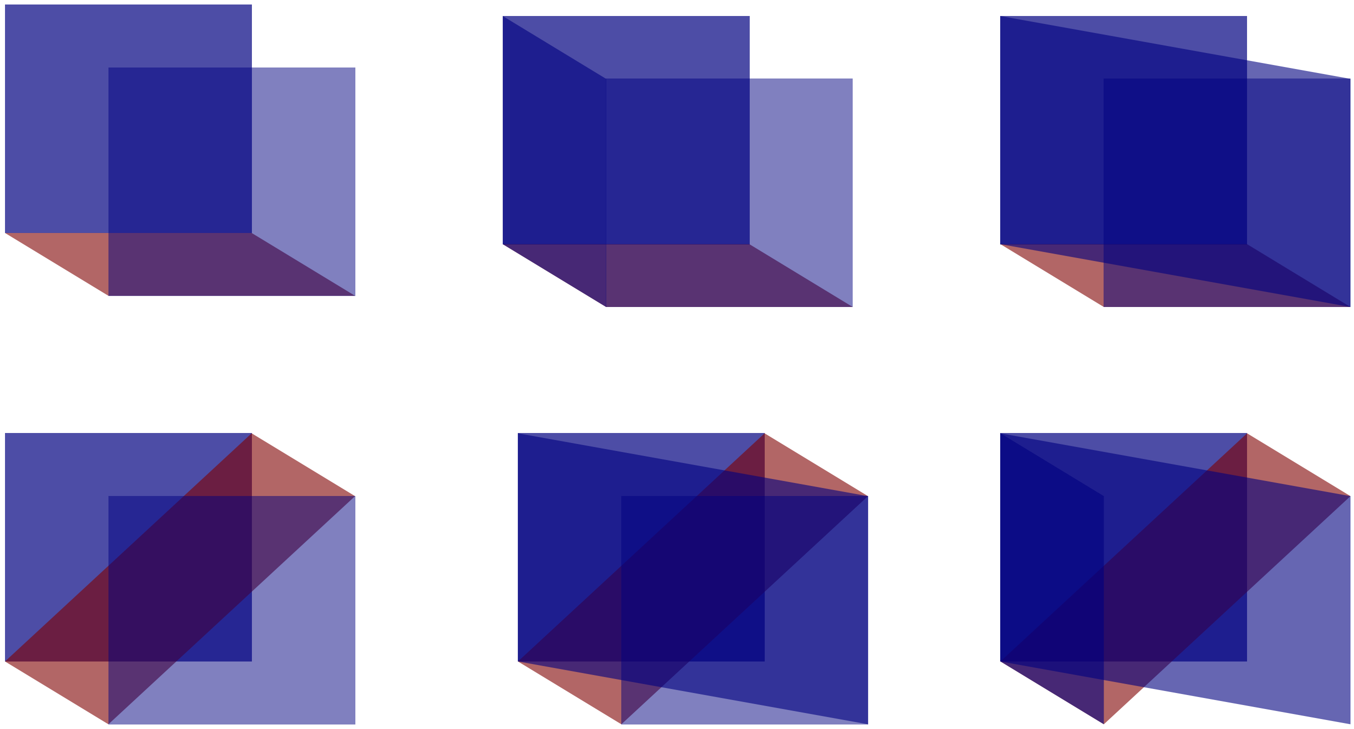


hcom[r⤳r']{A}(M) [..., r_i=r'_i⟶y.N_i, ...] ∈ A

hcom[r⤳r]{A}(M) ≐ M ∈ A

hcom[r⤳r']{A}(M)[..., r_i=r_i⟶y.N_i, ...]
  ≐ N_i⟨r'/y⟩ ∈ A

Kan 2/2: Homogeneous Comp.



Kan Circle

```
coe[r~>r'](_.S1)(M) ↦ M
```

Kan Circle

```
coe[r~>r'](_.S1)(M) ↦ M
```

```
hcom[r~>r'](S1)(M)[…] ↦ fhcom[r~>r'](M)[…]
```

formal homo.
composition

Kan Circle

```
coe[r~>r']{_.S1}(M) ↦ M
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```
hcom[r~>r']{S1}(M)[…] ↦ fhcom[r~>r'](M)[…]
```

formal homo. composition

```
fhcom[r~>r](M)[…] ↦ M
```

Kan Circle

```
coe[r⤳r']{_.S1}(M) ↦ M
```

```
hcom[r⤳r']{S1}(M)[…] ↦ fhcom[r⤳r'](M)[…]
```

formal homo.
composition

```
fhcom[r⤳r](M)[…] ↦ M
```

```
r≠r'   rᵢ=r'ᵢ (the first i)
──────────────────────────────────────────────
fhcom[r⤳r'](M)[…, rᵢ=r'ᵢ→y.Nᵢ, …] ↦ Nᵢ<r'/y>
```

Kan Circle

```
coe[r⇝r'](_.S1}(M) ↦ M
```

```
hcom[r⇝r']{S1}(M)[…] ↦ fhcom[r⇝r'](M)[…]
```

formal homo. composition

```
fhcom[r⇝r](M)[…] ↦ M
```

```
r!=r'   r_i=r'_i (the first i)
--------------------------------------------
fhcom[r⇝r'](M)[…, r_i=r'_i→y.N_i, …] ↦ N_i<r'/y>
```

```
r!=r'   r_i!=r'_i for all i
----------------------------
fhcom[r⇝r'](M)[…] val
```

Kan Circle

`S1elim` needs to handle `fcom`

Kan Circle

S1elim needs to handle **fcom**

```
r≠r'  rᵢ≠r'ᵢ
─────────────────────────────────────────
S1elim(a.A, fhcom[r⇝r'](M)[…], B, x.L)
  ↦ com[r⇝r']{y.A[fhcom[r⇝y](M)[…]/a}
        (S1elim(M, B, x.L))[…]
```

S1elim(composition) ↦ composition(S1elim)

Cubical Stability

Dimension substs. do not commute with evaluation!

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Dimension substs. do not commute with evaluation!

```
S1elim(a.A,
 loop{x}, B, y.L)   |————————→  L<x/y>
                                   |
                                   |  <0/x>
                                   ↓

                                 L<0/y>
```

Cubical Stability

Dimension substs. do not commute with evaluation!

```
S1elim(a.A,
 loop{x}, B, y.L)    |-------->  L<x/y>

        |                          |
        |  <0/x>                   |  <0/x>
        v                          v

S1elim(a.A,
 base, B, y.L)    |--> B   <=??=>  L<0/y>
```

Cubical Stability

Dimension substs. do not commute with evaluation!

```
S1elim(a.A,
 loop{x}, B, y.L)   |-----------> L<x/y>

        |                            |
        | <0/x>                      | <0/x>
        ↓                            ↓

S1elim(a.A,
 base, B, y.L)   |--> B   <=??=> L<0/y>
```

Restrict our theory to only cubically stable parts

Cubical Type Theory

stability: consider every substitution

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$$A \doteq B \text{ type } [\Psi]$$

dim context

A and B stably recognize the same stable values and have stably equal Kan structures

(see our arXiv papers)

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stability: consider every substitution

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dim context

A and B stably recognize the same stable values and have stably equal Kan structures

$$M \doteq N \in A\ [\Psi]$$

$A \doteq A$ type $[\Psi]$,
M and N stably eval to M' and N',
A stably treats M' and N' as the same

(see our arXiv papers)

Variables

Nuprl/...	Coq/Agda/...
Vars range over closed terms	Vars are indet.
Defined by transition b/w closed terms	Defined by conversion b/w open terms

exp vars dim vars

cubical computational TT

arXiv papers

CHTT Part I [AHW 2016]

Cartesian cubical + computational

CHTT Part II [AH 2017]

Dependent types

CHTT Part III [AFH 2017]

Univalent Kan universes

Strict equality

CHTT Part IV [AFH 2017]

Higher inductive types

Proof Assistants

RedPRL

In Nuprl style

redprl.org

redtt

(Work in progress)

github.com/RedPRL/redtt

---

yacctt

Proof of concept

modified from cubicaltt

github.com/mortberg/yacctt

Conclusion

We extended Nuprl semantics
by cubical structure which
justifies key features of HoTT

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Best of the two worlds!

We also built proof assistants

redprl.org
github.com/RedPRL/redtt
github.com/mortberg/yacctt