

The value of errors in proofs

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Plan

Proofs and computations

The value of errors in computations

The value of errors in proofs

Scientific impact: CS, Math, Physics, Optimization,...

Practical impact: Crypto, Clouds, Blockchains,...

Conceptual impact: Paradoxical properties of proofs

The value of the complexity theory methodology

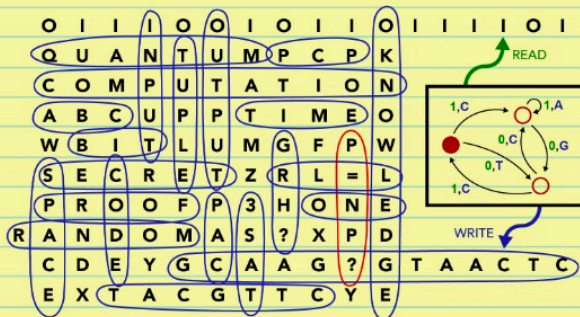
Modelling, classification, completeness, impracticality...

Book ad

MATHEMATICS + COMPUTATION

A THEORY REVOLUTIONIZING
TECHNOLOGY AND SCIENCE

Avi Wigderson



- Published by Princeton University Press
- Free (forever) on my website
- Comments welcome!

Proofs and Computations

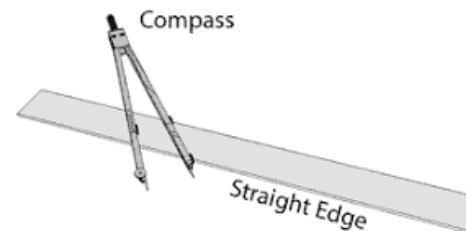
Two points in a very long history

[Euclid, 300BC]: *The Elements*

{Proofs of theorems in Plane Geometry,
deducible from 5 simple axioms}

=

{Constructions of planar point sets using
Straightedge and Compass



Turn of the 20th century

Hilbert's dream: Truth = Provability = Computability

...shattered

[Gödel '31] Incompleteness Thm (= is wrong)

[Turing '36] Undecidability Thm (= is wrong)

Def: algorithms \leftrightarrow Turing Machines

Corollary:
Computer
revolution

Set $S \leftrightarrow$ Decision Problem { "is $x \in S$?" }

$R = \{ \text{Sets computable by finite algorithms} \}$

$RE = \{ \text{Sets provable to finite algorithms} \}$

Thm: $R \neq RE$

Halting = {TMs M which halt on empty input}

Polynomial time

$P = \{ \text{Sets computable by efficient algorithms} \}$

$NP = \{ \text{Sets provable to efficient algorithms} \}$

Otherwise, no cryptography!

Open: $P \neq NP?$

Examples: claims, arguments, proofs,
proof systems, provers, verifiers...

What is true? In real life? In math?
What is a convincing argument?

Claim:
"x ∈ S"

Volume comparison



Left



Right

Claim: Left > Right

Verification:

Fill Left with water (to the rim) and pour to Right
[if spills, ACCEPT, else, REJECT]

Sudoku

Claim: This puzzle is solvable

		8	6					
							6	
			4	8			2	3
		5		9				8
	4	9					2	1
2				4		7		
3	6			2	9			
	1							
						5	1	

Argument:

9	2	8	6	1	3	4	5	7
4	7	3	9	5	2	8	6	1
1	5	6	4	8	7	9	2	3
7	3	5	2	9	1	6	4	8
6	4	9	7	3	8	2	1	5
2	8	1	5	4	6	7	3	9
3	6	7	1	2	9	5	8	4
5	1	2	8	7	4	3	9	6
8	9	4	3	6	5	1	7	2

Verification: Check each row, column, square,
AND that consistent with input. **ACCEPT/REJECT**

Composite numbers

Claim: 147573952588676412927 composite

Argument: 193707721, 761838257287

Verification: Check if

$$193707721 \times 761838257287 = 147573952588676412927$$

Again....

Volume comparison



Claim: Left > Right

Verification: General Procedure

Fill Left with water (to the rim) and pour to Right
[if spills, ACCEPT, else, REJECT]

Sudoku

Claim: This puzzle is solvable

		8	6					
							6	
			4	8			2	3
		5			9			8
	4	9					2	1
2				4		7		
3	6			2	9			
	1							
						5	1	

Argument:

General

9	2	8	6	1	3	4	5	7
4	7	3	9	5	2	8	6	1
1	5	6	4	8	7	9	2	3
7	3	5	2	9	1	6	4	8
6	4	9	7	3	8	2	1	5
2	8	1	5	4	6	7	3	9
3	6	7	1	2	9	5	8	4
5	1	2	8	7	4	3	9	6
8	9	4	3	6	5	1	7	2

Verification: Check each row, column, square, AND that consistent with input. **ACCEPT/REJECT**

Efficient algorithm: simple pattern matching

1			2	3	4			12		6				7	
		8				7			3			9	10	6	11
	12			10			1		13		11			14	
3			15	2			14				9			12	
13				8			10		12	2		1	15		
	11	7	6				16				15			5	13
			10		5	15			4		8			11	
16			5	9	12			1						8	
	2						13			12	5	8			3
	13			15		3			14	8		16			
5	8			1				2				13	9	15	
		12	4		6	16		13			7				5
	3			12				6			4	11			16
	7			16		5		14			1			2	
11	1	15	9			13			2				14		
	14				11		2			13	3	5			12

Composite numbers

Claim: 147573952588676412927 composite

Argument: 193707721, 761838257287

Crypto rests on the difficulty of finding such

Verification: Check if *General*

$193707721 \times 761838257287 = 147573952588676412927$

Efficient algorithm: simple arithmetic

Deductive proof systems

e.g. Peano Arithmetic

Numerous
others

Objects: Formulas/expressions over integers (A, B, \dots)

Axioms: E.g.

- $x+y = y+x$
- $x+1 > x$
- $(x+y)z = xz+yz$
- Induction Principle

You've got to believe/trust something!

Proofs are reductions of complex statements to simple truths via simple local sound steps

Deduction rules: E.g. if $A, A \rightarrow B$ true, then B is true.

Argument: A_1, A_2, \dots, A_m

Verification: Check that each A_i is an axiom, or follows from previous ones by a deduction rule.

Theorems:

- There are infinitely many primes
- Fermat's last theorem: no solution to $x^n+y^n=z^n, n>2$

Essentials of proof systems

Completeness: True claims have proofs

Soundness: False claims don't

Easy to check: Distinguishing convincing and faulty arguments by an efficient Verifier algorithm

A complexity theoretic view

Proof System [Cook-Reckhow '79]

An efficient Verifier $V(\text{claim}, \text{argument})$ satisfies:

Completeness: If claim is true then, for *some* argument
 $V(\text{claim}, \text{argument}) = \text{ACCEPT}$
(in which case $\text{claim} = \text{theorem}$, $\text{argument} = \text{proof}$)

Soundness: If claim is false then, for *every* argument
 $V(\text{claim}, \text{argument}) = \text{REJECT}$

T_V : the set of theorems in this system.

[CR'79] **NP** = $\{T_V : V \text{ deterministic}\}$

Probabilistic computation & error

"Axiom": Nature provides free access to randomness
(so, let algorithms make random choices!)

Def: A deterministic algorithm A computes a function f if
for all x , $A(x) = f(x)$ **always**

Def: A probabilistic algorithm B computes a function f if
for all x , $B(x) = f(x)$ **WHP** (eg $> 2/3$) **errors in algs**

Error can be efficiently reduced arbitrarily!

$$\Pr[B(x) \neq f(x)] < 1/3 \rightarrow \forall k, \Pr[B_k(x) \neq f(x)] < \exp(-k)$$

Rationale for allowing errors: (1) "Axiom" reasonable, and
(2) We tolerate uncertainty in life, why not in algs?

Value of allowing errors: Solve many more problems, ++

Probabilistic Proof System

[Babai '85, Goldwasser-Micali-Rackoff '85]

probabilistic

An efficient Verifier $V(\text{claim}, \text{argument})$ satisfies:

Completeness: If claim is true then, for *some* argument

$V(\text{claim}, \text{argument}) = \text{ACCEPT}$ **always**

(in which case $\text{claim}=\text{theorem}$, $\text{argument}=\text{proof}$)

Soundness: If claim is false then, for *every* argument

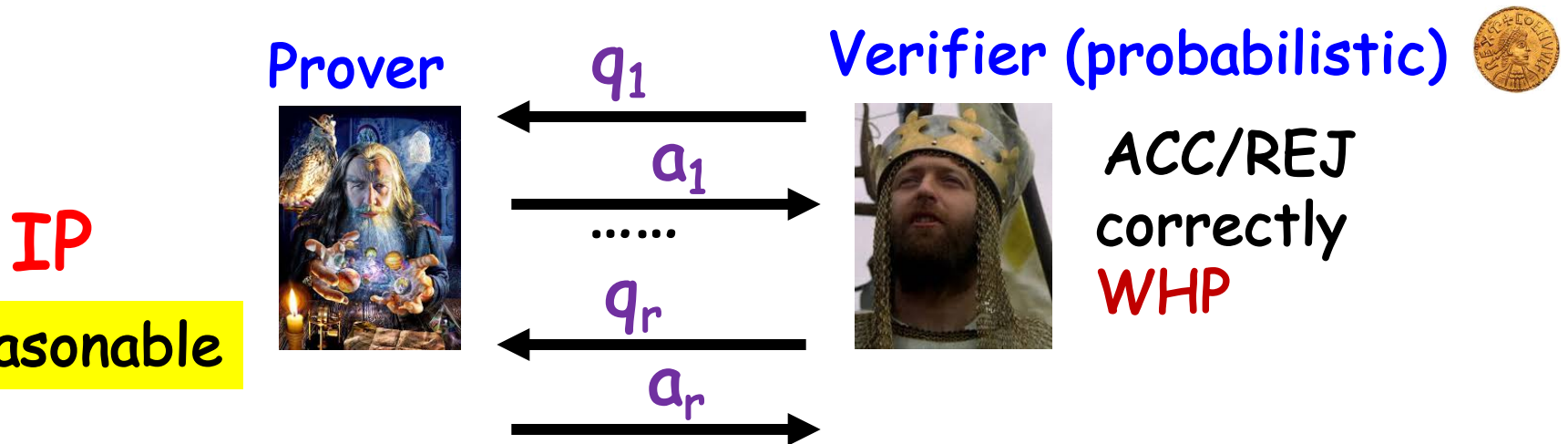
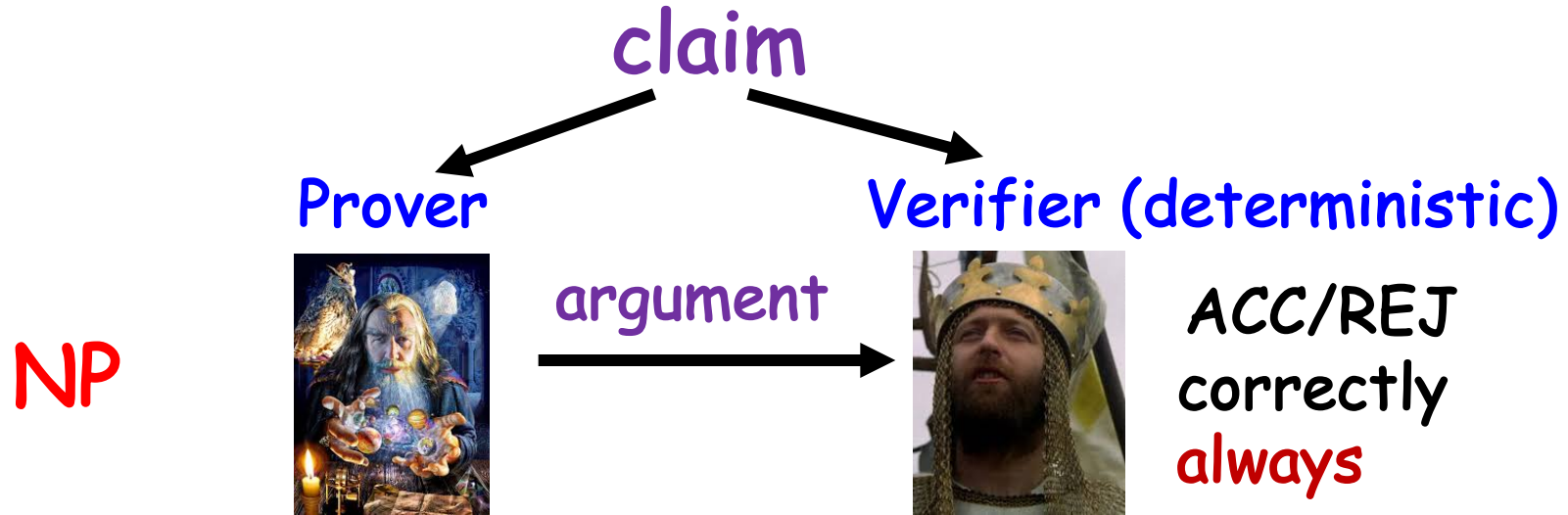
$V(\text{claim}, \text{argument}) = \text{REJECT}$ **WHP** **errors in proofs**

T_V : the set of theorems in this system.

IP $\triangleq \{T_V : V \text{ probabilistic} + \textit{interactive}\}$

IP = (Probabilistic) Interactive Proofs

[Babai'85, Goldwasser-Micali-Rackoff'85]



A revolutionary scientific notion!

Value of errors in proofs: Impact of interactive proofs

Conceptual

Scientific

Mathematical

Technological

Proof with paradoxical properties

ZK: Convincing proofs need not convey information

PCP: Convincing proofs need not be read

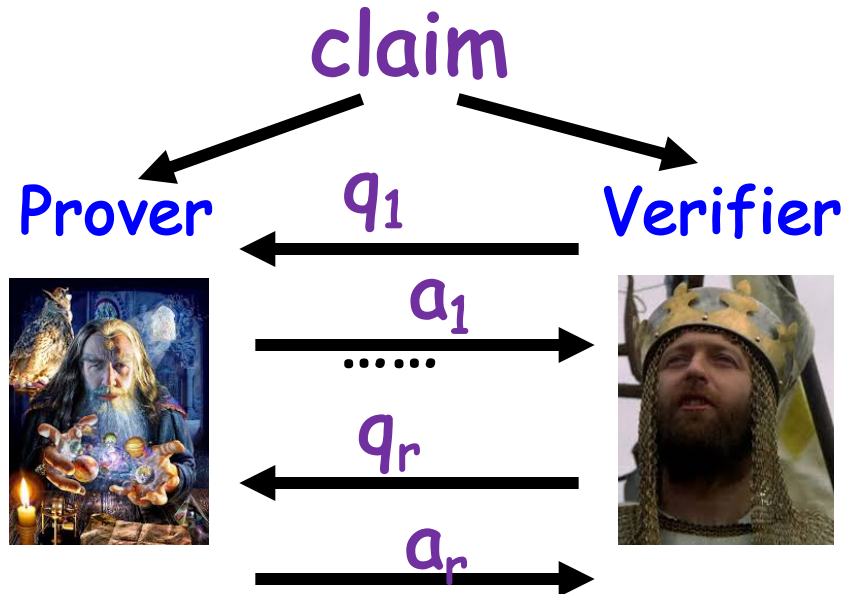
The amazing journey from **ZK** to **PCP**

ZKIP: Zero-Knowledge Interactive Proofs

[Goldwasser-Micali-Rackoff '85]

ZKIP

Formal def
non-trivial!



ZKIP =
IP +

V ACC. \rightarrow
V learns
nothing else

Possible? Can a convincing proof be uninformative?

[Goldreich-Micali-Wigderson '86]

1-way functions exist \rightarrow NP \subseteq ZKIP

Every proof can be made into a ZK proof!

Crypto is used! Is it necessary?

ZK impacts

Crypto: [Goldreich-Micali-Wigderson '87, ...]

Cryptographic protocol design, completeness thm

Practical applications:

Anonymous cash, Blockchains, Public ledgers ...

Physical ZK proofs:

[Barak-Glaser-Goldstone'14] Nuclear disarmament

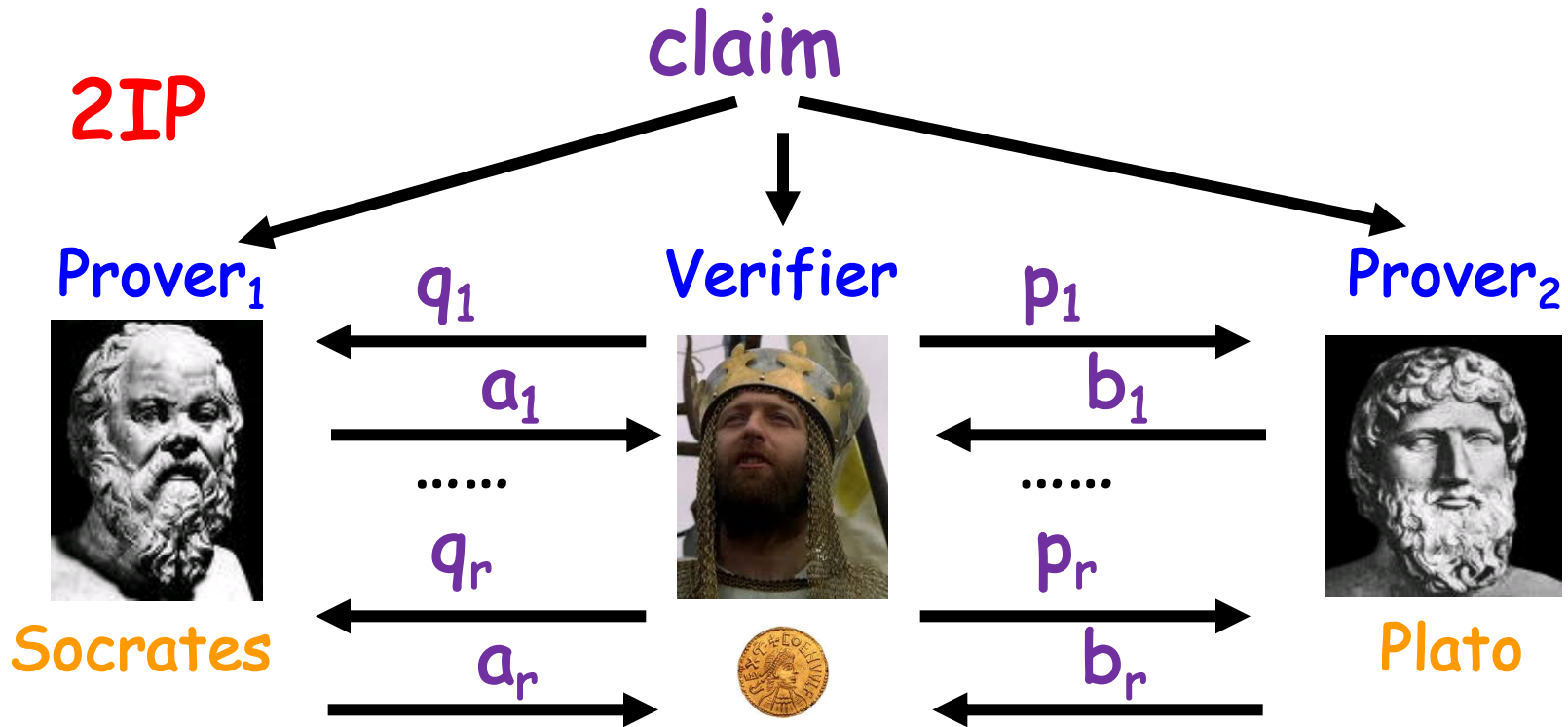
[Fisch-Freund-Naor '14] Anonymous DNA testing,...

New proof systems:

MIP: allowing multiple provers

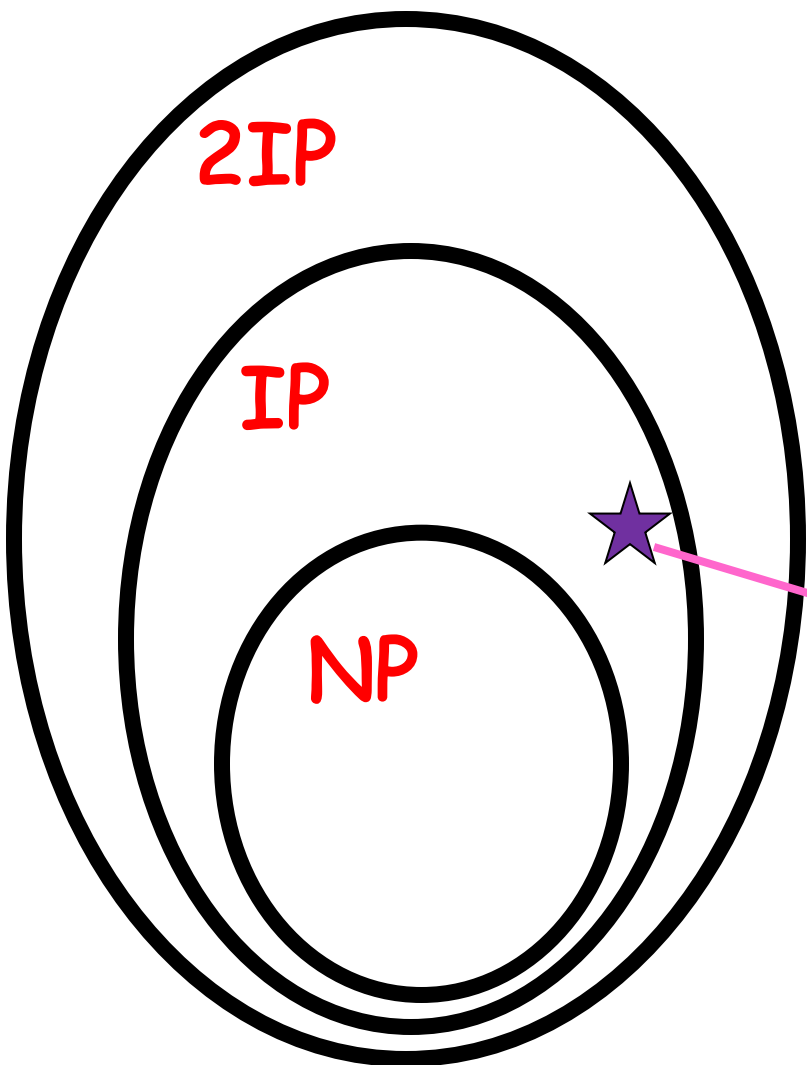
2IP: 2-Prover Interactive Proofs

[BenOr-Goldwasser-Kilian-Wigderson '89]



[BGKW '89] $NP \subseteq ZK$ 2IP
Physical separation replaces
computational assumptions

What is the power of **Randomness** and **Interaction** in Proofs?



Trivial inclusions

$$IP \subseteq PSPACE \quad \text{Polynomial Space}$$

$$2IP \subseteq NEXP \quad \text{Nondeterministic Exponential Time}$$

Few nontrivial examples

Graph non-isomorphism

.....

Few years of stalemate

wonders of polynomials

Avalanche of Characterizations

+ Conceptual meaning

[Lund-Fortnow-Karloff-Nisan, Shamir '90]

$IP = PSPACE$

Winning strategies are efficiently verifiable!

[Babai-Fortnow-Lund '91]

$2IP = NEXP$

Intractable problems are efficiently verifiable!

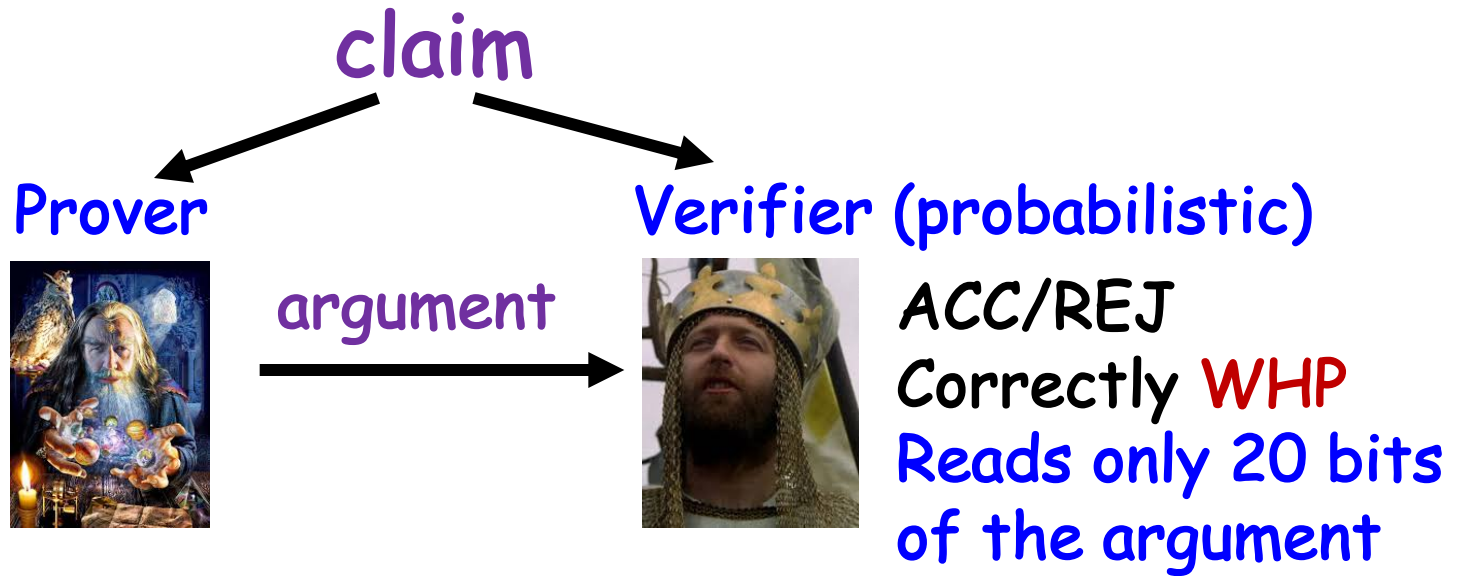
[Arora-Lund-Motwani-Safra-Sudan-Szegedy'92] $PCP = NP$

Written proofs verifiable from constant-size snapshots!

Same for transcripts of program execution

No crypto!

PCP (Probabilistically Checkable Proofs)



NP=PCP

Possible? Finding a single bug in a 100-page proof?

Yes!! Every proof can be turned into a PCP!

Optimization Hardness of approximation!

Coding theory

Complexity theory,...

Technology cloud computing, blockchains,...

Quantum computation

"Axiom": Nature provides access to quantum phenomena

[Manin '80, Feynman '82] Suggest building computers, that manipulate quantum superpositions with unitary operations.

[..., Deutsch '85, Bernstein-Vazirani '97,...] Formalize it.

BQP: efficient quantum algorithms (> probabilistic ones)

[Shor '94] Factoring, Discrete Log \in **BQP**

Frenzy attempts to develop:

- Supporting technology (billions invested)
- "Post-quantum" cryptography (e.g. harness assumptions)
- New quantum algorithms (not so much...)
- New models (plenty)

Quantum proof systems

Verifier is an efficient quantum algorithm

One prover: IP^* , Many provers: MIP^*

[Jain-Ji-Updahyay-Watrous'09] $IP^* = PSPACE$

[..., Ji-Natarajan-Vidick-Wright-Yuen'20] $MIP^* = RE$

Quantum Information theory: power of entanglement

Halting \equiv approx. the value of a non-local game!

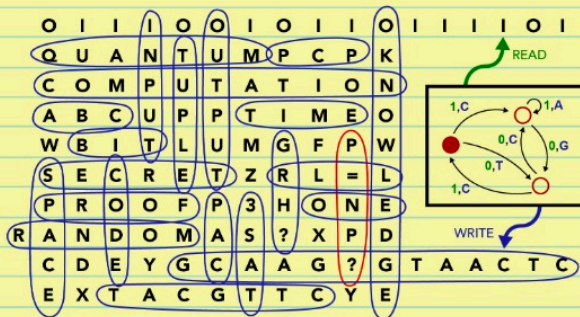
Math \rightarrow Connes' embedding conjecture in von-Neumann algebras is false!

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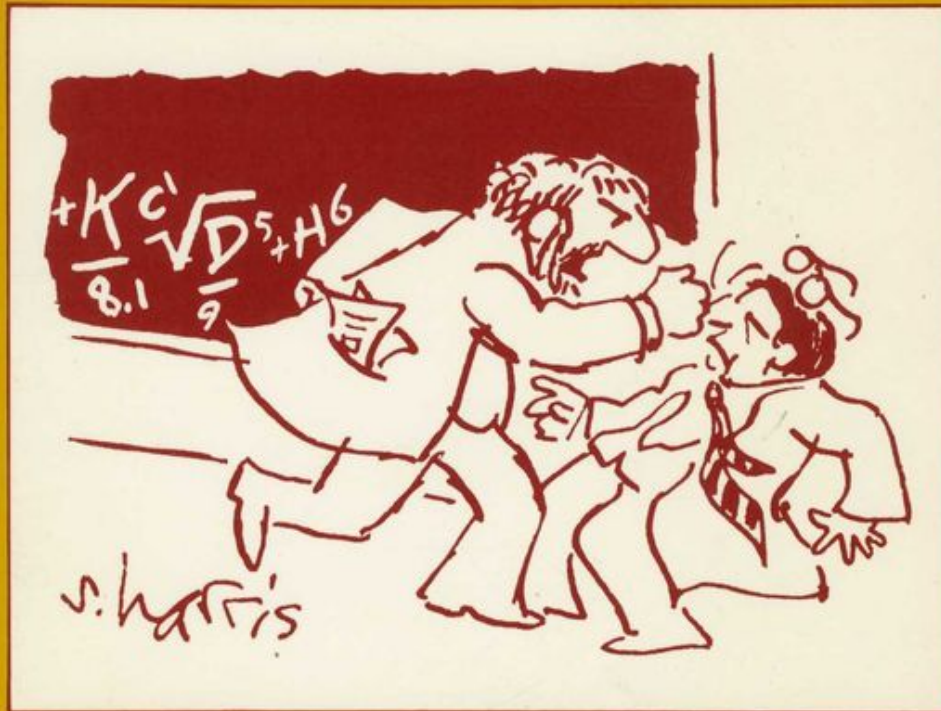
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"YOU WANT PROOF?
I'LL GIVE YOU PROOF!"



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There's no proof the FOREWORD is BY ALBERT EINSTEIN

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