

# Uniform Direct Product Theorems: Simplified, Optimized, and Derandomized

[Extended Abstract]

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## ABSTRACT

The classical Direct-Product Theorem for circuits says that if a Boolean function  $f : \{0, 1\}^n \rightarrow \{0, 1\}$  is somewhat hard to compute on average by small circuits, then the corresponding  $k$ -wise direct product function  $f^k(x_1, \dots, x_k) = (f(x_1), \dots, f(x_k))$  (where each  $x_i \in \{0, 1\}^n$ ) is significantly harder to compute on average by slightly smaller circuits. We prove a *fully uniform* version of the Direct-Product Theorem with information-theoretically *optimal* parameters, up to constant factors. Namely, we show that for given  $k$  and  $\epsilon$ , there is an efficient randomized algorithm  $A$  with the following property. Given a circuit  $C$  that computes  $f^k$  on at least  $\epsilon$  fraction of inputs, the algorithm  $A$  outputs with probability at least  $3/4$  a list of  $O(1/\epsilon)$  circuits such that at least one of the circuits on the list computes  $f$  on more than  $1 - \delta$  fraction of inputs, for  $\delta = O((\log 1/\epsilon)/k)$ . Moreover, each output circuit is an  $AC^0$  circuit (of size  $\text{poly}(n, k, \log 1/\delta, 1/\epsilon)$ ), with oracle access to the circuit  $C$ .

Using the Goldreich-Levin decoding algorithm [5], we also get a *fully uniform* version of Yao's XOR Lemma [18] with *optimal* parameters, up to constant factors. Our results simplify and improve those in [10].

Our main result may be viewed as an efficient approximate, local, list-decoding algorithm for direct-product codes (encoding a function by its values on all  $k$ -tuples) with opti-

mal parameters. We generalize it to a family of “derandomized” direct-product codes, which we call *intersection codes*, where the encoding provides values of the function only on a subfamily of  $k$ -tuples. The quality of the decoding algorithm is then determined by sampling properties of the sets in this family and their intersections. As a direct consequence of this generalization we obtain the first derandomized direct product result in the uniform setting, allowing hardness amplification with only constant (as opposed to a factor of  $k$ ) increase in the input length. Finally, this general setting naturally allows the decoding of concatenated codes, which further yields nearly optimal derandomized amplification.

## Categories and Subject Descriptors

F.2.0 [Analysis of Algorithms and Problem Complexity]: General; G.3 [Probability and Statistics]: probabilistic algorithms

## General Terms

Theory

## Keywords

Direct Product Theorem, Direct Product Code, XOR Code

## 1. INTRODUCTION

Applications such as cryptography and derandomization require reliably hard problems, ones that cannot be solved by any fast algorithm with even a non-trivial advantage over random guessing. Direct-product theorems are a primary tool in hardness amplification, allowing one to convert problems that are somewhat hard into problems that are more reliably hard. In a direct-product theorem, we start with a function  $f$  such that any feasible algorithm has a non-negligible chance of failing to compute  $f(x)$  given a random  $x$ . We then show that no feasible algorithm can, given multiple instances of the problem  $x_1, \dots, x_k$ , compute all of the values  $f(x_i)$ , with even a small probability of success. (Usually, the  $x_i$ 's are chosen independently, but there are also derandomized direct-product theorems where the  $x_i$ 's are chosen pseudo-randomly.) Many strong direct product theorems are known for non-uniform models, such as Boolean

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