

Lawvere theories and C-systems¹

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Abstract

In this paper we consider the class of l-bijective C-systems, i.e., C-systems for which the length function is a bijection. The main result of the paper is a construction of an isomorphism between two categories - the category of l-bijective C-systems and the category of Lawvere theories.

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1 Introduction

In this paper we consider the simplest class of C-systems (contextual categories) - the ones for which the length function $l : Ob(CC) \rightarrow \mathbb{N}$ is a bijection. It turns out that such C-systems are very closely related to Lawvere theories. We construct functors from the category of Lawvere theories to the category of l-bijective C-systems and a functor in the opposite direction and show in Theorem 6.1 that these functors are mutually inverse isomorphisms of the corresponding categories. We emphasize the unexpected aspect of this result which is that we obtain not simply an equivalence but an actual isomorphism of the categories.

This is essentially the first result on C-systems that includes the description of their homomorphisms.

Since this paper as well as other papers in the series on C-systems is expected to play a role in the mathematically rigorous construction of the simplicial univalent representation of the UniMath language and the Calculus of Inductive Constructions and since such a construction itself can not rely on the univalent foundations the paper is written from the perspective of the Zermelo-Fraenkel formalism.

Constructions and proofs of the paper do not use the axiom of excluded middle or the axiom of choice.

We use the diagrammatic order in writing compositions, i.e., for $f : X \rightarrow Y$ and $g : Y \rightarrow Z$ we write $f \circ g$ for the composition of f and g .

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We do not make precise the concept of a *universe* that we use for some of the statements of the paper. It would be certainly sufficient to assume that U is a Grothendieck universe. However, it seems likely that sets U satisfying much weaker conditions can be used both for the statements and for the proofs of our results.

2 The category of Lawvere theories

As usual, we let \mathbb{N} denote the set of natural numbers. For $m \in \mathbb{N}$ let $\mathbf{stn}(m) = \{i \in \mathbb{N} \mid 0 \leq i < m\}$ be the standard set with m elements. Let

$$\mathit{Mor}(\mathbb{F}) = \cup_{n,m \in \mathbb{N}} \mathit{Fun}(\mathbf{stn}(n), \mathbf{stn}(m))$$

where $\mathit{Fun}(X, Y)$ is the set of functions from X to Y . Since every function has a well defined domain and codomain we can define a category \mathbb{F} with the set of objects \mathbb{N} and the set of morphisms $\mathit{Mor}(\mathbb{F})$ such that for each n and m the set

$$\mathbb{F}(m, n) := \{f \in \mathit{Mor}(\mathbb{F}) \mid \mathit{dom}(f) = m \text{ and } \mathit{codom}(f) = n\}$$

equals to $\mathit{Fun}(\mathbf{stn}(m), \mathbf{stn}(n))$ and composition, when restricted to these subsets, is the composition of functions.

For $m, n \in \mathbb{N}$ let $\iota_1^{m,n} : \mathbf{stn}(m) \rightarrow \mathbf{stn}(m+n)$ and $\iota_2^{m,n} : \mathbf{stn}(n) \rightarrow \mathbf{stn}(m+n)$ be the injections of the initial segment of length m and the concluding segment of length n .

Definition 2.1 [2015.11.24.def1] *A Lawvere theory structure on a category T is a functor $L : \mathbb{F} \rightarrow T$ such that the following conditions hold:*

1. [2017.05.31.it1] *L is a bijection on the sets of objects,*
2. [2017.05.31.it2] *$L(0)$ is an initial object of T ,*
3. [2017.05.31.it3] *for any $m, n \in \mathbb{N}$ the square*

$$\begin{array}{ccc} L(0) & \longrightarrow & L(n) \\ \downarrow & & \downarrow L(\iota_2^{m,n}) \\ L(m) & \xrightarrow{L(\iota_1^{m,n})} & L(m+n) \end{array} \quad \text{is a pushout square.}$$

A Lawvere theory is a pair (T, L) where T is a category and L is a Lawvere theory structure on T .

Let us denote the set of Lawvere theory structures on T by $Lw(T)$. Lawvere theories can be constructed in the following two steps.

Let T be a category and $L : \mathbb{F} \rightarrow T$ a functor that satisfies the conditions 2.1.2 and 2.1.3, but not necessarily 2.1.1. Define a new category T_L as follows. We set $Ob(T_L) = \mathbb{N}$ and

$$\mathit{Mor}(T_L) = \coprod_{m,n \in \mathbb{N}} \mathit{Mor}_T(L(m), L(n))$$

Then $\mathit{Mor}_{T_L}(m, n)$ is the set of iterated pairs $((m, n), f)$ where $f \in \mathit{Mor}_T(L(m), L(n))$ and by the usual abuse of notations we may identify it with the set $\mathit{Mor}_T(L(m), L(n))$. Consider the functor data $L' : \mathbb{F} \rightarrow T_L$ given on objects by $L'_{Ob} = Id_{\mathbb{N}}$ and on morphism by $L'_{Mor}(a) = ((m, n), L(a))$ for $a \in \mathit{Mor}_{\mathbb{F}}(m, n)$.

Lemma 2.2 [2017.05.31.11] *The functor data (L'_{Ob}, L'_{Mor}) is a functor and moreover a Lawvere theory structure on T_L .*

Proof: The functor axioms for L imply immediately the functor axioms for L' . It is easy to prove that L' satisfied 2.1.2 and 2.1.3 since L does. In addition L' is a bijection (identity) on objects. Therefore L' is a Lawvere theory structure.

The second step allows one to construct easily functors $L : \mathbb{F} \rightarrow T$ that satisfy 2.1.2 and 2.1.3. Let T be any category, $\mathbf{0}$ an initial object of T , X and object of T and suppose that for all Y in T we are given a pushout square of the form

$$\begin{array}{ccc} \mathbf{0} & \longrightarrow & X \\ \downarrow & & \downarrow \iota_2^{Y,X} \\ Y & \xrightarrow{\iota_1^{Y,X}} & Y \amalg X \end{array}$$

Example 2.3 [2017.05.28.ex1] Lawvere theories are ubiquitous. For example, given any object X in a category with an initial object $\mathbf{0}$ and coproducts of the form $(Y + X, \iota_1^{Y,X}, \iota_2^{Y,X})$ for all Y we obtain a Lawvere theory in the sense of Def. 2.1 by taking the opposite to the subcategory of objects of the form $\mathbf{0}, X, X + X, (\dots(X + X) \cdots + X)$ with the inductively defined functor L . By duality the same construction works when we have a final object pt and binary products $(Y \times X, \pi_1^{Y,X}, \pi_2^{Y,X})$.

Problem 2.4 [2015.12.18.prob2] *For a universe U , to construct a category $LW(U)$ of Lawvere theories in U .*

Construction 2.5 [2015.12.18.constr2] Following Lawvere [4, p.61] we define a morphism from a Lawvere theory $\mathbf{T}_1 = (T_1, L_1)$ to a Lawvere theory $\mathbf{T}_2 = (T_2, L_2)$ as a functor $G : T_1 \rightarrow T_2$ such that $L_1 \circ G = L_2$. We let $Hom_{LW}(\mathbf{T}_1, \mathbf{T}_2)$ denote the subset in the set of functors from T_1 to T_2 that are morphisms of the Lawvere theories.

Note that here one uses the equality rather than isomorphism of functors. The composition of morphisms is defined as composition of functors. The identity morphism is the identity functor. The associativity and the left and right unity axioms follow immediately from the corresponding properties of the composition of functors.

We let $Ob(LW(U))$ denote the set of Lawvere theories in U and $Mor(LW(U))$ the set

$$Mor(LW(U)) = \coprod_{\mathbf{T}_1, \mathbf{T}_2 \in Ob(LW(U))} Hom_{LW}(\mathbf{T}_1, \mathbf{T}_2)$$

Together with the obvious domain, codomain, identity and composition functions the pair of sets $Ob(LW(U))$ and $Mor(LW(U))$ forms a category that we denote $LW(U)$ and call the category of Lawvere theories in U .

Note that a morphism from \mathbf{T}_1 to \mathbf{T}_2 in this category is not a morphism of Lawvere theories but an iterated pair $((\mathbf{T}_1, \mathbf{T}_2), G)$ where G is a morphism of Lawvere theories. However, there is an obvious bijection

$$Mor_{LW(U)}(\mathbf{T}_1, \mathbf{T}_2) \rightarrow Hom_{LW}(\mathbf{T}_1, \mathbf{T}_2)$$

and by the common abuse of notation every time we have an expression which denotes an element of one of these sets in a position where an element of the other is expected it is assumed to be replaced by its image under this bijection.

The following lemma will be used below in Construction 5.9.

Lemma 2.6 [2015.11.28.14] *Let T be a category and $L : \mathbb{F} \rightarrow T$ a functor such that the following conditions hold:*

1. $L(0)$ is an initial object of T ,
2. for any $m \in \mathbb{N}$ the square

$$\begin{array}{ccc}
 \text{[2015.11.28.eq2]} & L(0) & \longrightarrow & L(1) \\
 & \downarrow & & \downarrow L(\iota_2^{m,1}) \\
 & L(m) & \xrightarrow{L(\iota_1^{m,1})} & L(m+1)
 \end{array} \tag{1}$$

is a pushout square.

Then for any $m, n \in \mathbb{N}$ the square

$$\begin{array}{ccc}
 L(0) & \longrightarrow & L(n) \\
 \downarrow & & \downarrow L(\iota_2^{m,n}) \\
 L(m) & \xrightarrow{L(\iota_1^{m,n})} & L(m+n)
 \end{array}$$

is a pushout square.

Proof: Let $m, n \in \mathbb{N}$. Consider first the diagram

$$\begin{array}{ccccc}
 L(0) & \longrightarrow & L(n) & \xrightarrow{L(\iota_2^{m,n})} & L(m+n) \\
 \downarrow & & \downarrow L(\iota_1^{n,1}) & & \downarrow L(\iota_1^{m+n,1}) \\
 L(1) & \xrightarrow{L(\iota_2^{n,1})} & L(n+1) & \xrightarrow{L(\iota_2^{m,n+1})} & L(m+n+1)
 \end{array}$$

Then first square is the reflection relative to the diagonal of a square of the form (1) and therefore it is a pushout square.

We have $\iota_2^{n,1} \circ \iota_2^{m,n+1} = \iota_2^{m+n,1}$. Therefore the large square is the reflection relative to the diagonal of a square of the form (1) and therefore it is a pushout square.

The right hand side square is commutative.

By the general properties of pushout squares which are obtained from the similar properties of the pullback squares by duality we conclude that the right hand side square is a pushout square.

To prove the lemma proceed now by induction on n .

For $n = 0$ the horizontal arrows are isomorphisms and since the square commutes it is a pushout square.

For $n = 1$ it is a square of the form (1).

For the successor consider the diagram

$$\begin{array}{ccccc}
 L(0) & \longrightarrow & L(n) & \xrightarrow{L(\iota_1^{n,1})} & L(n+1) \\
 \downarrow & & \downarrow L(\iota_2^{m,n}) & & \downarrow L(\iota_2^{m,n+1}) \\
 L(m) & \xrightarrow{L(\iota_1^{m,n})} & L(m+n) & \xrightarrow{L(\iota_1^{m+n,1})} & L(m+n+1)
 \end{array}$$

The first square is pushout by the inductive assumption. The second square is pushout by the first part of the proof. Therefore the ambient square is pushout. Since $\iota_1^{m,n} \circ \iota_1^{m+n,1} = \iota_1^{m,n+1}$ this completes the proof of the lemma.

3 The category of l-bijective C-systems

For the definition of a C-system see [2], [3] (where they are called contextual categories) as well as [6]. A C-system structure on a category CC is a six-tuple $cs = (l, \text{pt}, \text{ft}, \text{p}, \text{q}, \text{s})$ where l, ft and p are functions, pt an element of $Ob(CC)$, q a partial function on pairs and s a partial function. To be a C-system structure these objects must satisfy the conditions of [6, Definitions 2.1 and 2.3]. One of the main structures of a C-system is that for any X such that $l(X) > 0$ and any $f : Y \rightarrow \text{ft}(X)$ there is given an object $f^*(X)$ and a morphism $\text{q}(f, X) : f^*(X) \rightarrow X$ such that

$$[\mathbf{2017.05.31.eq1}] \text{ft}(f^*(X)) = Y \tag{2}$$

and the square

$$\begin{array}{ccc} f^*(X) & \xrightarrow{\text{q}(f,X)} & X \\ \text{p}_{f^*(X)} \downarrow & & \downarrow \text{p}_X \\ Y & \xrightarrow{f} & \text{ft}(X) \end{array}$$

is a pullback. Squares of this form will be called the *canonical squares* of a C-system.

Definition 3.1 [$\mathbf{2015.11.24.def2}$] A l-bijective C-system is a C-system such that the length function $Ob(CC) \rightarrow \mathbb{N}$ is a bijection.

We let $Cs_{\mathbb{N}}(CC)$ denote the set of l-bijective C-system structures on a category CC .

Problem 3.2 [$\mathbf{2015.12.18.prob3}$] Let U be a universe. To construct a category $\mathbf{CSys}(U)$ of C-systems in U .

Construction 3.3 A morphism of C-systems is a functor between the underlying categories that is compatible with the corresponding C-system structures. In particular a morphism of C-systems should commute with the length functions. For a detailed definition see [5, Definition 3.1]. We let $Hom_{CS}(CC_1, CC_2)$ denote the set of homomorphisms from the C-system CC_1 to the C-system CC_2 .

That the composition of functors that are homomorphisms is again a homomorphism is stated in [5, Lemma 3.2]. That the identity functor is a homomorphism is very easy to prove. The associativity and the left and right unity axioms for the composition of homomorphisms follow directly from the similar properties of the composition of functors.

Repeating the approach that we used with Lawvere theories we obtain the category $\mathbf{CSys}(U)$ of C-systems in U .

We let $\mathbf{CSys}_{\mathbb{N}}(U)$ denote the full subcategory in $\mathbf{CSys}(U)$ that consists of l-bijective C-systems.

The following result shows that l-bijective C-systems are abundant in the world of C-systems.

Theorem 3.4 [$\mathbf{2015.12.26.th1}$] Let C be a C-system. Let X be an object of C such that $l(X) = 1$. Let C_X be the smallest C-subsystem of C that is a full subcategory of C and that contains X . Then C_X is an l-bijective C-system.

Proof: Recall that we let \mathbf{pt} denote the object of C of length 0 and for an object Y of C we let π_Y denote the unique morphism $Y \rightarrow \mathbf{pt}$. Define objects X^{*n} for $n \in \mathbb{N}$ inductively as follows:

1. $X^{*0} = \mathbf{pt}$,
2. $X^{*1} = X$,
3. $X^{*(n+1)} = (\pi_{X^{*n}})^*(X)$.

Let C_X be the full subcategory of C that is generated by objects X^{*n} for $n \in \mathbb{N}$. Let us show that it is a C -subsystem of C . Since it is a full subcategory, i.e., contains all morphisms between its objects it is sufficient to show that it is closed under the operations of C -systems that generate objects.

We have $\mathbf{pt} \in C_X$.

For $X^{*n} \in C_X$ we have $\mathbf{ft}(X^{*n}) \in C_X$. Indeed, $\mathbf{ft}(X) = \mathbf{pt}$ because $l(X) = 1$ and $\mathbf{ft}(X^{*(m+1)}) = X^{*m}$ by (2).

For $f : X^{*m} \rightarrow X^{*n}$ we have $f^*(X^{*(n+1)}) \in C_X$. Indeed, by definition we have $X^{*(n+1)} = (\pi_{X^{*n}})^*(X)$. Therefore

$$f^*(X^{*(n+1)}) = f^*((\pi_{X^{*n}})^*(X)) = (f \circ \pi_{X^{*n}})^*(X) = \pi_{X^{*m}}^*(X) = X^{*(m+1)}$$

To finish the proof of the theorem it remains to show that the restriction of l to C_X is a bijection which is straightforward.

4 A functor from Lawvere theories to l -bijective C -systems

Problem 4.1 [2015.11.24.probl] For a category T to construct a function

$$LC : Lw(T) \rightarrow Cs_{\mathbb{N}}(T^{op})$$

from the Lawvere theory structures on T to the l -bijective C -system structures on T^{op} .

Construction 4.2 [2015.11.24.constr1] Let $CC = T^{op}$. We need to construct a l -bijective C -system structure on CC . We set:

The length function $l = L^{-1}$.

The distinguished final object \mathbf{pt} is $L(0)$.

The map $\mathbf{ft} : Ob(CC) \rightarrow Ob(CC)$ maps \mathbf{pt} to \mathbf{pt} and any object X such that $l(X) > 0$ to $L(l(X) - 1)$.

For \mathbf{pt} the morphism $\mathbf{p}_{\mathbf{pt}}$ is the identity. For X such that $l(X) > 0$ the morphism $\mathbf{p}_X : X \rightarrow \mathbf{ft}(X)$ is $L(\iota_1^{l(X)-1,1})$.

To define $\mathbf{q}(f, X)$ observe first that for any X such that $l(X) > 0$ we have a pullback square in CC of the form

$$\begin{array}{ccc}
 X & \xrightarrow{L(\iota_2^{l(X)-1,1})} & L(1) \\
 \mathbf{p}_X \downarrow & & \downarrow \\
 \mathbf{ft}(X) & \longrightarrow & L(0)
 \end{array} \tag{3}$$

Given $f : Y \rightarrow X$ we set

$$f^*(X) = L(l(Y) + 1)$$

Since (3) is a pullback square and $L(0)$ is a final object there is a unique morphism $q(f, X) : f^*(X) \rightarrow X$ such that

$$[2015.12.02.eq7] q(f, X) \circ p_X = p_{f^*(X)} \circ f \quad (4)$$

and

$$[2015.12.02.eq8] q(f, X) \circ L(\iota_2^{l(X)-1,1}) = L(\iota_2^{l(Y),1}) \quad (5)$$

Let us check the conditions of [6, Definition 2.1] that will show that we obtained a C0-system. We have $l^{-1}(0) = L(0) = \{\text{pt}\}$. For X such that $l(X) > 0$ we have $l(\text{ft}(X)) = l(L(l(X) - 1)) = l(X) - 1$. We also have $\text{ft}(\text{pt}) = \text{pt}$. The object $\text{pt} = L(0)$ is final.

The square

$$[2015.11.24.eq2] \begin{array}{ccc} f^*(X) & \xrightarrow{q(f,X)} & X \\ \downarrow & & \downarrow p_X \\ Y & \xrightarrow{f} & \text{ft}(X) \end{array} \quad (6)$$

commutes by the construction of $q(f, X)$.

If $f = Id_{\text{ft}(X)}$ then $l(Y) = l(X)$ and therefore $f^*(X) = X$. Therefore $q(f, X) \circ p_X = p_X$ and $q(f, X) \circ L(\iota_2^{l(X)-1,1}) = L(\iota_2^{l(X)-1,1})$ which proves that $q(f, X) = Id_X$.

Given $g : Z \rightarrow Y$ we have to verify that $q(g \circ f, X) = q(g, f^*(X)) \circ q(f, X)$. We have $(g \circ f)^*(X) = L(l(Z) + 1) = g^*(f^*(X))$. Taking into account that (3) is a pullback square, it remains to verify two equalities

$$q(g \circ f, X) \circ p_X = q(g, f^*(X)) \circ q(f, X) \circ p_X$$

and

$$q(g \circ f, X) \circ L(\iota_2^{l(X)-1,1}) = q(g, f^*(X)) \circ q(f, X) \circ L(\iota_2^{l(X)-1,1})$$

For the first equality we have

$$\begin{aligned} q(g \circ f, X) \circ p_X &= p_{(g \circ f)^*(X)} \circ g \circ f = p_{g^*(f^*(X))} \circ g \circ f = \\ &= q(g, f^*(X)) \circ p_{f^*(X)} \circ f = q(g, f^*(X)) \circ q(f, X) \circ p_X \end{aligned}$$

and for the second

$$q(g \circ f, X) \circ L(\iota_2^{l(X)-1,1}) = L(\iota_2^{l(Z),1}) = q(g, f^*(X)) \circ L(\iota_2^{l(Y),1}) = q(g, f^*(X)) \circ q(f, X) \circ L(\iota_2^{l(X)-1,1})$$

According to [6, Proposition 2.4] it remains to show that the squares (6) are pullback squares.

Consider the diagram

$$\begin{array}{ccccc} f^*(X) & \xrightarrow{q(f,X)} & X & \xrightarrow{L(\iota_2^{l(X)-1,1})} & L(1) \\ p_{f^*(X)} \downarrow & & p_X \downarrow & & \downarrow \\ Y & \xrightarrow{f} & \text{ft}(X) & \longrightarrow & L(0) \end{array}$$

where both the right hand side square and the outside square are of the form (3) and in particular are pullback squares and the left hand side square has been shown to be commutative. Therefore, the left hand side square is a pullback square. We conclude, by [6, Proposition 2.4] that there exists a unique s such that $(l, \text{pt}, \text{ft}, p, q, s)$ is a C-system structure.

Lemma 4.3 [2015.12.06.prob1] *Let $G : (T_1, L_1) \rightarrow (T_2, L_2)$ be a morphism of Lawvere theories. Then the functor G^{op} is a homomorphism of C-systems $(T_1^{op}, LC(L_1)) \rightarrow (T_2^{op}, LC(L_2))$.*

Proof:[2015.12.06.constr1] For convenience we will write H instead of G^{op} . In view of [5, Lemma 3.4] it is sufficient to verify that H is compatible with the length function, distinguished final object, ft map, \mathfrak{p} -morphisms and \mathfrak{q} -morphisms.

The fact that $l_1 = H \circ l_2$ is equivalent to the fact that $L_1 \circ G = L_2$.

The fact that $H(L_1(0)) = L_2(0)$ follows from the same property of G .

The fact that $H(\text{ft}(X)) = \text{ft}(H(X))$ again follows from the same property of G .

The fact that $H(\mathfrak{p}_X) = \mathfrak{p}_{H(X)}$ follows from the fact that $L_1 \circ G = L_2$ on objects and on morphisms of the form $l_1^{n,1}$.

It remains to verify that for X such that $l(X) > 0$ and $f : Y \rightarrow \text{ft}(X)$ one has

$$H(\mathfrak{q}(f, X)) = \mathfrak{q}(H(f), H(X))$$

where the right hand side is defined because H is compatible with l and ft .

Morphism $\mathfrak{q}(H(f), H(X))$ is defined as the unique morphism such that

$$\mathfrak{q}(H(f), H(X)) \circ \mathfrak{p}_{H(X)} = \mathfrak{p}_{H(f)^*(H(X))} \circ H(f)$$

and

$$\mathfrak{q}(H(f), H(X)) \circ L_2(l_2^{l_2(H(X))-1,1}) = L_2(l_2^{l_2(H(Y)),1})$$

Therefore we need to verify the same equalities for the morphism $H(\mathfrak{q}(f, X))$. For the first equality we have

$$H(\mathfrak{q}(f, X)) \circ \mathfrak{p}_{H(X)} = H(\mathfrak{q}(f, X)) \circ H(\mathfrak{p}_X) = H(\mathfrak{q}(f, X) \circ \mathfrak{p}_X) = H(\mathfrak{p}_{f^*(X)} \circ f) = \mathfrak{p}_{H(f^*(X))} \circ H(f)$$

It remains to show that $H(f^*(X)) = H(f)^*(H(X))$. It follows from the fact that $f^*(X) = L_1(l_1(X) + 1)$ and $H(f)^*(H(X)) = L_2(l_2(H(X)) + 1)$.

For the second equality we have

$$\begin{aligned} H(\mathfrak{q}(f, X)) \circ L_2(l_2^{l_2(H(X))-1,1}) &= H(\mathfrak{q}(f, X)) \circ H(L_1(l_1^{l_1(X)-1,1})) = H(\mathfrak{q}(f, X) \circ L_1(l_1^{l_1(X)-1,1})) = \\ &= H(L_1(l_1^{l_1(Y),1})) = L_2(l_2^{l_2(H(Y)),1}) \end{aligned}$$

This completes the construction.

Problem 4.4 [2015.12.08.prob1] To construct, for any universe U , a functor $LC_U : LW(U) \rightarrow \mathbf{CSys}_{\mathbb{N}}(U)$.

Construction 4.5 [2015.12.08.constr1] We set LC_{Ob} to be the function that takes a Lawvere theory to the opposite category of its underlying category with the C-system structure defined by Construction 4.2. We set LC_{Mor} to be the function that takes a functor G that is a morphism of Lawvere theories to G^{op} . It is well defined by Lemma 4.3. That the functions (LC_{Ob}, LC_{Mor}) form a functor, i.e., commute with the identity morphisms and compositions is straightforward.

Remark 4.6 [2015.11.24.rem1] Let (T, L) be a Lawvere theory. Consider T^{op} . The morphism $p : L(1) \rightarrow L(0)$, pullback squares (3) and the final object $L(0)$ make T^{op} into a universe category (T^{op}, p) (see [5]). It is easy to prove that the C-system of Construction 4.2 is isomorphic to the C-system $CC(T^{op}, p)$ of this universe category. However this isomorphism is not an equality since the set of objects of the category $CC(T^{op}, p)$ is not equal to the set of objects of T^{op} . Indeed, an object of $CC(T^{op}, p)$ is a sequence of the form

$$(\mathbf{pt}; \pi_{L(0)}, \dots, \pi_{L(n)})$$

where $\pi_{L(i)}$ is the unique morphism $L(i) \rightarrow L(0)$.

5 A functor from l-bijective C-systems to Lawvere theories

Problem 5.1 [2015.11.24.prob2] For a category CC to construct a function

$$CL : Cs_{\mathbb{N}}(CC) \rightarrow Lw(CC^{op})$$

To perform a construction we will need a number of lemmas and intermediate constructions. Let us fix a category CC and a l-bijective C-system structure $cs = (l, \mathbf{pt}, \mathbf{ft}, \mathbf{p}, \mathbf{q}, \mathbf{s})$ on CC . We will often write CC both for the category and for the C-system (CC, cs) .

Problem 5.2 For $n \in \mathbb{N}$ and $i = 0, \dots, m-1$ to construct a morphism $\pi_i^m : l^{-1}(m) \rightarrow l^{-1}(1)$ in CC .

Construction 5.3 [2015.11.28.constr1] By induction on m .

For $m = 0$ there are no morphisms to construct.

For $m = 1$ we set $\pi_0^1 = Id_{l^{-1}(1)}$.

For the successor consider the canonical square:

$$\begin{array}{ccc}
 l^{-1}(m+1) & \xrightarrow{\mathbf{q}(\pi, l^{-1}(1))} & l^{-1}(1) \\
 \text{[2015.11.28.eq1]} \downarrow \mathbf{p}_{l^{-1}(m+1)} & & \downarrow \mathbf{p}_{l^{-1}(1)} \\
 l^{-1}(m) & \xrightarrow{\pi} & l^{-1}(0)
 \end{array} \tag{7}$$

where we use π to denote the unique morphisms from objects of CC to the final object $l^{-1}(0)$. We set

$$\pi_i^{m+1} = \begin{cases} \mathbf{p}_{l^{-1}(m+1)} \circ \pi_i^m & \text{for } i < m \\ \mathbf{q}(\pi, l^{-1}(1)) & \text{for } i = m \end{cases}$$

Problem 5.4 For any $m, n \in \mathbb{N}$ and a function $f : \mathbf{stn}(m) \rightarrow \mathbf{stn}(n)$ to construct a morphism $L_f : l^{-1}(n) \rightarrow l^{-1}(m)$ in CC .

Construction 5.5 [2015.11.28.constr2] By induction on m .

For $m = 0$ we set $L_f = \pi$.

For $m = 1$ we set $L_f = \pi_{f(0)}^n$.

For the successor consider $f : \text{stn}(m+1) \rightarrow \text{stn}(n)$ and the square (7). We define L_f as the unique morphism such that:

$$[\mathbf{2015.11.30.eq2}] L_f \circ \mathbf{p}_{|-1(m+1)} = L_{\iota_1^{m,1} \circ f} \quad (8)$$

and

$$[\mathbf{2015.11.30.eq3}] L_f \circ \mathbf{q}(\pi, |-1(1)) = L_{\iota_2^{m,1} \circ f} \quad (9)$$

where, let us recall,

$$\iota_1^{m,1} : \text{stn}(m) \rightarrow \text{stn}(m+1)$$

$$\iota_2^{m,1} : \text{stn}(1) \rightarrow \text{stn}(m+1)$$

are the morphism that define the representation

$$\text{stn}(m+1) = \text{stn}(m) \amalg \{m+1\}$$

Lemma 5.6 [2015.11.28.11] *Let $m, n \in \mathbb{N}$ and $f : \text{stn}(m) \rightarrow \text{stn}(n)$. Then for any $i = 0, \dots, m-1$ one has*

$$L_f \circ \pi_i^m = \pi_{f(i)}^n$$

Proof: By induction on m .

For $m = 0$ there is nothing to prove.

For $m = 1$ we need to prove that $L_f \circ \pi_0^n = \pi_{f(0)}^n$. By construction, $\pi_0^n = Id_{|-1(1)}$ and $L_f = \pi_{f(0)}^n$ which implies the goal.

For the successor $f : \text{stn}(m+1) \rightarrow \text{stn}(n)$ and the square (7) for $m+1$.

If $i = m$ then $\pi_i^{m+1} = \mathbf{q}(\pi, |-1(1))$ and by the construction of L_f we have $L_f \circ \mathbf{q}(\pi, |-1(1)) = \pi_{f(m)}^n$.

If $i < m$ then $\pi_i^{m+1} = \mathbf{p}_{|-1(m+1)} \circ \pi_i^m$. Therefore

$$L_f \circ \pi_i^{m+1} = L_f \circ \mathbf{p}_{|-1(m+1)} \circ \pi_i^m = L_{\iota_1^{m,1} \circ f} \circ \pi_i^m$$

and by the inductive assumption

$$L_{\iota_1^{m,1} \circ f} \circ \pi_i^m = \pi_{f(\iota_1^{m,1}(i))}^n = \pi_{f(i)}^n$$

The lemma is proved.

Lemma 5.7 [2015.11.28.12] *Let $m, n \in \mathbb{N}$ and $f, g : |-1(n) \rightarrow |-1(m)$ are two morphisms such that for all $i = 0, \dots, m-1$ one has*

$$f \circ \pi_i^m = g \circ \pi_i^m$$

Then $f = g$.

Proof: By induction on m .

For $m = 0$, $|-1(m)$ is a final object and $f = g$.

For $m = 1$, $\pi_0^1 = Id_{|-1(1)}$ and $f = f \circ Id = g \circ Id = g$.

For the successor the square (7) for $m+1$. Since the square is a pullback square it is sufficient to show that

$$f \circ \mathbf{p}_{|-1(m+1)} = g \circ \mathbf{p}_{|-1(m+1)}$$

and

$$f \circ \mathbf{q}(\pi, l^{-1}(1)) = g \circ \mathbf{q}(\pi, l^{-1}(1))$$

The second equality follows from the fact that $\mathbf{q}(\pi, l^{-1}(1)) = \pi_m^{m+1}$.

The first equality follows by the inductive assumption since for $i = 0, \dots, m-1$ we have

$$(f \circ \mathbf{p}_{l^{-1}(m+1)}) \circ \pi_i^m = f \circ \pi_i^{m+1}$$

and

$$(g \circ \mathbf{p}_{l^{-1}(m+1)}) \circ \pi_i^m = g \circ \pi_i^{m+1}$$

The lemma is proved.

Lemma 5.8 [2015.11.28.13]

1. for any $m \in \mathbb{N}$ one has $L_{Id_{\text{stn}(m)}} = Id_{l^{-1}(m)}$,
2. for $k, m, n \in \mathbb{N}$ and $f : \text{stn}(k) \rightarrow \text{stn}(m)$, $g : \text{stn}(m) \rightarrow \text{stn}(n)$ one has

$$L_{f \circ g} = L_g \circ L_f$$

Proof: Both cases follow in a straightforward way from Lemmas 5.7 and 5.6.

We can now provide a construction for Problem 5.1.

Construction 5.9 [2015.11.30.constr1] We need to construct a Lawvere theory structure on CC^{op} , i.e. a functor $L : \mathbb{F} \rightarrow CC^{op}$ satisfying the conditions of Definition 2.1. We define the object part of L as l^{-1} . We define the morphism part of L as $L_{Mor}(f) = L_f$. Lemma 5.8 shows that L is a covariant functor to CC^{op} .

The first condition of Definition 2.1 is obvious. The second condition as well (it follows from the axioms of a C-system). To prove the third condition we first apply Lemma 2.6. It remains to prove that squares of the form (1) are pushout squares in CC^{op} or, equivalently, that squares of the form

$$\begin{array}{ccc}
 l^{-1}(m+1) & \xrightarrow{L_{l_2^{m,1}}} & l^{-1}(1) \\
 \text{[2015.11.30.eq1]}_{L_{l_1^{1,m}}} \downarrow & & \downarrow \\
 l^{-1}(m) & \longrightarrow & l^{-1}(0)
 \end{array} \tag{10}$$

in CC are pullback squares. We will do it by showing that the square (10) equals to the canonical square of the C-system structure for the pair $(\pi_{l^{-1}(m)}, l^{-1}(1))$. The right hand side vertical morphism, $\pi_{l^{-1}(1)}$, is a unique morphism from $l^{-1}(1)$ to $l^{-1}(0)$ and since $\text{ft}(l^{-1}(1)) = l^{-1}(0)$ it equals $\mathbf{p}_{l^{-1}(1)}$.

It remains to show that

$$\text{[2015.12.02.eq1]}_{L_{l_1^{m,1}}} = \mathbf{p}_{l^{-1}(m+1)} \tag{11}$$

and

$$\text{[2015.12.02.eq2]}_{L_{l_2^{m,1}}} = \mathbf{q}(\pi_{l^{-1}(m)}, l^{-1}(1)) \tag{12}$$

These equalities follow from the equalities (8) and (9) for $f = Id_{l^{-1}(m+1)}$ because of Lemma 5.8(1). The construction is completed.

Next we will show that our function on objects extends to a functor from the category of l -bijective C -systems to the category of Lawvere theories. First we need the following lemma.

Lemma 5.10 [2015.12.08.13] *Let $H : CC_1 \rightarrow CC_2$ be a homomorphism of l -bijective C -systems. Then for any $n \in \mathbb{N}$ and $i = 0, \dots, n-1$ one has*

$$\text{[2015.12.08.eq3]} H(\pi_i^n) = \pi_i^n \tag{13}$$

Proof: Note that since we have $l_2(H(X)) = l_1(X)$ both sides of (13) are morphisms from $l_2^{-1}(n)$ to $l_2^{-1}(1)$.

The proof is by induction on n .

For $n = 0$ there are no equations to prove.

For $n = 1$ we have $\pi_0^1 = Id_{l^{-1}(1)}$ and the statement of the lemma follows from the identity axiom of the definition of a functor.

For the successor we have two cases. For $i < n$ we have

$$H(\pi_i^{n+1}) = H(\mathfrak{p}_{l^{-1}(n+1)} \circ \pi_i^n) = H(\mathfrak{p}_{l^{-1}(n+1)}) \circ H(\pi_i^n) = \mathfrak{p}_{l^{-1}(n+1)} \circ \pi_i^n = \pi_i^{n+1}$$

where the third equality uses the inductive assumption. For $i = n$ we have

$$H(\pi_n^{n+1}) = H(\mathfrak{q}(\pi, l^{-1}(1))) = \mathfrak{q}(\pi, l^{-1}(1)) = \pi_n^{n+1}$$

The lemma is proved.

Lemma 5.11 [2015.12.08.12] *Let $G : (CC_1, cs_1) \rightarrow (CC_2, cs_2)$ be a homomorphism of C -systems. Then the functor $G^{op} : CC_1^{op} \rightarrow CC_2^{op}$ is a morphism of Lawvere theories $(CC_1^{op}, CL(cs_1)) \rightarrow (CC_2^{op}, CL(cs_2))$.*

Proof: Let $CL(cs_1) = L_1$ and $CL(cs_2) = L_2$. We need to show that $L_1 \circ G^{op} = L_2$. The equality between the object components of these functors follows from the fact that a homomorphism of C -systems is compatible with the length functions. For the morphism component it is more convenient to consider equivalent equation

$$L'_1 \circ G = L'_2$$

where $L'_i : \mathbb{F}^{op} \rightarrow CC_i$. Then we have to show that for any $f : \text{stn}(m) \rightarrow \text{stn}(n)$ one has $G(L_{1,f}) = L_{2,f}$. Both sides of this equality are morphisms $l_2^{-1}(n) \rightarrow l_2^{-1}(m)$. By Lemma 5.7 it is sufficient to show that $G(L_{1,f}) \circ \pi_i^n = L_{2,f} \circ \pi_i^n$ for all $i = 0, \dots, n-1$. We have

$$G(L_{1,f}) \circ \pi_i^n = G(L_{1,f}) \circ G(\pi_i^n) = G(L_{1,f} \circ \pi_i^n) = G(\pi_{f(i)}^n) = \pi_{f(i)}^n$$

where we used Lemma 5.6 and twice Lemma 5.10. On the other hand

$$L_{2,f} \circ \pi_i^n = \pi_{f(i)}^n$$

again by Lemma 5.6. This completes the proof of Lemma 5.11.

Problem 5.12 [2015.12.08.prob2] *For any universe U to construct a functor $CL_U : \mathbf{CSys}_{\mathbb{N}}(U) \rightarrow LW(U)$.*

Construction 5.13 [2015.12.08.constr3] The object component of CL_U takes a C-system (CC, cs) to the Lawvere theory $(CC^{op}, CL(cs))$ where $CL(cs)$ is defined by Construction 5.9.

The morphism component takes a homomorphism $G : (CC_1, cs_1) \rightarrow (CC_2, cs_2)$ to G^{op} . It is well defined by Lemma 5.11.

The identity and composition axioms are straightforward from the corresponding properties of functor composition and its compatibility with the function $G \mapsto G^{op}$.

6 Isomorphism theorem

Theorem 6.1 [2015.12.08.th1] *For any universe U , Constructions 4.5 and 5.13 define mutually inverse isomorphisms between the categories of Lawvere theories in U and $\mathbb{1}$ -bijective C-systems in U .*

Proof: Let us show first that the object components of the functors LC_U and CL_U are mutually inverse bijections of sets. We write LC for $(LC_U)_{Ob}$ and similarly for CL .

Let us consider first the composition $LC \circ CL$ of the function of Construction 4.2 with the function of Construction 5.9.

Let Lw be a Lawvere theory structure on T and let

$$Lw' = (LC \circ CL)(Lw) = CL(LC(Lw))$$

We have $Lw = (L : \mathbb{F} \rightarrow T)$ and $Lw' = (L' : \mathbb{F} \rightarrow T)$. We have to prove that $L'(f) = L(f)$. To distinguish L' and L more clearly we will use the notation L'_f for $L'(f)$ as we did in Construction 5.9.

On objects we have

$$L'(n) = \mathbb{1}^{-1}(n) = (L^{-1})^{-1}(n) = L(n)$$

Let $f : \text{stn}(m) \rightarrow \text{stn}(n)$ be a morphism in \mathbb{F} .

Since L and L' coincide on objects both morphisms $L(f)$ and L'_f are of the form $L(n) \rightarrow L(m)$ when considered as contravariant functors to $CC = T^{op}$.

The proof is by induction on m .

If $m = 0$, $L(m) = L'(m)$ is a final object and any two morphisms with it as a codomain are equal.

If $m = 1$ we have by Constructions 5.9 and 5.5 that

$$L'(f) = L'_f = \pi_{f(0)}^n$$

We need to show that

$$L(f) = \pi_{f(0)}^n$$

We prove these equalities by induction on n .

If $n = 0$ then f exists only if $m = 0$ and we have already considered this case.

Let $(\mathbb{1}, \text{pt}, \text{ft}, \text{p}, \text{q}, \text{s})$ be the components of $LC(L)$. If $n = 1$ then by Construction 5.3 we have $\pi_{f(0)}^n = \pi_0^n = Id_{\mathbb{1}^{-1}(1)}$ and since L is a functor and $f = Id_{\mathbb{1}^{-1}(0)}$ we have $L(f) = Id_{\mathbb{1}^{-1}(1)}$.

For the successor we should consider diagram (7). We have that $\pi_{f(0)}^{n+1}$ is given by:

$$\pi_{f(0)}^{n+1} = \begin{cases} \text{p}\mathbb{1}^{-1}(n+1) \circ \pi_{f(0)}^n & \text{for } f(0) < n \\ \text{q}(\pi, \mathbb{1}^{-1}(1)) & \text{for } f(0) = n \end{cases}$$

Assume that $f(0) < n$. Then, by the inductive assumption we have $\pi_{f(0)}^n = L(f)$. By (11) we have $\mathfrak{p}_{l^{-1}(n+1)} = L(\iota_1^{n,1})$ and since L is a functor we get

$$\pi_{f(0)}^{n+1} = L(\iota_1^{n,1}) \circ \pi_{f(0)}^n = L(\iota_1^{n,1}) \circ L(g)$$

where $g : \mathbf{stn}(1) \rightarrow \mathbf{stn}(n)$ is such that $g(0) = f(0)$. Since L is a contravariant functor we get

$$L(\iota_1^{n,1}) \circ L(g) = L(g \circ \iota_1^{n,1}) = L(f)$$

For $f(0) = n$ we have $\pi_{f(0)}^{n+1} = \mathfrak{q}(\pi, l^{-1}(1)) = L(\iota_2^{n,1}) = L(f)$ where the second equality is by (12).

We have to consider now the case of the successor of m .

The morphism L'_f for $f : \mathbf{stn}(m+1) \rightarrow \mathbf{stn}(n)$ is defined in (8) and (9) as the unique morphism such that

$$L'_f \circ \mathfrak{p}_{l^{-1}(m+1)} = L'_{\iota_1^{m,1} \circ f} \quad (14)$$

and

$$L'_f \circ \mathfrak{q}(\pi, l^{-1}(1)) = L'_{\iota_2^{m,1} \circ f} \quad (15)$$

By the inductive assumption we have

$$L'_{\iota_1^{m,1} \circ f} = L(\iota_1^{m,1} \circ f)$$

$$L'_{\iota_2^{m,1} \circ f} = L(\iota_2^{m,1} \circ f)$$

It remains to prove that

$$[\mathbf{2015.12.02.eq5}] L(f) \circ \mathfrak{p}_{l^{-1}(m+1)} = L(\iota_1^{m,1} \circ f) \quad (16)$$

and

$$[\mathbf{2015.12.02.eq6}] L(f) \circ \mathfrak{q}(\pi_{l^{-1}(m)}, l^{-1}(1)) = L(\iota_2^{m,1} \circ f) \quad (17)$$

By (11) and (12) we have

$$\mathfrak{p}_{l^{-1}(m+1)} = L'_{\iota_1^{1,m}} = L(\iota_1^{1,m})$$

$$\mathfrak{q}(\pi_{l^{-1}(m)}, l^{-1}(1)) = L'_{\iota_2^{m,1}} = L(\iota_2^{m,1})$$

Where the second equalities are by the inductive assumption on m . This proves equalities (16) and (17).

Let us consider now the composition $CL \circ LC$. Fix a category CC in U . We need to show that for a l -bijective C -system structure cs one has $LC(CL(cs)) = cs$. Let $CL(cs) = (L)$ and let

$$cs = (l, \mathfrak{pt}, \mathfrak{ft}, \mathfrak{p}, \mathfrak{q}, \mathfrak{s})$$

and

$$LC(CL(cs)) = (l', \mathfrak{pt}', \mathfrak{ft}', \mathfrak{p}', \mathfrak{q}', \mathfrak{s}')$$

Then on objects $L = l^{-1}$ and $l' = L^{1-}$. Therefore $l = l'$.

Next $\mathfrak{pt} = l^{-1}(0) = (l')^{-1}(0) = \mathfrak{pt}'$. Similarly we see that $\mathfrak{ft} = \mathfrak{ft}'$ since on a l -bijective C -system \mathfrak{ft} is determined by l .

For X such that $l(X) > 0$ we have by Construction 4.2 that $\mathbf{p}'_X = L(\iota_1^{l(X)-1,1})$. Together with (11) we obtain

$$\mathbf{p}'_X = L(\iota_1^{l(X)-1,1}) = \mathbf{p}_{l^{-1}(l(X))}$$

since the C-system is l -bijective we have $X = l^{-1}(l(X))$ and therefore $\mathbf{p}' = \mathbf{p}$.

The morphism \mathbf{q}' is defined in Construction 4.2 as the unique morphism such that (4) and (5) hold. In our notation these equations take the form:

$$\mathbf{q}'(f, X) \circ \mathbf{p}'_X = \mathbf{p}'_{f^*(X)} \circ f \tag{18}$$

and

$$\mathbf{q}'(f, X) \circ L(\iota_2^{l(X)-1,1}) = L(\iota_2^{l(Y),1}) \tag{19}$$

We need to check that the same equations hold for \mathbf{q} . For the first one it follows immediately from the fact that $\mathbf{p}' = \mathbf{p}$.

To prove the second one consider equation (12). Applying this equation to (19) for \mathbf{q} and using the fact that of C-systems are l -bijective we get

$$\mathbf{q}(f, X) \circ \mathbf{q}(\pi_{\text{ft}(X)}, l^{-1}(1)) = \mathbf{q}(\pi_Y, l^{-1}(1))$$

which a particular case of the composition axiom for \mathbf{q} (see [6, Definition 2.1(7)]).

This proves that CL_U and LC_U are mutually inverse bijections on the sets of objects of our categories.

The fact that they give mutually inverse functions on morphisms between each pair of objects is straightforward. Indeed

$$CL_{U, Mor}(G) = G^{op}$$

and

$$LC_{U, Mor}(G) = G^{op}$$

as functors and since the mappings from morphisms of Lawvere theories and homomorphisms of C-systems to functors between the corresponding categories are injective we see that $CL_{U, Mor}$ and $LC_{U, Mor}$ are mutually inverse bijections.

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