## Lecture 14X: Pro-Objects

## March 19, 2018

Let  $\mathcal{C}$  and  $\mathcal{D}$  be categories which admit finite limits. We let  $\operatorname{Fun}^{\operatorname{lex}}(\mathcal{C}, \mathcal{D})$  denote the full subcategory of  $\operatorname{Fun}(\mathcal{C}, \mathcal{D})$  spanned by those functors which are *left exact*: that is, which preserve finite limits.

**Definition 1.** Let  $\mathcal{C}$  be an essentially small category which admits finite limits. We let  $\operatorname{Pro}(\mathcal{C})$  denote the category  $\operatorname{Fun}^{\operatorname{lex}}(\mathcal{C},\operatorname{Set})^{\operatorname{op}}$ . We will refer to the objects of  $\operatorname{Pro}(\mathcal{C})$  as  $\operatorname{pro-objects}$  of  $\mathcal{C}$ , and to  $\operatorname{Pro}(\mathcal{C})$  as  $\operatorname{the}$  category of  $\operatorname{pro-objects}$  of  $\mathcal{C}$ .

**Remark 2.** Let  $\mathcal{C}$  be a category which admits finite limits. For each object  $C \in \mathcal{C}$ , the functor  $D \mapsto \operatorname{Hom}_{\mathcal{C}}(C,D)$  preserves finite limits, and can therefore be regarded as an object of  $\operatorname{Pro}(\mathcal{C})$ . The Yoneda embedding  $C \mapsto \operatorname{Hom}_{\mathcal{C}}(C,\bullet)$  induces a fully faithful functor  $\mathcal{C} \to \operatorname{Pro}(\mathcal{C})$ . In what follows, we will generally abuse notation by identifying  $\mathcal{C}$  with its essential image in  $\operatorname{Pro}(\mathcal{C})$ .

**Remark 3.** In the category of sets, the formation of finite limits commutes with filtered colimits. It follows that the full subcategory  $\operatorname{Fun}^{\operatorname{lex}}(\mathcal{C},\operatorname{Set})$  is closed under filtered colimits in  $\operatorname{Fun}(\mathcal{C},\operatorname{Set})$ . In particular, the category  $\operatorname{Fun}^{\operatorname{lex}}(\mathcal{C},\operatorname{Set})$  admits filtered colimits, so that  $\operatorname{Pro}(\mathcal{C}) = \operatorname{Fun}^{\operatorname{lex}}(\mathcal{C},\operatorname{Set})^{\operatorname{op}}$  admits filtered limits.

**Example 4.** Let  $\mathcal{C}$  be a small category which admits finite limits. Suppose we are given a diagram  $\{C_{\alpha}\}_{{\alpha}\in\mathcal{A}^{\mathrm{op}}}$  indexed by (the opposite of) a filtered category  $\mathcal{A}$ . Then we can also regard  $\{C_{\alpha}\}$  as a diagram in the category  $\mathrm{Pro}(\mathcal{C})$  (via the Yoneda embedding), where we can take the inverse limit. We will abuse notation by denoting this inverse limit also by  $\{C_{\alpha}\}$ . When viewed as a functor from  $\mathcal{C}$  to the category of sets, it is given by the construction  $D \mapsto \varinjlim_{\alpha} \mathrm{Hom}_{\mathcal{C}}(C_{\alpha}, D)$ .

**Remark 5.** Let  $\mathcal{C}$  be an essentially small category which admits finite limits, and let  $F:\mathcal{C}\to \mathcal{S}$ et be a functor. Then F has a canonical presentation

$$\varinjlim_{(C,\eta)\in\mathcal{A}}\mathrm{Hom}_{\mathfrak{C}}(C,\bullet)$$

has a colimit of corepresentable functors, indexed by the category  $\mathcal{A}$  whose objects are pairs  $(C, \eta)$  where  $C \in \mathcal{C}$  and  $\eta \in F(C)$ , where

$$\operatorname{Hom}_{\mathcal{A}}((C, \eta), (C', \eta')) = \{ f \in \operatorname{Hom}_{\mathcal{C}}(C', C) : F(f)(\eta') = \eta \}.$$

If the functor F preserves finite limits, then the category  $\mathcal{A}$  is filtered. It follows that every object of  $Pro(\mathcal{C})$  has a (canonical) presentation as a filtered limit of objects of  $\mathcal{C}$ .

**Remark 6.** Let  $\mathcal{C}$  be an essentially small category which admits finite limits. From the above discussion, we see that the category  $Pro(\mathcal{C})$  can be described more informally as follows:

- The objects of Pro( $\mathcal{C}$ ) are diagrams  $\{C_{\alpha}\}$  in  $\mathcal{C}$ , indexed by (the opposite of) a small filtered category.
- Given two such diagrams  $\{C_{\alpha}\}\$  and  $\{D_{\beta}\}\$ , we have

$$\operatorname{Hom}_{\operatorname{Pro}(\mathfrak{C})}(\{C_{\alpha}\},\{D_{\beta}\}) = \varprojlim_{\beta} \operatorname{Hom}_{\operatorname{Pro}(\mathfrak{C})}(\{C_{\alpha}\},D_{\beta}) = \varprojlim_{\beta} \varinjlim_{\alpha} \operatorname{Hom}_{\mathfrak{C}}(C_{\alpha},D_{\beta}).$$

Remark 7. The category  $\operatorname{Pro}(\mathcal{C})$  can be characterized by a universal property. Let  $\mathcal{D}$  be any category which admits small filtered limits, and let  $\operatorname{Fun}'(\operatorname{Pro}(\mathcal{C}), \mathcal{D})$  be the full subcategory of  $\operatorname{Fun}(\operatorname{Pro}(\mathcal{C}), \mathcal{D})$  spanned by those functors which preserve small filtered limits. Then composition with the inclusion  $\mathcal{C} \hookrightarrow \operatorname{Pro}(\mathcal{C})$  induces an equivalence of categories  $\operatorname{Fun}'(\operatorname{Pro}(\mathcal{C}), \mathcal{D}) \to \operatorname{Fun}(\mathcal{C}, \mathcal{D})$ . In other words, every functor  $f: \mathcal{C} \to \mathcal{D}$  admits an essentially unique extension to a functor  $F: \operatorname{Pro}(\mathcal{C}) \to \mathcal{D}$  which preserves small filtered limits.

Let  $\mathcal{C}$  be an essentially small category which admits finite limits and let  $\mathcal{I}$  be any small category. Since  $\operatorname{Pro}(\mathcal{C})$  admits small filtered limit, the functor category  $\operatorname{Fun}(\mathcal{I},\operatorname{Pro}(\mathcal{C}))$  also admits small filtered limits (which are computed pointwise). Consequently, the inclusion functor  $\operatorname{Fun}(\mathcal{I},\mathcal{C}) \hookrightarrow \operatorname{Fun}(\mathcal{I},\operatorname{Pro}(\mathcal{C}))$  admits an essentially unique extension to a functor

$$\operatorname{Pro}(\operatorname{Fun}(\mathfrak{I},\mathfrak{C})) \to \operatorname{Fun}(\mathfrak{I},\operatorname{Pro}(\mathfrak{C}))$$

which preserves small filtered limits. We will use the following standard result:

**Proposition 8.** Let C be an essentially small category which admits finite limits and let I be a finite poset. Then the map

$$\operatorname{Pro}(\operatorname{Fun}(I,\mathcal{C})) \to \operatorname{Fun}(I,\operatorname{Pro}(\mathcal{C}))$$

is an equivalence of categories. In particular, every diagram  $I \to \text{Pro}(\mathfrak{C})$  can be written as a filtered limit of diagrams  $I \to \mathfrak{C}$ .

**Example 9.** Applying Proposition 8 in the case  $I = \{0 < 1\}$ , we see that every morphism  $f : C \to D$  in  $Pro(\mathcal{C})$  can be obtained as the limit of a filtered diagram of morphisms  $\{f_{\alpha} : C_{\alpha} \to D_{\alpha}\}$  between objects of  $\mathcal{C}$ .

**Corollary 10.** Let C be an essentially small category which admits finite limits. Then the category Pro(C) admits finite limits. Moreover, the inclusion  $C \hookrightarrow Pro(C)$  preserves finite limits.

*Proof.* Let  $\{C_i\}_{i\in I}$  be a finite diagram in  $\mathcal{C}$  having a limit  $C\in\mathcal{C}$ , and let  $\{D_\alpha\}$  be a filtered diagram in  $\mathcal{C}$  which we identify with an object of  $Pro(\mathcal{C})$ . Then

$$\operatorname{Hom}_{\operatorname{Pro}(\mathfrak{C})}(\{D_{\alpha}\}, C) \simeq \varinjlim_{\alpha} \operatorname{Hom}_{\mathfrak{C}}(D_{\alpha}, C)$$

$$\simeq \varinjlim_{\alpha} \varprojlim_{i} \operatorname{Hom}_{\mathfrak{C}}(D_{\alpha}, C_{i})$$

$$\simeq \varprojlim_{i} \varinjlim_{\alpha} \operatorname{Hom}_{\mathfrak{C}}(D_{\alpha}, C_{i})$$

$$\simeq \varprojlim_{i} \operatorname{Hom}_{\operatorname{Pro}(\mathfrak{C})}(\{D_{\alpha}\}, C).$$

where we have invoked the fact that filtered colimits commute with finite limits in the category of sets. This proves that the inclusion  $\mathcal{C} \hookrightarrow \operatorname{Pro}(\mathcal{C})$  preserves finite limits. In particular,  $\operatorname{Pro}(\mathcal{C})$  has an initial object. To complete the proof, it will suffice to show that every diagram  $C \to D \leftarrow E$  in  $\operatorname{Pro}(\mathcal{C})$  admits a fiber product. Using Proposition 8, we can realize our diagram as a filtered limit of diagrams  $\{C_{\alpha} \to D_{\alpha} \leftarrow E_{\alpha}\}$  in  $\mathcal{C}$ . Then the filtered diagram  $\{C_{\alpha} \times_{D_{\alpha}} E_{\alpha}\}$  represents a fiber product  $C \times_D E$  in the category  $\operatorname{Pro}(\mathcal{C})$ .

We will be particularly interested in studying  $Pro(\mathcal{C})$  in the case where  $\mathcal{C}$  is a pretopos.

**Proposition 11.** Let  $\mathcal{C}$  be a category which admits finite limits. Assume that every morphism  $f: X \to Z$  in  $\mathcal{C}$  factors as a composition  $X \xrightarrow{g} Y \xrightarrow{h} Z$ , where g is an effective epimorphism and h is a monomorphism. Then every morphism in  $\operatorname{Pro}(\mathcal{C})$  factors as a composition  $X \xrightarrow{g} Y \xrightarrow{h} Z$ , where g is an effective epimorphism and h is a monomorphism.

Proof. Let  $f: X \to Z$  be a morphism in  $\operatorname{Pro}(\mathcal{C})$ , which we can realize as a filtered limit of morphisms  $\{f_\alpha: X_\alpha \to Z_\alpha\}$  in  $\mathcal{C}$ . Factor each  $f_\alpha$  as a composition  $X_\alpha \xrightarrow{g_\alpha} Y_\alpha \xrightarrow{h_\alpha} Z_\alpha$ , where  $g_\alpha$  is an effective epimorphism and  $h_\alpha$  is a monomorphism. This factorization is functorial, so we can regard  $Y = \{Y_\alpha\}$  as a pro-object of  $\mathcal{C}$  equipped with morphisms  $g: X \to Y$  and  $h: Y \to Z$  with  $f = h \circ g$ . Note that  $Y \times_Z Y \simeq \{Y_\alpha \times_{Z_\alpha} Y_\alpha\} \simeq \{Y_\alpha\} = Y$ , so that h is a monomorphism in  $\operatorname{Pro}(\mathcal{C})$ . We will complete the proof by showing that g is an effective epimorphism in  $\operatorname{Pro}(\mathcal{C})$ . For this, we wish to show that for each object  $C \in \operatorname{Pro}(\mathcal{C})$ , the diagram

$$\operatorname{Hom}_{\operatorname{Pro}(\mathcal{C})}(Y,C) \to \operatorname{Hom}_{\operatorname{Pro}(\mathcal{C})}(X,C) \rightrightarrows \operatorname{Hom}_{\operatorname{Pro}(\mathcal{C})}(X \times_Y X,C)$$

is an equalizer. Writing C as a filtered limit of objects of C, we can assume that  $C \in C$ . In this case, the diagram above is given by a filtered colimit of diagrams

$$\operatorname{Hom}_{\mathfrak{C}}(Y_{\alpha}, C) \to \operatorname{Hom}_{\mathfrak{C}}(X_{\alpha}, C) \rightrightarrows \operatorname{Hom}_{\mathfrak{C}}(X_{\alpha} \times_{Y_{\alpha}} X_{\alpha}, C).$$

We conclude by observing that each of these diagrams is an equalizer (since  $g_{\alpha}$  is an effective epimorphism in  $\mathcal{C}$ ), and the collection of equalizer diagrams in Set is closed under filtered colimits.

**Remark 12.** The proof of Proposition 11 shows that a morphism  $f: X \to Y$  in Pro( $\mathcal{C}$ ) is a monomorphism (effective epimorphism) if and only if it can be realized as a filtered limit of morphisms  $\{f_{\alpha}: X_{\alpha} \to Y_{\alpha}\}$  which are monomorphisms (effective epimorphisms) in  $\mathcal{C}$ .

**Remark 13.** In the situation of Proposition 11, suppose that the formation of images in  $\mathcal{C}$  is compatible with pullback (or equivalently, the collection of effective epimorphisms is stable under pullback). Then the category  $\text{Pro}(\mathcal{C})$  has the same property: any diagram  $X \xrightarrow{f} Y \xleftarrow{g} Z$  can be realized as a filtered limit of diagrams  $\{X_{\alpha} \xrightarrow{f_{\alpha}} Y_{\alpha} \xleftarrow{g_{\alpha}} Z_{\alpha}\}$ , in which case we have

$$\begin{array}{rcl} \operatorname{Im}(X\times_YZ\to Z) & \simeq & \{\operatorname{Im}(X_\alpha\times_{Y_\alpha}Z_\alpha\to Z_\alpha)\\ & \simeq & \{\operatorname{Im}(X_\alpha\to Y_\alpha)\times_{Y_\alpha}Z_\alpha\}\\ & \simeq & \operatorname{Im}(X\to Y)\times_YZ. \end{array}$$

Let  $\mathcal{C}$  be an essentially small category which admits finite limits. Then  $\operatorname{Fun}^{\operatorname{lex}}(\mathcal{C},\operatorname{Set})$  is closed under limits in Set, and therefore admits small limits. It follows that the category  $\operatorname{Pro}(\mathcal{C})$  admits small colimits. Moreover, the inclusion functor  $\mathcal{C} \hookrightarrow \operatorname{Pro}(\mathcal{C})$  preserves all colimits which exist in  $\mathcal{C}$  (this is immediate from the definitions).

**Proposition 14.** Let C be an essentially small category which admits finite limits and finite coproducts. Then the category Pro(C) admits finite coproducts, given by the formula

$${C_{\alpha}}\coprod {D_{\beta}} = {C_{\alpha}\coprod D_{\beta}}.$$

*Proof.* It suffices to observe that for any object  $E \in \mathcal{C}$ , we have

$$\varinjlim_{\alpha,\beta} \operatorname{Hom}_{\mathfrak{C}}(C_{\alpha} \coprod D_{\beta}, E) \simeq (\varinjlim_{\alpha} \operatorname{Hom}_{\mathfrak{C}}(C_{\alpha}, E)) \times (\varinjlim_{\beta} \operatorname{Hom}_{\mathfrak{C}}(D_{\beta}, E)).$$

Given objects  $C, D \in \text{Pro}(\mathcal{C})$ , we can use Proposition 8 to write  $C = \{C_{\alpha}\}$  and  $D = \{D_{\alpha}\}$  as limits of diagrams indexed by the same category. In this case, the coproduct  $C \coprod D$  is given by  $\{C_{\alpha} \coprod D_{\alpha}\}$ .

**Remark 15.** In the situation of Proposition 14, suppose that the formation of coproducts in  $\mathcal{C}$  is preserved by pullback. Then the same is true in  $\operatorname{Pro}(\mathcal{C})$ . Given morphisms  $f: C \to X$ ,  $g: D \to X$ , and  $h: Y \to X$  in  $\operatorname{Pro}(\mathcal{C})$ , we can apply Proposition 8 to realize f, g, and h as filtered limits of maps  $f_{\alpha}: C_{\alpha} \to X_{\alpha}$ ,  $g_{\alpha}: D_{\alpha} \to X_{\alpha}$ , and  $h_{\alpha}: Y_{\alpha} \to X_{\alpha}$  (indexed by the same category), so that both  $(C \coprod D) \times_X Y$  and  $(C \times_X Y) \coprod (D \times_X Y)$  are represented by the diagram

$$\{(C_\alpha \amalg D_\alpha) \times_{X_\alpha} Y_\alpha\} \simeq \{(C_\alpha \times_{X_\alpha} Y_\alpha) \amalg (D_\alpha \times_{X_\alpha} Y_\alpha)\}.$$

**Remark 16.** In the situation of Proposition 14, suppose that coproducts in  $\mathcal{C}$  are disjoint. Then, for every pair of objects  $C = \{C_{\alpha}\}$  and  $D = \{D_{\alpha}\}$  in  $\mathcal{C}$ , we deduce that

$$C \times_{C \coprod D} D \simeq \{C_{\alpha} \coprod_{C_{\alpha} \coprod D_{\alpha}} D_{\alpha}\} = \{\emptyset\}$$

is an initial object of  $Pro(\mathcal{C})$ : that is, coproducts are disjoint in  $Pro(\mathcal{C})$ .

Combining the above results, we obtain the following:

**Proposition 17.** Let C be an essentially small coherent category with disjoint coproducts (for example, a pretopos). Then Pro(C) is also a coherent category with disjoint coproducts.

Warning 18. It is not true that if  $\mathcal{C}$  is a pretopos, then  $Pro(\mathcal{C})$  is also a pretopos. For example, let  $\mathcal{C}$  be the category of finite sets. Then the category  $Pro(\mathcal{C})$  of profinite sets can be identified with the category of Stone spaces: that is, the category whose objects are totally disconnected compact Hausdorff spaces, and whose morphisms are continuous maps. Let  $C \in Pro(\mathcal{C})$  be the Cantor set, which we identify with the collection of infinite sequences  $(n_1, n_2, n_3, \ldots)$  where  $n_i \in \{0, 1\}$ . The construction

$$(n_1, n_2, n_3, \ldots) \mapsto \sum \frac{n_i}{2^i}$$

defines a continuous surjection  $C \to [0,1]$ , and the fiber product  $R = C \times_{[0,1]} C$  can be regarded as an equivalence relation on C in the category of Stone spaces. However, this equivalence relation is *not* effective: given any Stone space X, a continuous map  $C \to X$  which equalizes the two projection maps  $R \rightrightarrows C$  must factor through a continuous map  $[0,1] \to X$ . Such a map is automatically constant (since X is totally disconnected), so that  $C \times_X C = C \times C$  is larger than R.