## Math 155 (Lecture 8)

## September 19, 2011

**Definition 1.** Let G be a graph. A path in G is a sequence of vertices  $v_0, \ldots, v_n$  such that  $v_i$  is adjacent to  $v_{i-1}$  for  $1 \le i \le n$ . In this case, we also say  $(v_0, \ldots, v_n)$  is a path from  $v_0$  to  $v_n$ . We say that a path is simple if the vertices  $v_i$  are disjoint. We say that a simple path  $(v_0, \ldots, v_n)$  is a cycle if  $n \ge 2$  and  $v_0$  is adjacent to  $v_n$ .

We say that a graph G is *connected* if every pair of vertices can be connected by a path. We say that G is a *tree* if it is connected and contains no cycles.

The problem we will consider in this lecture is the following:

Question 2. How many trees are there having the set of vertices  $\{1, \ldots, n\}$ ?

The answer turns out to be  $n^{n-2}$ . There are many different proofs of this result. In this lecture, we will give a proof using the composition formula for exponential generating functions. First, it is convenient to introduce a slight variation on Question 2.

Question 3. How many ways are there to choose a tree having the set of vertices  $\{1, \ldots, n\}$ , together with an ordered pair of (possibly identical) vertices of the tree?

It is clear that the answer to Question 3 differs from the answer to Question ?? by a factor of  $n^2$ . We will show that the answer to Question 3 is  $n^n$ . We already know another counting problem whose answer is  $n^n$ : namely, it is the number of maps from the set  $\{1, 2, ..., n\}$  to itself. To discuss these counting problems in more detail, it is convenient to introduce some species:

**Definition 4.** Let  $S_{\text{End}}$  denote the *species of endomorphisms*: for each finite set I,  $S_{\text{End}}[I]$  is the set of all maps from I to itself.

We define a species  $S_{2-\text{tree}}$  as follows:

- If I is nonempty, then  $S_{2-\text{tree}}[I]$  is the collection of all triples (T, i, j) where i and j are elements of I, and T is a tree with vertex set I.
- If  $I = \emptyset$ , then  $S_{2-\text{tree}}[I] = \{*\}.$

We would like to show that for every nonempty finite set I, the sets  $S_{\text{End}}[I]$  and  $S_{2-\text{tree}}[I]$  have the same number of elements (our ad-hoc definition of  $S_{2-\text{tree}}[\emptyset]$  was made to ensure that this is also correct when I is empty). Equivalently, we would like to show that the exponential generating functions of  $S_{\text{End}}$  and  $S_{2-\text{tree}}$ coincide.

Warning 5. The species  $S_{\text{End}}$  and  $S_{2-\text{tree}}$  are *not* isomorphic. For example, when  $I = \{1, 2\}$ , then the sets  $S_{\text{End}}[I]$  and  $S_{2-\text{tree}}[I]$  are acted on by the symmetric group  $\Sigma_2$ . This action has no fixed points on  $S_{2-\text{tree}}[I]$ , but has two fixed points on  $S_{\text{End}}[I]$ . This is what makes Question 2 interesting: though it has a simple answer, there is no obvious bijective approach.

Let's begin by analyzing the species  $S_{\text{End}}$ . Suppose that I is a finite set and we are given a map  $\pi : I \to I$ . What could I look like? It might be a permutation of I. Let  $S_{\text{perm}}$  denote the species of permutations, so that  $S_{\text{perm}}[I]$  is the set of permutations of I. However, there are many maps from I to itself which are not permutations. For example, at the other extreme, there are constant maps from I to itself. **Definition 6.** Let  $\pi$  be a map from a finite set I to itself. We will say that  $\pi$  is *nilpotent* if some power of  $\pi$  is a constant map. In other words,  $\pi$  is nilpotent if there exists an element  $i \in I$  such that for each  $j \in I$ , we have  $\pi^n(j) = i$  for all sufficiently large n. We let  $S_{nil}$  denote the *species of nilpotent endomorphisms*: for every finite set I, we let  $S_{nil}[I]$  denote the set of all nilpotent maps from I to itself.

**Remark 7.** If  $\pi : I \to I$  is a nilpotent map, then there is a unique element  $i \in I$  such that  $\pi^n(j) = i$  for  $n \gg 0$ . We will call *i* the *attractor* of  $\pi$ .

Note that if I has more than one element, then no map  $\pi : I \to I$  can be both a permutation and nilpotent (otherwise, some power of  $\pi$  would be a constant permutation). We now show that, in some sense, every map is made out of nilpotent maps and permutations.

**Proposition 8.** There is an isomorphism of species  $S_{\text{End}} \simeq S_{\text{perm}} \circ S_{\text{nil}}$ .

*Proof.* Let I be a finite set. By definition, an element of  $S_{\text{perm}}$  corresponds to the following data:

- (a) An equivalence relation  $\sim$  on the set I.
- (b) A permutation  $\sigma$  of the set  $I/\sim$  of equivalence classes.
- (c) For each equivalence class  $J \subseteq I$ , a nilpotent map  $\tau_J : J \to J$ .

Let us first describe how, given this data, we can construct a map  $\pi : I \to I$ . Each of the nilpotent maps  $\tau_J$  has an attractor, which we will denote by  $i_J$ . Let  $I_0 \subseteq I$  be the collection of all these attractors. The composition  $I_0 \to I \to I/\sim$  is a bijection (that is,  $I_0$  contains exactly one element of each equivalence class), so we may regard  $\sigma$  as a permutation of the set  $I_0$ . We now set

$$\pi(i) = \begin{cases} \sigma(i) & \text{if } i \in I_0\\ \tau_J(i) & \text{if } i \in J - \{i_J\} \end{cases}$$

Conversely, suppose that we start with an arbitrary map  $\pi : I \to I$ . We will say that an element  $i \in I$  is *periodic* if there exists an integer n > 0 such that  $\pi^n(i) = i$ . Let  $I_0$  denote the collection of periodic elements of I. The map  $\pi$  carries periodic elements to periodic elements (note that if  $\pi^n(i) = i$ , then  $\pi^n \pi(i) = \pi^{n+1}(i) = \pi(i)$ ), and restricts to a permutation on the subset  $I_0 \subseteq I$ . For each element  $i \in I$ , the sequence

$$i, \pi(i), \pi^2(i), \ldots$$

must eventually repeat (since I is finite), and therefore contains a periodic element of I. We will denote the first periodic element of the sequence by r(i). The map r determines an equivalence relation on I: let us write  $i \sim j$  if r(i) = r(j). Note that each equivalence class contains a unique element of  $I_0$ . Let J be an arbitrary equivalence class containing an element  $i_J \in I_0$ . For each element  $j \in I_0$ , we have  $\pi^n(j) = i_J$ for some integer  $n \geq 0$ , and  $\pi^m(j) \notin I_0$  for m < n. If  $j \neq i_J$ , then n > 0, so that  $\pi^{n-1}(\pi(j)) = i_J$  (and  $\pi^m(\pi_j) \notin I_0$  for m < n-1). It follows that  $\pi(j) \in J$ , so that  $\pi$  determines a map

$$\tau_J: J - \{i_J\} \to J.$$

This extends to a nilpotent map from J to itself, if we set  $\tau_J(i_J) = i_J$ . Then the triple  $(\sim, \sigma, \{\tau_J\}_{J \in I/\sim})$  is an element of  $(S_{\text{perm}} \circ S_{\text{nil}})[I]$ .

The above constructions determine maps of species

$$S_{\text{perm}} \circ S_{\text{nil}} \to S_{\text{End}}$$
  
 $S_{\text{perm}} \circ S_{\text{nil}} \leftarrow S_{\text{End}}.$ 

We leave it to the reader as an exercise to show that these maps really are inverse to one another, and therefore determine an isomorphism  $S_{\text{perm}} \circ S_{\text{nil}} \simeq S_{\text{End}}$ .

We would now like to relate the above discussion to trees. First, we need a few basic facts from graph theory.

**Lemma 9.** Let G be a tree containing vertices v and v'. Then there is a unique simple path  $v = v_0, v_1, v_2, \ldots, v_n = v'$  from v to v'.

*Proof.* Since G is connected, there is at least one path  $v = v_0, v_1, \ldots, v_n = v'$  from v to v'. The number n is called the *length* of the path; let us assume that n has been chosen as small as possible. We claim that this path is automatically simple: if  $v_i = v_j$  for i < j, then

$$v = v_0, \dots, v_{i-1}, v_i = v_j, v_{j+1}, \dots, v_n = v'$$

is a shorter path from v to v'. This proves existence.

Now we prove uniqueness. Choose a path  $v = v_0, \ldots, v_n = v'$  of minimal length. We will use induction on n to show that every other simple path  $v = w_0, \ldots, w_m = v'$  coincides with  $(v_0, \ldots, v_n)$ . The case n = 0is trivial: if v = v', then every simple path from v to v' is automatically of length 0 (since the vertices appearing in the path must be distinct).

By assumption, we have  $v_n = v' = w_m$ . Let *i* be the smallest positive integer such that  $v_i \in \{w_1, \ldots, w_m\}$ , and write  $v_i = w_j$ . Then the sequence

$$v_0, v_1, \ldots, v_i = w_j, w_{j-1}, \ldots, w_1$$

is a simple path from  $v_0$  to  $w_1$ . If i + j > 2, this is a cycle. We must therefore have i = j = 1. This proves that  $v_1 = w_1$ . We then have two simple paths  $(v_1, \ldots, v_n)$  and  $(w_1, \ldots, w_m)$  from  $v_1 = w_1$  to v', which must coincide by the inductive hypothesis.

**Definition 10.** A rooted tree is a tree T together with a choice of vertex  $r \in T$ , called the root of T. We let  $S_{1-\text{tree}}$  denote the species of rooted trees: for each finite set I, we let  $S_{1-\text{tree}}[V]$  denote the set of pairs (r, T), where  $r \in V$  and T is a tree with vertex set V.

**Proposition 11.** The species  $S_{1-\text{tree}}$  and  $S_{\text{nil}}$  are isomorphic.

*Proof.* Let V be a finite set. We will show that there is a canonical bijection between  $S_{1-\text{tree}}[V]$  and  $S_{\text{nil}}[V]$ . First, suppose we are given a tree T with vertex set V and a choice of root r. We define a function  $\pi: V \to V$  as follows. For each element  $v \neq r$  in V, let  $v = v_0, v_1, \ldots, v_n = r$  be the unique simple path from v to the root r. Then we set  $\pi(v) = v_1$ . If v = r, we let  $\pi(r) = r$ . It is easy to see that the function  $\pi$  is nilpotent (with attractor r).

Conversely, suppose that we are given a nilpotent function  $\pi : V \to V$ , and let r be the attractor of  $\pi$ . We can make V into a graph by declaring that a pair of distinct vertices v and w are adjacent if either  $v = \pi(w)$  or  $w = \pi(v)$ . We claim that this graph is always a tree. To prove this, suppose we are given a cycle

$$v_0, v_1, \ldots, v_n$$
.

For  $1 \le i \le n$ , we have either  $v_i = \pi(v_{i-1})$  or  $\pi(v_i) = v_{i-1}$ . Note that in the second case, since  $v_{i+1} \ne v_{i-1}$ , we must also have  $\pi(v_{i+1}) = v_i$ . It follows that the path is given by

$$v_0, \pi(v_0), \pi^2(v_0), \dots, \pi^p(v_0) = w_0, w_1, w_2, \dots, w_q$$

where p + q = n and the sequence  $\{w_i\}$  satisfies  $\pi(w_i) = w_{i-1}$ . If the path is a cycle, then either  $\pi(v_0) = w_q$ or  $\pi(w_q) = v_0$ . In the first case, we must have p = 0 (otherwise  $v_1 = v_n$ ), so that  $\pi^{q+1}w_q = w_q$ , which implies that  $w_q = r$  and therefore q = 0, a contradiction. In the second case we must have q = 0 (otherwise  $v_0 = v_{n-1}$ ) in which case we have  $\pi^{p+1}v_0 = v_0$ , which implies that  $v_0 = r$  so that p = 0, again a contradiction.

We again leave it to the reader to show that these constructions are inverse to one another, and give an isomorphism of species  $S_{1-\text{tree}} \simeq S_{\text{nil}}$ .

Combining Propositions 8 and 11, we get an isomorphism of species

$$S_{\text{End}} \simeq S_{\text{perm}} \circ S_{1-\text{tree}}.$$

To complete our analysis, we will need the following analogous fact:

**Proposition 12.** Let  $S_{\text{lin}}$  denote the species of linear orderings. There is an isomorphism of species  $S_{2-\text{tree}} \simeq S_{\text{lin}} \circ S_{1-\text{tree}}$ .

Assuming Proposition 12 for the moment and using the composition formula, we obtain

$$F_{S_{\text{End}}}(x) = F_{S_{\text{perm}}}(F_{S_{1-\text{tree}}}(x))$$
  
$$F_{S_{2-\text{tree}}}(x) = F_{S_{\text{lin}}}(F_{S_{1-\text{tree}}}(x)).$$

We have already seen that, although the species  $S_{\text{perm}}$  and  $S_{\text{lin}}$  are different, they have the same exponential generating function  $\sum_{n>0} \frac{n!}{n!} x^n = \frac{1}{1-x}$ . We therefore obtain

$$F_{S_{\text{End}}}(x) = \frac{1}{1 - F_{S_{1-\text{tree}}}(x)} = F_{S_{2-\text{tree}}}(x).$$

This proves that the number of elements of  $S_{2-\text{tree}}[I]$  is the same as the number of elements of  $S_{\text{End}}[I]$  for every finite set I, thereby giving  $n^{n-2}$  as the answer to Question 2.

Let us now prove Proposition 12. For each finite set V, we must construct a bijection from the set  $S_{2-\text{tree}}[V]$  to the set  $(S_{\text{lin}} \circ S_{1-\text{tree}})[V]$ . It is clear what to do if V is empty (in that case, both sides have one element). Let us therefore assume that V is nonempty, so that an element of  $S_{2-\text{tree}}[V]$  corresponds to a triple (T, v, v'), where T is a tree with vertex set V, and  $v, v' \in V$  are elements. To this, we wish to associate the following data:

- (a) An equivalence relation  $\sim$  on the set V.
- (b) A linear ordering of the set of equivalence classes  $V/\sim$ .
- (c) For each equivalence class  $W \subseteq V$ , a rooted tree with vertex set W.

We first invoke Lemma 9: since T is a tree, there is a unique simple path

$$v = v_0, v_1, \dots, v_n = v'$$

from v to v'. Let  $V_0 = \{v_0, \ldots, v_n\}$  be the set of vertices along this path. For any vertex  $w \in V$ , there exists a simple path  $w = w_0, \ldots, w_m$  which ends in a vertex  $w_m$  which belongs to  $V_0$ . In fact, we can assume that  $w_i \notin V_0$  for i < m (otherwise, end the path at  $w_i$  instead). In this case, we claim that the path  $w_0, \ldots, w_m$ is unique. Suppose otherwise: that is, suppose that we have another simple path  $w = w'_0, w'_1, \ldots, w'_{m'}$ with  $w'_{m'} \in V_0$  and  $w'_j \notin V_0$  for j < m'. If  $w'_{m'} = w_m$ , this contradicts Lemma 9. Otherwise, we may suppose that  $w_m = v_i$  and  $w'_{m'} = v_j$  for i < j. Let k be the largest integer such that  $w_k$  appears in the path  $w'_0, \ldots, w'_{m'}$  (such an integer exists, since  $w_0 = w = w'_0$ ). Writing  $w_k = w'_{k'}$ , we obtain a cycle  $w_k, w_{k+1}, \ldots, w_m = v_i, v_{i+1}, \ldots, v_j = w'_{m'}, w'_{m'-1}, \ldots, w'_{k'+1}$ , contradicting our assumption that T is a tree.

For each  $w \in V$ , let  $\tau(w) \in V_0$  be the endpoint of the unique path constructed above. Let us define an equivalence relation on V by writing  $w \sim w'$  if  $\tau(w) = \tau(w')$ . Then the composition  $V_0 \to V \to V/ \sim$  is a bijection: that is,  $V_0$  contains a unique element of each equivalence class. Writing  $V_0 = \{v_0, v_1, \ldots, v_n\}$ , we obtain a linear ordering on the set  $V/ \sim$ . Each equivalence class  $W \in V/ \sim$  has the form  $\tau^{-1}\{v_i\}$  for some  $0 \leq i \leq n$ . Note that if  $w \in \tau^{-1}v_i$ , there is a simple path  $w = w_0, \ldots, w_m = v_i$  from w to  $v_i$  which contains no other vertex of  $V_0$ . Each of the paths  $w_j, w_{j+1}, \ldots, w_m$  has the same property, so that  $\tau(w_j) = v_i$  for  $0 \leq j \leq m$ . It follows that w can be connected to  $v_i$  by a path lying entirely in W, so that W determines a connected subgraph of T. This subgraph cannot contain any cycles (since T does not contain any cycles) and is therefore a tree (which has a natural choice of root, given by the vertex  $v_i \in V_0$ ).

We leave it to the reader to verify that this construction determines a bijection of  $S_{2-\text{tree}}[V]$  with  $(S_{\text{lin}} \circ S_{1-\text{tree}})[V]$  (if it is not clear, try drawing a picture for yourself).