

abstract

Members Seminar
Topic:

Speaker:

Affiliation:

Date:

Time/Room:

I will give an introduction to the problem of parallel repetition of two-prover games and its applications and related results in theoretical computer science (the PCP theorem, hardness of approximation), mathematics (the geometry of foams, tiling the space \mathbb{R}^n) and, if time allows, physics (Bell inequalities, the EPR paradox).

In a two-prover (alternatively, two-player) game, a referee chooses questions (x,y) according to a (publicly known) distribution, and sends x to the first player and y to the second player. The first player responds by $a=a(x)$ and the second by $b=b(y)$ (without communicating with each other). The players jointly win if a (publicly known) predicate $V(x,y,a,b)$ holds. The value of the game is the maximal probability of success that the players can achieve, where the maximum is taken over all protocols $a=a(x)$, $b=b(y)$.

A parallel repetition of a two-prover game is a game where the players try to win n copies of the original game simultaneously. More precisely, the referee generates questions $x=(x_1,\dots,x_n)$, $y=(y_1,\dots,y_n)$, where each pair (x_i,y_i) is chosen independently according to the original distribution. The players respond by $a=(a_1,\dots,a_n)$ and $b=(b_1,\dots,b_n)$. The players win if they win simultaneously on all the coordinates, that is, if for every i , $V(x_i,y_i,a_i,b_i)$ holds.

The parallel repetition theorem states that for any two-prover game with value smaller than 1, parallel repetition reduces the value of the game in an exponential rate. Formally, for any two-prover game with value $1-\epsilon$ (for, say, $\epsilon < 1/2$), the value of the game repeated in parallel n times is at most $(1-\epsilon^3)^{\Omega(n/s)}$, where s is the answers' length (of the original game).

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I will discuss applications of the parallel repetition theorem and related results in theoretical computer science, mathematics, and, if time allows, physics.