

## abstract

Computer Science/Discrete Mathematics Seminar II  
Topic:

Speaker:

Affiliation:

Date:

Time/Room:

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The notion of exactly (or approximately) representing certain combinatorial properties of a graph  $G$  on a simpler graph is ubiquitous in combinatorial optimization. In this talk, I will introduce the notion of vertex sparsification. Here we are given a graph  $G = (V, E)$  and a set of terminals  $K \subset V$  and our goal is to find one single graph  $H = (K, E_H)$  on just the terminal set so that  $H$  approximately preserves the minimum cut between every bi-partition of the terminals. Standard results in combinatorial optimization are concerned with minimum cuts between "small" subsets of the terminals - .e.g. Mader's theorem implies that if we are interested in the minimum cuts between each pair of terminals, there is a graph  $H$  on just the set of terminals that exactly represents these  $\binom{K}{2}$  values. Yet in our setting, we are interested in minimum cuts between "large" subsets of terminals as well - e.g. the minimum bisection of the terminals. One might wonder whether good vertex sparsifiers exist. There are two orthogonal challenges that we must overcome - first, the minimum cuts that we are interested contain the answers to "hard" optimization problems and second, we must approximately preserve exponentially many values using only  $\binom{K}{2}$  degrees of freedom. Nevertheless, I will prove that good vertex sparsifiers do exist and in fact the approximation factor is independent of the size of the original graph (and is sub-logarithmic in the number of terminals). Moreover, the approximation factor can be improved to a constant for naturally arising graphs - namely those that exclude any fixed minor. I will give a number of results in this context which will give me an excuse to build connections to metric embeddings, fourier analysis and linear programming hierarchies. Lastly, I will also give a number of applications (all based on the pattern of using a good vertex sparsifier as a proxy for the original graph) including a master theorem for flow-cut gaps. Parts of this talk will be based on joint work with Tom Leighton, and Moses Charikar and Shi Li.

<http://math.ias.edu/files/seminars/Moitra.pdf>

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