

## **abstract**

COMPUTER SCIENCE AND DISCRETE MATHEMATICS SEMINAR I

Topic:

Speaker:

Affiliation:

Date:

Time/Room:

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A "sparsifier" of a graph is a weighted subgraph for which every cut has approximately the same value as the original graph, up to a factor of  $(1 \pm \epsilon)$ . Sparsifiers were first studied by Benczur and Karger (1996). They have wide-ranging applications, including fast network flow algorithms, fast linear system solvers, etc. Batson, Spielman and Srivastava (2009) showed that sparsifiers with  $O(n/\epsilon^2)$  edges exist, and they can be computed in time  $\text{poly}(n, \epsilon)$ .

We describe two new approaches to constructing sparsifiers. The first approach independently samples each edge  $uv$  with probability inversely proportional to the edge-connectivity between  $u$  and  $v$ . The fact that this approach produces a sparsifier resolves an open question of Benczur and Karger. The second approach samples uniformly random spanning trees. Both of our approaches produce sparsifiers with  $O(n \log^2(n) / \epsilon^2)$  edges. Our proofs are based on extensions of the Karger-Stein contraction algorithm which allow it to compute minimum "Steiner" cuts.