

## **abstract**

COMPUTER SCIENCE AND DISCRETE MATHEMATICS SEMINAR I

Topic:

Speaker:

Affiliation:

Date:

Time/Room:

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The Johnson-Lindenstrauss lemma states that for any  $n$  points in Euclidean space and error parameter  $0 < \epsilon < 1/2$ , there exists an embedding into  $k = O(\epsilon^{-2} \cdot \log n)$  dimensional Euclidean space so that all pairwise distances are preserved up to a  $1 + \epsilon$  factor. This lemma has applications in high-dimensional computational geometry (decreasing dimension makes many algorithms run faster), compressed sensing, and numerical linear algebra.

All known proofs of the lemma construct a distribution over linear mappings so that a random such mapping suffices with high probability. In this talk, I will present various proofs of the JL lemma satisfying, for example, (1) the support of the distribution is small, so that a random embedding can be selected with few random bits (e.g.  $O(\log n \log \log n)$  bits for constant  $\epsilon$ , which is suboptimal by a  $\log \log n$  factor), and (2) every embedding matrix in the support of the distribution is sparse (only  $O(\epsilon \cdot k)$  entries per column are non-zero), to speed up computation. I will also describe some open problems. This talk is based on joint works with Daniel Kane (Harvard).