

abstract

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Topic:

Speaker:

Affiliation:

Date:

Time/Room:

An epsilon-biased set X in $\{0,1\}^n$ is a set so that for every non-empty set T in $[n]$ the following holds. The random bit $B(T)$ obtained by selecting at random a vector x in X , and computing the mod-2 sum of its T -coordinates, has bias at most epsilon. Such sets may be viewed as generating matrices of binary error correcting codes of distance $(1/2 - \epsilon)$, as well as pseudorandom sets in the sense that all their nontrivial Fourier coefficients are at most epsilon in absolute value. Determining the smallest size of such an epsilon-biased sets X , and explicitly constructing them has a long history.

We shall survey three constructions of epsilon-bias sets, each is better in some range of parameters. The first construction is by Alon, Goldreich, Hastad and Peralta and gives an explicit epsilon-biased set X of size at most order $(n/\epsilon)^2$. The second construction, of Alon, Bruck, Naor, Naor and Roth, gives a set X of size at most roughly n/ϵ^3 . The third construction is by Ben-Aroya and Ta-Shma and gives roughly $(n/\epsilon^2)^{5/4}$ for epsilon at least $n^{-0.5}$ (which is better than previous constructions for $\epsilon = 1/n$, for example). The last two constructions (can) use algebraic-geometric codes. If time permits, we shall discuss the connection to algebraic geometry codes in more detail.