

abstract

COMPUTER SCIENCE/DISCRETE MATH I

Topic:

Speaker:

Affiliation:

Date:

Time/Room:

A given function $f(x)$ is "hard on average" with respect to an algorithm A , if $A(x)$ disagrees with $f(x)$ on "many" inputs x . Applications in cryptography and derandomization require functions that are "very hard on average" (essentially unpredictable) with respect to any efficient algorithm. An algorithmic procedure of converting a "somewhat hard on average" function $f(x)$ into a "much harder on average" function $g(x)$ is called hardness amplification.

A primary tool for hardness amplification is the classical Direct Product Theorem that essentially says the following: If a function $f(x)$ is "somewhat hard on average", then the function $f^k(x_1, \dots, x_k) = f(x_1) \dots f(x_k)$ (required to compute the value of f on each of k independent inputs x_1, \dots, x_k) is "much harder on average", where the amount of hardness increases exponentially fast with the parameter k . In the language of error-correcting codes this

basically means that the truth table of the direct-product function f^k can "tolerate" many more corruptions than the truth table of f . Of special interest is an efficient error-correcting (decoding) algorithm for such "Direct Product Codes".

Our main result is a simple, efficient decoding algorithm for Direct Product codes, which achieves information-theoretically optimal parameters (up to constant factors); thus it significantly improves on an earlier result of [Impagliazzo, Jaiswal, and Kabanets; FOCS'06]. We also define a more general class of "direct product"-like codes with efficient decoding algorithms, which in particular yields a certain "derandomized" version of the Direct Product Theorem.

Joint work with Russell Impagliazzo (UCSD & IAS), Ragesh Jaiswal (UCSD), and Avi Wigderson

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